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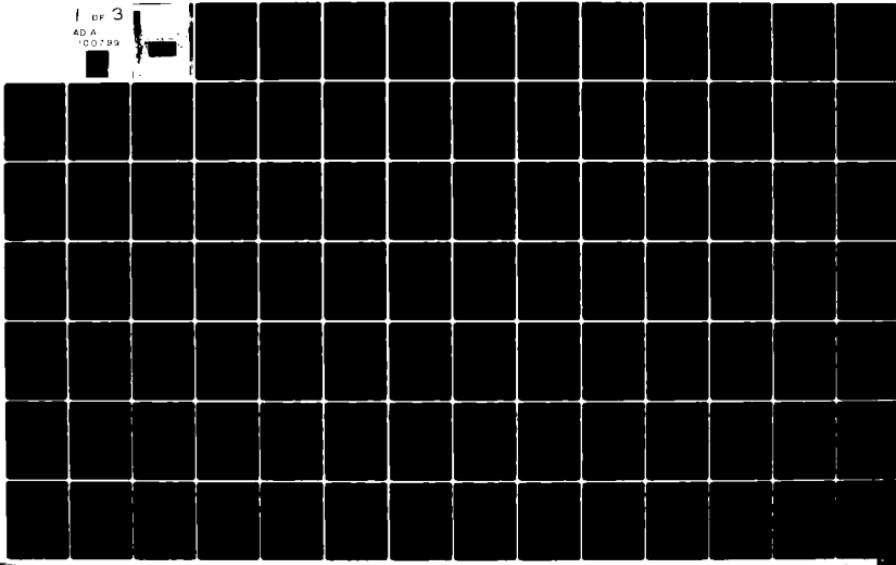
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ANALYSIS AND PERFORMANCE EVALUATION

OF ELECTROCARDIOGRAM

DATA COMPRESSION TECHNIQUES

THESIS

AFIT/GE/EE/80D-46 / Melvin D. Townsend  
Captain USAF

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AFIT/GE/EE/80D-46

ANALYSIS AND PERFORMANCE EVALUATION  
OF ELECTROCARDIOGRAM  
DATA COMPRESSION TECHNIQUES

THESIS

Presented to the Faculty of the School of Engineering  
of the Air Force Institute of Technology  
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Master of Science

by

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Approved of public release; distribution unlimited.

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## Contents

	<b>Page</b>
Acknowledgements . . . . .	i
List of Figures . . . . .	iv
Abstract . . . . .	vi
I. Introduction . . . . .	1
Background . . . . .	1
Problem . . . . .	3
Scope . . . . .	3
General Approach . . . . .	4
Sequence of Presentation . . . . .	5
II. Data Compression and the EKG . . . . .	8
Entropy Reduction . . . . .	9
Sampling . . . . .	10
Orthogonal Expansions and Filtering . . . . .	11
EKG Filtering Compression . . . . .	13
Transforms . . . . .	14
Fast Fourier Transform . . . . .	15
Discrete Karhunen-Loeve Transform . . . . .	17
Redundancy Reduction . . . . .	23
Predictors/Interpolators . . . . .	24
Difference Reductions . . . . .	26
Time Compression . . . . .	27
Entropy Encoding . . . . .	29
Summary . . . . .	31
III. Tolan and Dower Compression Techniques . . . . .	33
Tolan EKG Data Compressor . . . . .	39
Tolan Decorrelator . . . . .	39
Tolan Entropy Encoder . . . . .	40
Dower EKG Data Compressor . . . . .	47
Dower Decorrelator . . . . .	48
Dower Entropy Encoder . . . . .	51
Tolan versus Dower . . . . .	57
Other Microcomputer EKG Compressors . . . . .	59
Turning Point Algorithm . . . . .	59
Marquette Algorithm . . . . .	61
Chapter III Summary . . . . .	63
IV. EKG-Data Acquisition and Analysis System . . . . .	66

EKG-DAAS Hardware . . . . .	66
EKG-DAAS Software . . . . .	69
EKG-EXEC . . . . .	74
NOCPRS . . . . .	76
DISPLAY . . . . .	78
TOLAN . . . . .	80
Summary . . . . .	83
V. Experimental Procedure, Data Analysis, and Results . . . . .	
Experimental Procedure . . . . .	84
EKG Compression Measurement Parameters . . . . .	87
Data Analysis and Results . . . . .	92
TOLAN and DOWER Performance Comparison . . . . .	98
Summary . . . . .	99
VI. Summary, Conclusions and Recomendations . . . . .	101
Bibliography . . . . .	104
Appendix A: The Electrocardiogram . . . . .	
The Physiology and Electrical Characteristics of the Heart . . . . .	108
The Cardiac Cycle . . . . .	111
Heart Disease and the EKG . . . . .	113
Appendix B: Fundamentals of Information Theory . . . . .	
Information Source . . . . .	118
Transmission Channel . . . . .	121
Encoder . . . . .	125
Decoder . . . . .	127
Appendix C: Software Program Listings . . . . .	
EKG-EXEC . . . . .	130
DISPLAY . . . . .	143
PRSTAT . . . . .	151
NOCPRS . . . . .	166
TOLAN-A . . . . .	176
DECPRS . . . . .	189
ENTROPY . . . . .	199
Appendix D: EKG Experimental Data Listing . . . . .	206
Appendix E: Equipment Specifications . . . . .	231
Vita . . . . .	252

## List of Figures

Figure	Page
1 Typical EKG Waveform with Arrhythmia . . . . .	13
2 Typical EKG Redundancy Reduction Data Compressor.	23
3 Second Order Difference Operator. . . . .	28
4 Relative Frequency of 2nd Order Difference Operator	28
5 Nonstationary Waveform With Random Level and Slope	34
6 EKG Waveform With Moving Average Slope . . . . .	35
7 Simulated Variance Reduction Via Second Difference	37
8 Tolan Collection and Decorrelation Algorithm . .	41
9 Tolan Code and Code Tree . . . . . . . . . . .	42
10 Tolan Variable Length Encoder Algorithm . . . . .	44
11 Dower Collection and Decorrelation Algorithm . .	45
12 Dower EKG Data Compression System . . . . .	47
13 Three Hypothetical Sample Sequences . . . . .	50
14 Example of State Space Partition for 4 Symbols .	52
15 Dower Encoder Accumulator-Memory Buffer Interface	53
16 Dower VLC Algorithm . . . . . . . . . . .	56
17 Turning Point Patterns . . . . . . . . . . .	60
18 EXORCISER Component Module Layout . . . . .	68
19 EKG-DAAS Hardware Configuration . . . . .	69
20 EKG-DAAS Software Control Flowgraph . . . . .	71
21 EKG-DAAS Overlay Structure and Memory Map . . . .	72
22 EKG-EXEC Functional Flowchart . . . . .	75
23 NOCPRS Functional Flowchart . . . . .	77

24	DISPLAY Functional Flowchart . . . . .	79
25	TOLAN Functional Flowchart . . . . .	81
26	Experimental Data Collection Setup . . . . .	86
27	Original and Reconstructed EKG With Best Compression Ratio . . . . .	89
28	Original and Reconstructed EKG With Average Compression Ratio . . . . .	90
29	Original and Reconstructed EKG With Worst Compression Ratio . . . . .	91
30	Tolan Compression Ratio Breakdown . . . . .	95
31	Second Difference Distributions . . . . .	97
A1	The Heart Cross-Section . . . . .	109
A2	Wilson EKG Electrode System . . . . .	111
A3	Typical EKG Waveform . . . . .	113
B1	The Communication System Model . . . . .	117
B2	A Discrete Memoryless Channel . . . . .	121
B3	A Typical Rate Distortion Function . . . . .	124

Abstract

EKG data compression techniques were investigated for potential real time implementation on an 8 bit Motorola 6800 microprocessor. Research indicated entropy reduction transform techniques such as the Fast Fourier Transform and the discrete Karhunen-Loeve Transform were not feasible for implementation on the 6800. Two redundancy reduction (RR) techniques (TOLAN and DOWER) utilizing 2nd order difference operations in conjunction with variable length encoding were studied in detail. One such RR technique (TOLAN) was fully implemented and tested with "in vivo" EKG data. Analysis revealed compression ratios ranging from 1.25:1 to 2.26:1. Investigation of the poor performance of the compression algorithm showed significant degradation of the 2nd order difference "decorrelator" due to a noisy collection environment. It was concluded that real time EKG data compression is feasible on the 6800 but that time compression techniques which store a zero value sequence counter versus the value of zero are not efficient in a high noise environment.

ANALYSIS AND PERFORMANCE EVALUATION  
OF ELECTROCARDIOGRAM DATA COMPRESSION TECHNIQUES

I. Introduction

Background

Currently the USAF School of Aerospace Medicine (USAFSAM) is receiving more than thirty thousand electrocardiogram (EKG) data records a year from Air Force flight personnel worldwide. This data is presently recorded on paper strip charts for immediate medical analysis and long term storage (via microfilm).

Advances in computerized biomedical analysis has generated a need for digitized storage of the EKG waveform to allow computerized interpretation and comparison of present and past cardiographic data. As computerized diagnosis becomes more accurate, a long term digital record of historical EKG data will enable the cardiology staff at USAFSAM to identify developing heart disease before it becomes a danger to a flight crew member or his fellow crewman.

Data compression of the sampled EKG has been an area of active research since the late nineteen sixties. References (1), (7), (12), (26), (28), (29), (32), (33), and (35) are representative of the research efforts performed in the last ten years. The reasons for compressing EKG data are twofold: 1) digital storage costs are rapidly approaching

analog storage costs and ; 2) increased use of computer aided diagnosis requires large digital data bases.

In the past, EKG data compression has generally been performed at a central computer facility. The data is normally collected in a physician's office or clinic and transmitted in an analog format over a standard telephone link to the central computer. The data is then digitized and input to a computer for diagnosis and storage (in compressed form).

With the current advances in microprocessor technology, sampling and compression at the collection site is now a viable alternative. State of the art digital communication systems can operate at 9600 bits per second over the standard 3 kilohertz (KHz) bandwidth telephone channel. With error detection and correction protocols (Ref 27), high fidelity digital transmission of the compressed EKG to the central processing center appears to be the wave of the future (Ref 21:253-254).

The Department of Defense (DOD) is currently installing a computerized EKG interpretation system at centralized US military medical centers worldwide. To collect and transmit this EKG data, remote medical clinics will use a commercial EKG "cart" containing a Motorola 6800 microprocessor. This EKG cart performs internal data compression, record formatting, and error protective "channel" encoding for digital transmission to the central medical center.

As an independent study, USAFSAM has sponsored this

thesis to investigate the general field of microprocessor based EKG data compression. The results of this study are to be used for comparison against the data compression achieved in the DOD system and to create a measurement baseline to evaluate future microcomputer based EKG data compression systems.

#### Problem

The problem addressed by this thesis was the analysis of currently available EKG data compression algorithms for implementation in a microprocessor based computational environment. Additionally, a performance measure was to be developed by which differing EKG data compression techniques could be compared.

#### Scope

In this thesis, an EKG Data Acquisition and Analysis System (EKG-DAAS) was constructed. The EKG-DAAS collects 3 leads of an amplified EKG (i.e., Analog/Digital (A/D) dynamic range = $\pm$  5 volts), and samples the data at a operator controlled rate between 300-700 hertz. Data is digitized at 12 bit precision and subsequently rounded to 8 bits for uniform truncation error performance (Ref 24:413-418).

The EKG-DAAS hardware was developed around a Motorola Exorciser microcomputer (6800 microprocessor) with associated A/D converter, disk memory, computer terminal, and

hard copy printer. The EKG-DAAS hardware is controlled via a software program called EKG-EXEC. EKG-EXEC is an assembly language program which performs terminal, printer, and disk input/output (I/O) operations as well as providing a supporting structure for the data compression and analysis software. Because of the lack of a high order language (e.g. FORTRAN, PASCAL, etc.) and the desire for maximum program execution speed, all programming done in this thesis is in assembly language. This has proven to be a major limitation.

Only one EKG data compression algorithm was completed and implemented on the EKG-DAAS. "In vivo" EKG data was taken, however, and EKG data compression performance analyzed.

#### General Approach

To accomplish the objectives in the problem statement, the literature was first searched for EKG compression algorithms whose implementation and execution on a 6800 microprocessor was considered feasible. As a result of this literature search and private correspondence (Ref 11 and 31), two compression routines were found. These two EKG data compression algorithms will hereafter be referred to as the Tolan and Dower methods (Ref 31,12).

With the Tolan and Dower algorithms identified, construction of the EKG-DAAS was begun by assembling the data acquisition subsystem. This data acquisition system

ccnists of an Analog/Digital converter and a interrupt sampling clock. The A/D was calibrated and coding of the EKG-EXEC software started. Several months of effort resulted in the EKG-EXEC software and programming of the Tolan compression algorithm initiated. Upon completion of the Tolan compression routines, the Dower algorithm was analyzed but not implemented.

Based on the results of the Tolan compression algorithm, a performance measure was formulated which compares achieved compression against an approximate maximum compression computed from the data statistics. The Dower compression algorithm was analyzed for similarity with the Tolan technique and an estimated performance figure calculated with respect to the compression measure or "metric".

#### Sequence of Presentation

Chapter 2 begins the thesis development with a general survey of the field of data compression. Data compression is shown to be divided into two subclasses (Entropy Reduction and Redundancy Reduction) and each subclass is defined. Several EKG compression techniques found in the literature were included in the ER category. These ER techniques were described and their performance advantages and limitations analyzed. Chapter 2 proceeds by describing redundancy reduction and several RR electrocardiogram compression algorithms are also analyzed. Chapter 2

concludes with the determination that implementation of the identified ER compression routines would require programming efforts beyond the scope of this thesis. Hence this thesis is limited to redundancy reduction compression.

Chapter 3 describes, in detail, the Tolan and Dower redundancy reduction algorithms whose implementation on the Exorciser was considered feasible. The Tolan algorithm is described first followed by a discussion of the Dower method. Two other EKG compression techniques, discovered during the research of this thesis, are also summarized.

Chapter 4 reviews the hardware and software configuration of the EKG-DAAS. The hardware system is described first and the specifications of the Exorciser microcomputer, A/D converter, and I/O peripherals are presented. Following the hardware description, the EKG-EXEC program is documented and the software design philosophy examined. Finally the Tolan compression module is discussed.

Chapter 5 outlines the results obtained using the EKG-DAAS with the Tolan algorithm. A detailed description of the experimental setup is presented along with a discussion of the performance parameters measured by EKG-EXEC. Finally a performance metric is described and the Tolan compression results compared against this metric. Chapter five ends with an analysis of how well the Dower compression technique would have performed against the performance metric. Chapter 6 concludes the thesis and

presents recommendations for future study.

This thesis contains five appendices. Appendix A surveys basic electrocardiology and is recommended to readers unfamiliar with this subject. Appendix B is a tutorial on Information and Coding Theory and is likewise recommended to those readers unfamiliar with this field. Appendix C contains a listing of the EKG-EXEC assembly language software as well as one BASIC program used on the Exorciser for data analysis. Appendix D contains data printouts of the data taken in the data collection experiment. Appendix E presents photocopies of the specification sheets for the equipment used in this thesis.

With the sequence of presentation outlined, attention now turns to the theoretical section of this thesis. The first subject is data compression theory.

## II. Data Compression and the EKG

Data compression is an operation in which data from an information source is "simplified" or "filtered" in a manner that produces an approximation of the original with at most some predefined amount of distortion. Some form of data compression is usually necessary when storage limitations, bandwidth requirements, or transmission channel capacity prohibit operation on the original data.

In general, data compression can be divided into two types of operations (Ref 10:4). The first operation, called entropy reduction (ER), is an irreversible transformation which reduces or compresses the data by mapping a source into an approximation of itself with a lower entropy rate. Sampling is an example of such a transformation. The second operation, known as redundancy reduction (RR) compresses the sampled data train by reducing, or eliminating, the redundancy existing in digital sequence. Since the redundant components of the data train contain no information about the source, RR is an "exact" data compression operation.

This chapter discusses data compression as it has been applied to the electrocardiogram (EKG). Two types of data compression strategies prevail. The first strategy involves two (or more) ER operations on the EKG data and attempts to compress the data by filtering selected components of discrete transforms. Because ER operations are irreversible, this type of data compression is sometimes

referred to as "inexact" compression. The second strategy performs data compression by reducing the redundancy present in the sampled EKG data train. This redundancy arises from two causes : 1) neighboring signal samples are not statistically independent and ; 2) the quantized signals amplitudes do not occur with equal probability.

Because of the speed limitations of the Motorola Exorciser microcomputer used in this thesis, only redundancy reduction algorithms were tested. This chapter, therefore, is intended as a review of the work that has been done in EKG data compression to allow comparison with the results obtained by this author.

#### Entropy Reduction

As is known from information theory (Ref 19), entropy  $\{H(X)\}$  is defined as a measure of the "randomness" or "uncertainty" of an information source. If the symbols emitted by the source are statistically independent, then the source is said to be "memoryless" and entropy is given by the equation:

$$H(X) = - \sum_{i=1}^N p_i \log_2 p_i \quad (\text{bits}) \quad (1)$$

where  $p_i$  represents the probability of occurrence of the  $i$ th symbol and  $N$  is the number of distinct symbols output by the source. For those "symbols" in a data set which occur

with zero probability, the term  $p_i \log_2 p_i$  is defined equal to zero.

Immediately obvious from Eq.(1) is the observation that  $H(X)$  is defined on a discrete probability distribution. In fact, Shannon (Ref 30) defined the entropy of a continuous source as equal to positive infinity. Thus an ER transformation must, in some manner, "discretize" a continuous waveform into a countable set of components and attach some probability to the elements of the set.

Sampling. The clearest example of this "discretization" is sampling, obviously the most important operation necessary for digital signal processing. In sampling, an electrical circuit (such as an Analog-to-Digital (A/D) converter) periodically measures the value of a signal  $x(t)$  and records the data as a numeric (usually binary) number. If the signal is sampled at least twice the highest frequency component of  $x(t)$ , and the duration of the sampling operation is long enough that "aliasing" effects (Ref 36:68-72) of "windowing" are negligible, then all of the "frequency domain" components of interest in  $x(t)$  will be preserved. Physical constraints, however, limit the accuracy of the amplitude measurement to some finite precision. Amplitude information residing below the sensitivity of the A/D converter is irretrievably lost.

As an example, let the A/D converter digitize to 8 bits. With 8 levels there are 256 possible outputs, each with a certain probability of occurrence. If the continuous

waveform fed to the A/D is the output of a stationary stochastic process with uniform statistics, then each numeric value will occur with probability of 1/256 (assuming the source max/min deviation  $\geq$  the A/D dynamic range). Assuming sample-to-sample independence, the source entropy is then calculated as

$$H(X) = \sum_{256} (1/256) \log_2 256 = 8. \quad (2)$$

Entropy has been "reduced" from  $+\infty$  to 8. If the stochastic source has "less random" statistics (like gaussian), entropy would be reduced even further.

All of the EKG compression routines discussed in this thesis are implemented on digital computers, hence the ER operation of sampling is always performed. In "inexact" EKG data compression, the next operation is another ER mapping in which the sample sequence is transformed into an alternate domain and filtered.

Orthogonal Expansions and Filtering. A bandlimited waveform  $x(t)$  can be expanded (over a given interval  $T$ ) as a linear combination of orthonormal basis functions  $\phi_n(t)$  (Ref 36:20-21). That is

$$x(t) = \sum_{n=0}^{\infty} a_n \phi_n \quad (3)$$

The functions  $\phi_n(t)$ , as  $n=0, 1, \dots, \infty$ , are defined as orthonormal if

$$\int_T \phi_i(t) \phi_k(t) dt = 1, \quad i=k \\ = 0 \quad \text{otherwise} \quad (4)$$

The orthonormal set  $\{\phi_n(t)\}$  is called complete (C.O.N.) if, for any given  $\epsilon > 0$ , there exists an  $N$  and a finite expansion

$$\tilde{x}(t) = \sum_{n=0}^{N-1} a_n \phi_n(t) \quad (5)$$

such that

$$\int_T |x(t) - \tilde{x}(t)|^2 dt < \epsilon \quad (6)$$

Given a complete, orthonormal set  $\{\phi_n(t)\}$ , then representation of any physical, noise-limited signal (over a prescribed interval  $T$ ) is possible with a finite set of weighting coefficients  $\{a_0, a_1, \dots, a_{N-1}\}$ .

If  $k$  ( $k < N$ ) terms in Eq.(5) are selectively suppressed (i.e.,  $a_0, a_1, \dots, a_k = 0$ ), then the signal  $x(t)$  is said to be filtered of those components and a compressed representation of  $x(t)$  is produced.

EKG Filtering Compression. A great deal of work (1),(2),(26),(33),(35),(21) has been done in attempting to "compress" EKG data by storing, or transmitting, selective components of an orthonormal expansion. This type of compression strategy is complicated by the complex characteristics of the EKG waveform. Foremost are the limitations imposed due to the "nonstationarity" or variability of the EKG source, both within any given waveform and the general population as a whole.

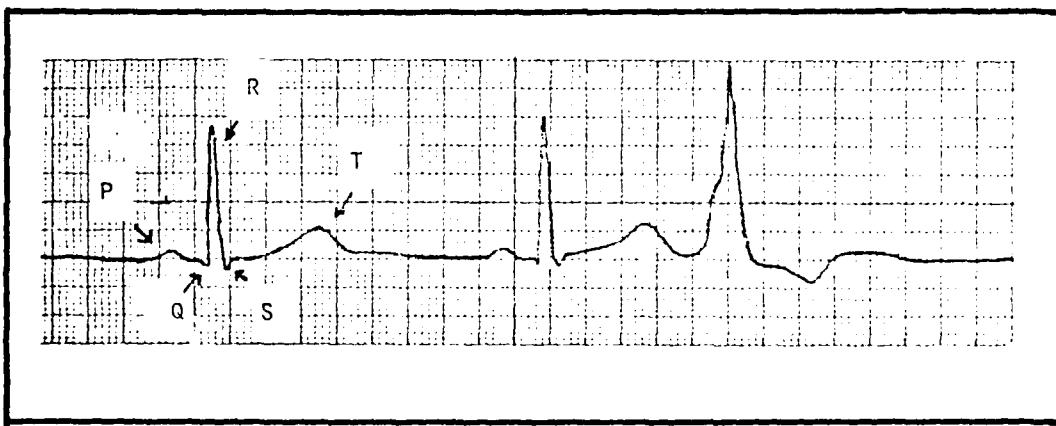


Fig 1. Typical EKG waveform with arrhythmia.

During an EKG collection, the heart rate of the person undergoing testing can vary. This variation, or arrhythmia, requires that a suitable reference point be established on

the P-QRS-T complex (Fig 1) around which the orthogonal expansion can be performed. Without this reference, the data falling within an expansion interval T (based purely on a constant  $\Delta t$ ) will appear to be "almost" random. If the "filtering" used to effect compression is based on an expected waveform within the transform window, improper registration will cause severe degradation in compression efficiency. The above situation is known as the "epoch" problem and is discussed by Womble (Ref 35:703).

Assuming that the expansion interval is aligned properly on some feature of the EKG (e.g. R wave), then the problem of variable P, QRS, and T complexes must be considered.

If the individual under test has a heart disease, then in many cases the P, QRS, and T segments of that person's EKG may vary significantly from the population "norm". Again, if the compression filtering is based on an assumed, or "normal" waveform, then significant reduction in compression efficiency can be expected.

These problems, among others, require that the EKG "filtering" compressor be robust enough to handle the variations possible within the waveform. With these limitations in mind, discussion will now proceed to compression transforms.

Transforms. Two major types of discrete filtering compression strategies, as applied to EKGs, will now be discussed. These are the Fast Fourier Transform (FFT) and

the discrete Karhunen-Loeve Transform (DKLT). These two techniques are representative of current research into EMG filtering compression and will serve as good basepoints against which to compare the redundancy reduction compression techniques outlined in chapter 3.

Fast Fourier Transform. The fast Fourier transform is a computationally efficient algorithm for calculating the discrete Fourier transform (DFT). The following discussion involves the DFT, but calculation by the FFT is implied.

Data compression using the DFT is achieved by zonal filtering in which "zones" of the transform sequence are selectively discarded. Usually these zones are defined by a "cutoff frequency"  $f_c$  and only those components less than or equal to  $f_c$  are saved. Data reconstruction is performed by computing the inverse DFT (FFT) on the filtered sequence and replacing the filtered components with zeroes. Using the symmetry relationships described in Oppenheim and Shaffer (Ref 24:103-105), then the transform sequence for N=8 is:

$$X(k) = \begin{bmatrix} X(0) \\ X(1) \\ X(2) \\ X(3) \\ X(4) \\ X(5) \\ X(6) \\ X(7) \end{bmatrix} = \begin{bmatrix} R(0) + j\text{Im}(0) \\ R(1) + j\text{Im}(1) \\ R(2) + j\text{Im}(2) \\ R(3) + j\text{Im}(3) \\ R(4) + j\text{Im}(4) \\ R(3) - j\text{Im}(3) \\ R(2) - j\text{Im}(2) \\ R(1) - j\text{Im}(1) \end{bmatrix} \quad (7)$$

where R is the real part, Im the imaginary part of the complex number X, and  $j=\sqrt{-1}$ .

Zonal filtering takes advantage of the symmetry in Eq.

(7) by saving those positive frequency components less than or equal to a certain cutoff frequency  $f_c$ . If  $f_c$  were chosen as equal to  $X(2)$ , then only  $X(0)$ ,  $X(1)$ , and  $X(2)$  would be saved. On reconstruction,  $X(3)$ ,  $X(4)$ , and  $X(5)$  would be set equal to zero (i.e. they have been filtered out) while  $X(6)$  and  $X(7)$  would be recreated from  $X(1)$  and  $X(2)$  using symmetry. The inverse DFT (FFT) now produces a "filtered" approximation of  $x(n)$ .

EKG data compression via filtered FFT spectra has been studied by Womble (Ref 35) and the TRW corporation (Ref 33). In both studies the distortion criterion used in establishing  $f_c$  was visible reproducibility with no detectable distortion.

In the TRW study, compression ratios (defined as bits in : bits out) from 2:1 to as high as 17:1 were obtained using zonal filtering. The high compression ratios were measured using highly rhythmic EKGs taken from an individual (an astronaut) in a low noise environment. This extraordinary compression ratio is far from "normal", however. Womble (Ref 35) has shown that, on the average, 40-80 terms of a 512 sample FFT (sampled at 500 Hz) are necessary to reproduce the EKG with acceptable visual distortion. Since the FFT requires storage of 2 numbers per term (real and imaginary components) then, in general, FFT zonal filtering compression ratios of 5:1 to 3:1 are more common.

The FFT is often used in digital signal processing

because of its speed and representation in the frequency domain. As will be seen from the data in chapter 5, compression ratios of 5 : 1 are usually higher than those obtainable by redundancy reduction techniques. Even so, the FFT is not the "optimal" transform for representing a sample sequence.

Discrete Karhunen-Loeve Transform. Because all of the filtering compression techniques are ER transformations, distortion upon reconstruction is inevitable. One would like a performance measure against which this distortion could be evaluated.

One common measure of performance is the mean square error defined by the equation:

$$e(M) = E\{(\underline{x} - \tilde{\underline{x}})^2\} \quad (8)$$

where  $E$  is the statistical expectation operator,  $\underline{x}$  is a discrete signal composed of  $N$  sample values, and  $\tilde{\underline{x}}$  is the estimate of  $\underline{x}$  via some orthogonal coordinate system. Given that a representation of  $\underline{x}$  is desired with less than  $N$  components, Ahmed and Rao (Ref 2:200-203) have shown that the DKLT is the optimum transform to minimize mean square error.

The DKLT can be described as follows: Let the orthonormal transform matrix  $[T]$  be defined as

$$[T]' = [ \emptyset_1, \emptyset_2, \dots, \emptyset_N ] \quad ([T]' = [T] \text{ transpose}) \quad (9)$$

and  $\phi_i$  are N dimensional, real valued basis vectors. Hence a uniform vector  $\underline{Y}$  can be created by

$$\underline{Y} = [T] \underline{X} \quad (10)$$

where  $\underline{X} = [x_1, x_2, \dots, x_N]$ ;  $\underline{Y} = [y_1, y_2, \dots, y_N]$ . Since the  $\phi_i$  are orthonormal, then

$$\underline{X} = y_1 \phi_1 + y_2 \phi_2 + \dots + y_N \phi_N = \sum_{i=1}^N y_i \phi_i \quad (11)$$

The goal is to optimally represent  $\underline{X}$  by a subset  $\{y_1, y_2, \dots, y_M\}$  where  $M < N$ . If the remaining  $N-M$  terms are represented by the constants  $b_i$ , then an estimate  $\tilde{\underline{X}}$  is defined such that

$$\tilde{\underline{X}} = \sum_{i=1}^M y_i \phi_i + \sum_{i=M+1}^N b_i \phi_i \quad (12)$$

An error vector,  $\Delta \underline{X}$  is now created where

$$\Delta \underline{X} = (\underline{X} - \tilde{\underline{X}}) = \sum_{i=M+1}^N (y_i - b_i) \phi_i \quad (13)$$

Now the mean square error defined in Eq.(8) can be redefined

as

$$\begin{aligned} e(M) &= E\{(\Delta \underline{x})' (\Delta \underline{x})\} \\ &= E\left(\sum_{i=M+1}^N \sum_{j=M+1}^N (y_i - b_i)(y_j - b_j)(\phi'_i \phi_j)\right) \\ &= \sum_{i=M+1}^N E\{(y_i - b_i)^2\} \end{aligned} \quad (14)$$

Ahmed and Rao (Ref 2:202) show that by minimizing Eq.(14)

$$e(M) = \sum_{i=M+1}^N [\phi'_i E\{(\underline{x} - \bar{\underline{x}})(\underline{x} - \bar{\underline{x}})'\} \phi_i] \quad (15)$$

where  $\bar{\underline{x}}$  is the mean of  $\underline{x}$ .

The expectation in Eq. (15) is recognized as the covariance of  $\underline{x}$ , hence Eq. (15) becomes

$$e(M) = \sum_{i=M+1}^N (\phi'_i [K_x] \phi_i) \quad (16)$$

where  $[K_x]$  is the covariance matrix of  $\underline{x}$ .

By minimizing Eq.(16), Ahmed and Rao (Ref 2:200-205)

show that the orthonormal functions  $\phi_i$  become the eigenfunctions of the covariance matrix  $[K_x]$  and the mean square error becomes

$$e(M) = \sum_{i=M+1}^N \lambda_i \quad (17)$$

where the  $\lambda_i$  are the eigenvalues of  $[K_x]$ . By expanding  $\underline{X}$  with the eigenfunctions corresponding to the  $M$  largest eigenvalues, then  $\underline{X}$  is filtered of  $N-M$  components with minimum mean square error. This is the discrete Karhunen-Loeve transform.

The DKLT has been studied by Ahmed (Ref 1) and Womble (Ref 35). Both studies compared the performance of EKG data compression using the DKLT against other orthogonal transforms.

Ahmed's tests utilized canine EKGs and clearly showed the optimality (in the "mean square sense") of the DKLT. Ahmed, however, deemed the DKLT to be too complex for practical implementation, and proceeded with the development of EKG compression strategies using suboptimal transforms. The details of these suboptimal expansions will not be given here, but the reader is referenced to Ahmed's paper (Ref 1).

Womble, on the otherhand, demonstrated that the DKLT is not difficult to implement. In his experiments, Womble

took 3 lead vectorcardiogram (VCG) data (see appendix A), hereafter referred to as X, Y, Z, and first transformed the data into a different, orthogonal coordinate system (U, V, W). The transformation was determined by solving for the eigenvalues of the 3 X 3 matrix:

$$S = 1/N \sum_{i=1}^N \begin{bmatrix} X(i) \\ Y(i) \\ Z(i) \end{bmatrix} [X(i), Y(i), Z(i)] \quad (18)$$

where N was chosen equal to 200 and X(i), Y(i), Z(i) are the X, Y, Z components of the VCG in the Frank (Ref 14) coordinate system.

Next an "average", or mean, heartbeat was calculated using 900 patients and the data (in the eigenvector coordinate system) was subtracted from the mean forming a "patient" vector  $\rho$ . The vector  $\rho$  is defined as

$$\rho = \begin{bmatrix} u(1) - u_m(1) \\ u(2) - u_m(2) \\ \vdots \\ u(N) - u_m(N) \\ v(1) - v_m(1) \\ v(2) - v_m(2) \\ \vdots \\ v(N) - v_m(N) \\ z(1) - z_m(1) \\ z(2) - z_m(2) \\ \vdots \\ z(N) - z_m(N) \end{bmatrix} \quad (19)$$

This patient vector  $\rho$  is expanded using the eigenfunctions

of the matrix

$$L = 1/N \sum_{i=1}^N (\rho_i)(\rho_i)' \quad (20)$$

where  $N$  is now a large number of patients being "averaged" (up to 300 reported in Ref 33). The eigenvectors of this tremendous matrix are calculated (once for all time), and then  $M$  ( $M \ll 200$ ) are used to expand the data vector  $\rho$ .

The above approach has allowed accurate representation of EKG sample sequences, properly aligned on the heartbeat, with an  $M=20$  DKLT coefficients per lead per second (Ref 35). At Womble's sample rate of 250 Hz, this represents a compression ratio of over 12 : 1 . With the most significant eigenfunctions stored from the solution of Eq. (20), then the compression of digitized EKG data by filtering the DKLT expansion on  $\rho$  is approaching feasibility on a microcomputer.

The preceding discussion was intended as a basic tutorial on some of the entropy reducing (ER) transformations currently under test for EKG data compression. As stated earlier, ER operations are "inexact" because some "information" is always discarded. By discarding data efficiently, zonal filtering can achieve compression ratios larger than those obtained using the redundancy reduction techniques implemented in this thesis.

### Redundancy Reduction

Redundancy reduction is an operation which performs data compression by indentifying and removing the redundant components of a digital sequence. This redundancy exists for two reasons : 1) the sample points are not statistically independent and ; 2) the quantized amplitude values of the sample train do not occur with equal probability. The redundant components of the data sequence carry no "information" (see appendix B), hence their removal does not affect the "message" content of the data.

This section of chapter two discusses redundancy reduction (RR) operations which have been applied to the EKG. This is done because the compression algorithms examined in this thesis are of the RR type, and the background given here will aid in the descriptions given in chapter three.

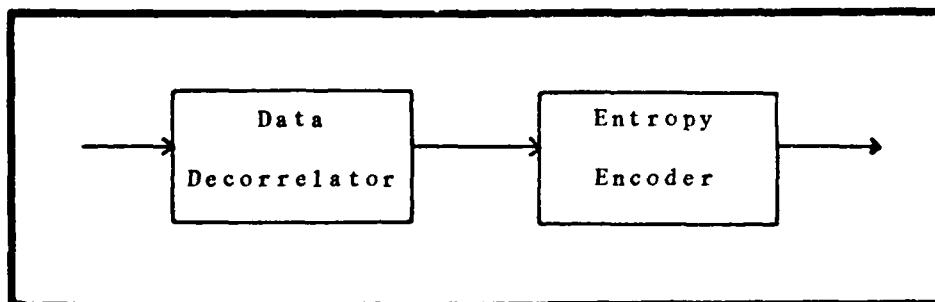


Fig 2. Typical EKG Redundancy Reduction Data Compressor.

The sampled EKG contains redundancy resulting from both

sample-to-sample correlation and unequal amplitude probability. As is illustrated by Figure 2, each of these redundancy components can be reduced by a different operator.

The decorrelator, as inferred by its title, transforms the data in a manner which decorrelates the data stream. The decorrelated residual sequence is then efficiently compressed by means of the entropy encoder.

To reduce the redundancy arising from intersample correlation, two basic approaches are used. The first approach utilizes linear predictor/interpolators (Ref 28,33) to produce an information bearing "residual" sequence. The second method generates the residual sequence by taking successive differences on the data stream. As will be shown, these residual sequences carry all of the source information with minimum intersample redundancy. The predictor/interpolator method is discussed first.

Predictors and Interpolators. A predictor estimates the next sample value of a sequence (i.e.  $x_n$ ) based on a linear combination of  $k$  past samples. That is

$$\tilde{x}_n = \sum_{i=1}^k a_i x_{n-i} \quad (21)$$

where  $a_i$  are coefficients chosen to minimize the mean square error  $\sigma_e^2 = E\{(X - \tilde{X})^2\}$  between the sample sequence  $X_n$  and the

predicted sequence  $\tilde{x}_n$ . An error, or residual, sequence  $e_n$  can then be formed where

$$e_n = x_n - \tilde{x}_n \quad (22)$$

thus

$$x_n = \tilde{x}_n + e_n \quad (23)$$

Conceptually, this technique decomposes a sample value into a part which is correlated with  $k$  past sample values (i.e. the redundant component) and a part which is uncorrelated with them. Linear mean square estimation theory (Ref 24:385-430) shows that the uncorrelated part ( $e_n$ ) may be retained alone with no loss of information.

In their research in EKG data compression, Ruttimann and Pipberger (Ref 28:616) prove that since  $E[e_n] = 0$ , then the mean square error  $\sigma_e^2$  is equal to the variance of  $e_n$ . With a second order predictor ( $k=2$ ), Ruttimann and Pipberger have demonstrated a variance reduction ( $\sigma_x^2/\sigma_e^2$ ) of over 25 : 1. This reduction in variance implies that the EKG residual sequence  $e_n$  is much more tightly clustered around a given mean than is the original sequence  $x_n$ . This clustering, as will be discussed later, enhances data compression by means of entropy, or source encoding.

In an analogous way, interpolators estimate a value of  $x_n$ . In the case of the interpolator, however, the estimate

of  $x_n$  consists of a linear combination of past and future samples. That is

$$\tilde{x}_n = \sum_{i=1}^k a_i x_{n-i} + \sum_{i=1}^m b_i x_{n+i} \quad (24)$$

where  $k$  past and  $m$  future values are used. As with the predictor, the coefficients  $a_i$  and  $b_i$  are again chosen to minimize the mean square error. As was the case with the second order predictor, Ruttiman and Pipberger (Ref 28:617) have found that a second order interpolator ( $k=1, m=1$ ) yields EKG residual sequences  $e_n$  with the most significant variance reductions (greater than 31 : 1).

EKG data compression utilizing predictors/interpolators have been studied by the TRW Corporation (Ref 33), as well as Ruttimann and Pipberger (Ref 28). In both cases, predictors/interpolators of order 2 seem to prevail. As is shown by the two groups above, second order systems have the smallest residual sequence variance hence are most amenable to entropy encoding. An alternate approach used to reduce intersample redundancy is by means of difference operations.

Difference Reduction. In difference reduction, the residual sequence is formed by taking successive differences of the sample data train. Because there is no multiplication (as in Eq.(24)), the difference operation is inherently a

simpler procedure than is prediction or interpolation. Difference sequences, however, are not optimized with respect to mean square error and less "efficient" redundancy reduction (decorrelation) occurs.

An example of an EKG difference reduction compressor is illustrated in Figure 3. This technique was implemented by Cox and Ripley (Ref 7) and tested against a data base of 45 patients. Figure 4 shows that Cox and Ripley's second difference decorrelator produced a sharply peaked relative frequency distribution. This "peakedness" or clustering of the residual sequence permits efficient entropy encoding. A variant of the difference reduction above is time compression.

Time Compression. EKG time compression, as implemented by Dower, Berghofer, and Stewart (Ref 12,29), uses a second order difference reduction but instead of keeping the value of zero (the most common), a run length counter ( $\Delta t$ ) is kept. This run length counter measures the number of repetitive  $\Delta^2 X$  sequence terms equal to zero. By saving (then encoding) only those  $\Delta^2 X \neq 0$ , and the  $\Delta t$  between the nonzero second differences, data compression can be realized.

Time compression is the method that has been implemented and studied extensively by this author. Chapter 3 discusses two different RR compression algorithms which both use second order time compression for the decorrelator. For now, discussion is focused on the second

operation identified in Figure 2, the entropy encoder.

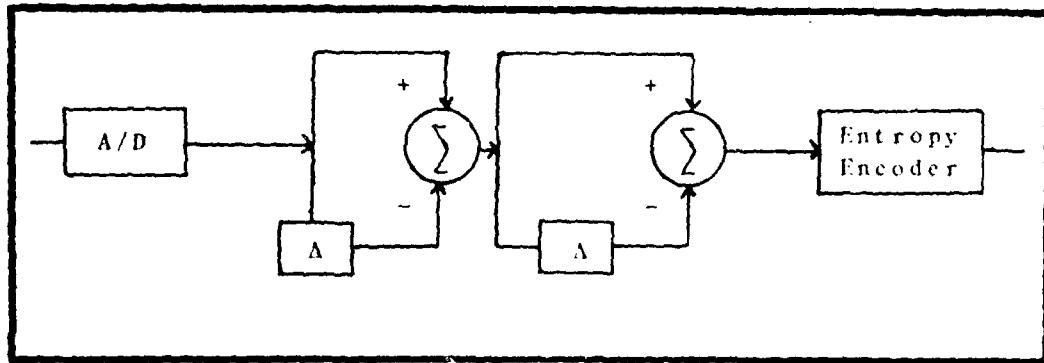


Fig 3. Second Order Difference Operator (from Ref 7:336).

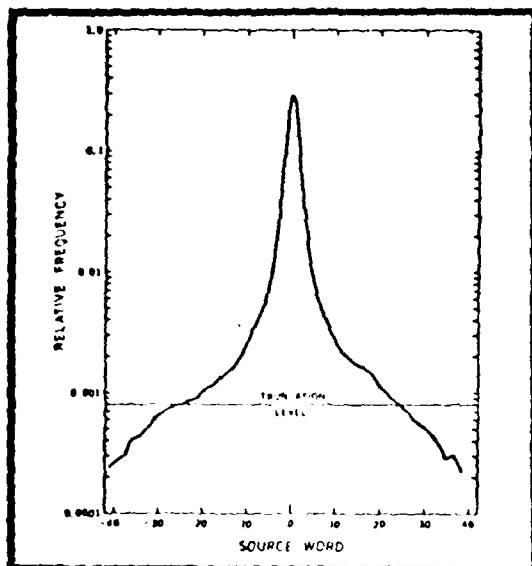


Fig 4. Relative Frequency of 2nd Order Difference Operator  
(Ref 7:336)

Entropy Encoding. Entropy, or source, encoding operates on the remaining redundancy in the residual sequence  $e_n$  resulting from unequal source symbol probability. How well  $e_n$  can be compressed is given by Shannon's noisless source coding theorem (Ref 30) which states that the average number of binary symbols per source output can be made to approach the entropy of the source and no less. Unfortunately, the EKG "source" has memory, hence the entropy of this source will be less than that which would be calculated from Eq.(1). In general, the true value of a nonstationary, memory source is difficult (or impossible) to calculate.

To circumvent this "memory" problem, the assumption is made in the literature (Ref 29,12, 7,35) that the "uncorrelated" nature of the residual sequence  $e_n$  approaches independence, hence the entropy of  $e_n$  as calculated by Eq.(1) is a good bound on the possible compression. Since the variance of  $e_n$  is sharply peaked around zero, entropy encoding via variable length coding (VLC) appears attractive.

In variable length coding, those values of  $e_n$  which occur most often are assigned the shortest code words (i.e. fewest code symbols per source symbol). As example, assume that both the source symbol alphabet and the code alphabet are binary. The values in the sequence  $e_n$  are the result of algebraic operations on fixed length numeric sample values and are hence fixed length binary numbers (e.g.  $L=8$ ). The variable length coder maps these fixed length numeric values

(binary) into uniquely decodable (UD) binary code words whose bit length is a function of the frequency of occurrence of the values in  $e_n$ .

Coding theory (Ref 19: 237-248) shows that, in principle, it is possible to have variable length codes with an average bit length ( $\bar{l}$ ) equal to the entropy of source. Hence with the right VLC, a compression ratio of  $L/\bar{l}$  is achievable.

A common VLC used in EKG data compression research (Ref 7, 28) is the Huffman code. Huffman, in 1952, developed an algorithm (Ref 16) for generating the optimal UD code (assuming stationary, memoryless source). This code is in a class of UD codes called prefix codes in which no code word is the prefix of any other. Unfortunately, a codeword must be assigned to every possible symbol which occurs, regardless of how infrequently. This means that although the average code word bit length approaches the entropy of the source, the longest code word can be substantially larger. In a straight Huffman code, these "long" code words can induce severe problems due to "buffer overflow" (Ref 17).

For practical implementation in EKG data compression, Ruttimann and Pipberger (Ref 28) and Cox and Ripley (Ref 7) constructed a modified Huffman code. In this modified code, the residual sequence source words are partitioned into a frequent set and an infrequent set (known as "else"). A Huffman code was then formed with all of the residual words

in the first set plus a special code word used as a prefix for any source word from the infrequent set. The prefix, when it occurs, is followed by a fixed length suffix which contains the value of the infrequent source word. The probability that "else" will occur must be kept small enough to maintain the efficiency of the truncated Huffman code.

Using the modified Huffman code in conjunction with a second order interpolator, Ruttimann and Pipberger (Ref 28) have attained compression ratios as high as 9 : 1. To get this ratio, however, significant digital signal processing was performed on the "raw" 8 bit data. This preprocessing involved digital filtering for noise reduction and Lagrange interpolation/decimation to produce an effective 200 Hz sample rate. Nonetheless, this represents a significant approach to the compression ratios attainable by the discrete Karhunen-Loeve transform mentioned previously.

Although optimal, Huffman codes are not the only VLC used in EKG data compression. Two different codes will be described in chapter three, one of which (Dower code) may possibly be "more optimal" than Huffman in the compression of quantized EKGs.

### Summary

This chapter has reviewed the theory of data compression and how this theory has been applied to the electrocardiogram. It was shown that data compression could

be partitioned into two types of operations: 1) entropy reducing transformations which map a data source into an approximation of itself with a lower entropy rate and ; 2) redundancy reduction operations which compress by removing redundancy resulting from sample-to-sample dependence and unequal source symbol probabilities.

The first operation was applied to the EKG via "filtering" of orthogonal expansions of the digitized EKG. Filtering transforms, especially the discrete Karhunen-Loeve transform, were shown to be very efficient "compressors" if the loss of the filtered components were tolerable. If this loss was not acceptable, then redundancy reduction operations are used.

In redundancy reduction techniques, the EKG is first processed by passing the digitized data through a decorrelator which reduced the redundant component caused by source symbol dependence. This decorrelator could be implemented with predictor/interpolators or difference operations. Next the data is encoded via entropy encoding operations, usually utilizing variable length codes. It was finally shown that a second order interpolator, followed by an "optimal" Huffman encoder could achieve compression ratios which approach those obtained by the discrete Karhunen-Loeve transform. Both of these last two techniques, however, require substantial computational overhead.

### III. Tolan and Dower EKG Compression Techniques

Two EKG data compression algorithms are studied in detail in this chapter. The first compression technique, hereafter referred to as the Tolan method, was conceived by Dr. Gil Tolan, MC from the USAF School of Aerospace Medicine (USAFSAM), Brooks AFB, Texas. The second compression approach, referred to as the Dower method, was developed by Mr. Roger Dower and Mr. Dave Berghofer from Shaughnessy Hospital, Vancouver, B.C., Canada. Both of the EKG compression procedures are redundancy reduction operations using time compression for decorrelation and variable length codes for entropy encoding.

This author had originally intended to implement and test both the Tolan algorithm and the Dower algorithm on the Motorola microcomputer (See Chapter 4). Unfortunately, difficulties with hardware failures and insufficient software tools (i.e., a high order language) prevented implementation of the Dower algorithm. Nonetheless, this chapter compares the design of both the Tolan and the Dower algorithms and illustrates their differences as well as their similarities. Based on the results of the Tolan compression approach (chapter 5), this comparison will illustrate the potential performance of the Dower compression strategy in a real time microcomputer environment. Before continuing with the Dower and Tolan algorithm descriptions, a short digression will be made.

In both the Tolan and The Dower compressors, a second

difference operation is used to produce an information bearing sequence with reduced component-to-component dependence. Why the second difference operation results in the lowest correlation is a question for closer scrutiny.

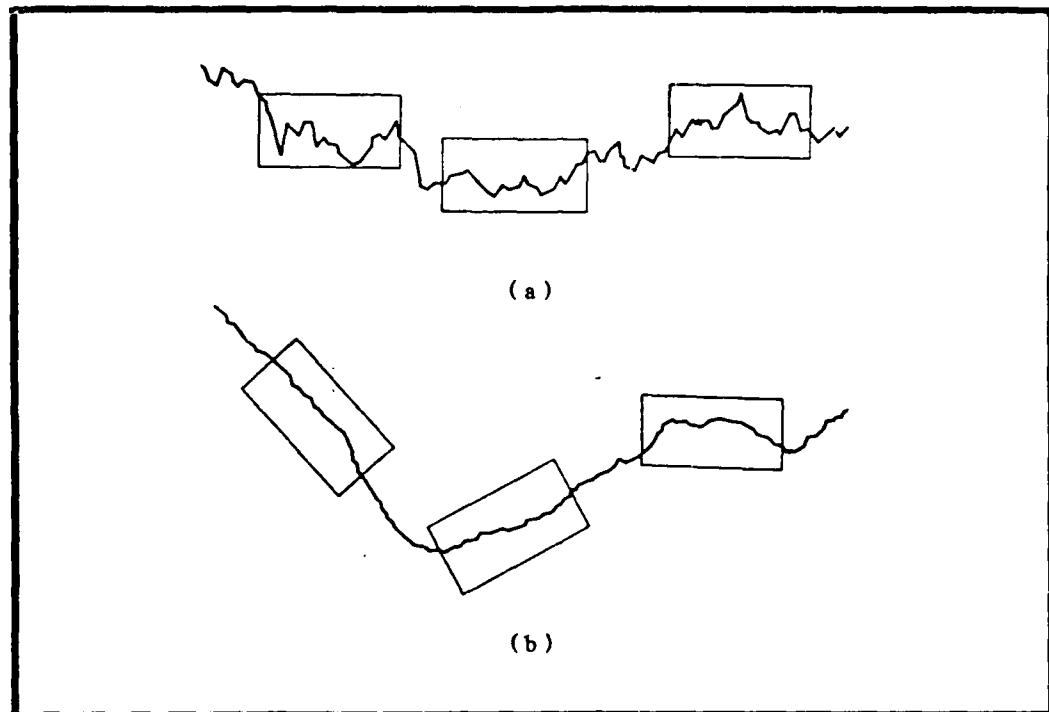


Fig 5. Nonstationary waveforms with random level and slope.  
(From ref 6:91)

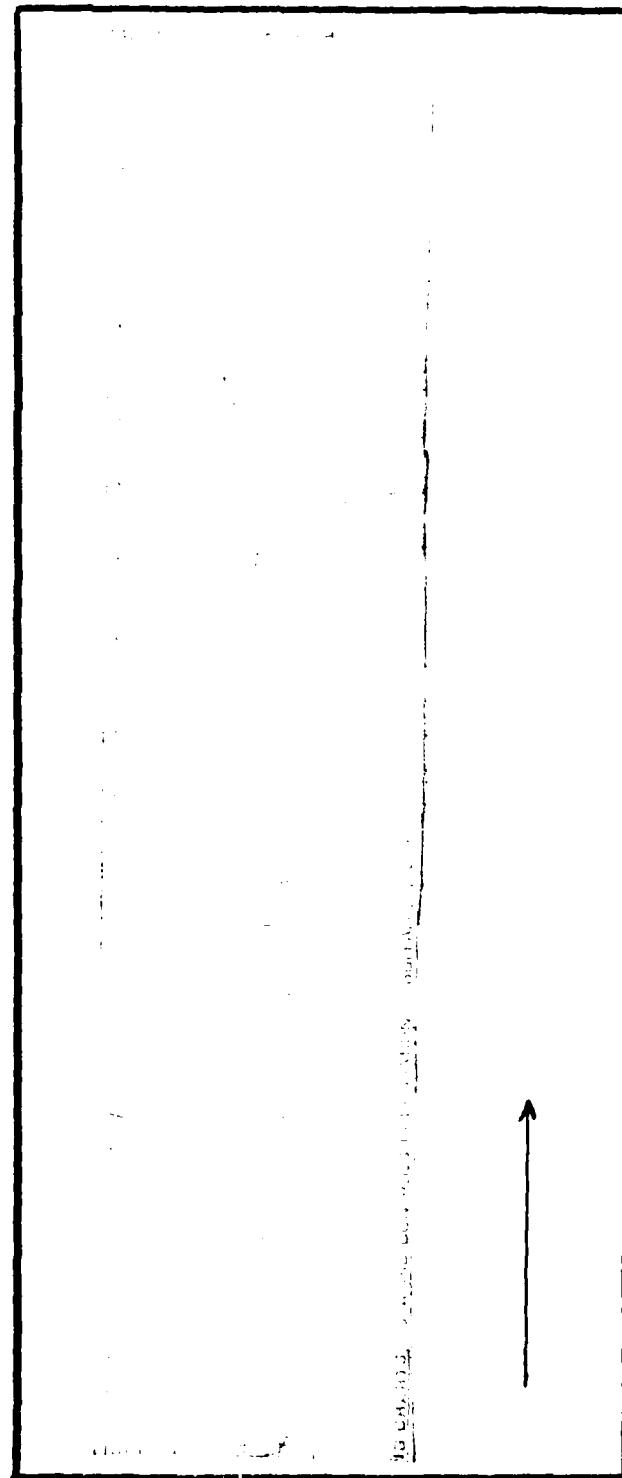
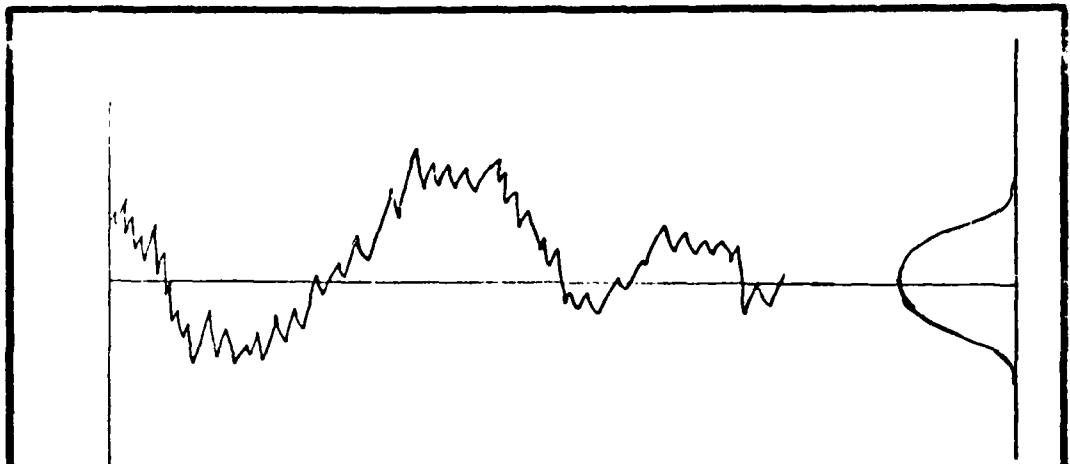


Fig 6. EKG waveform with moving average slope.

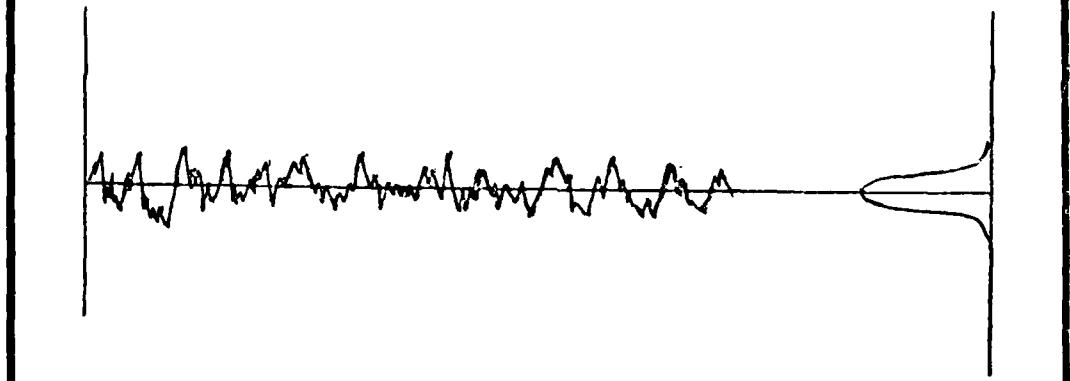
A nonstationary process (like the EKG) exhibits a wide amplitude distribution due to the random "wander" of the waveform as shown in Figure 6. This wander, or "moving average", can be induced by patient variation, EKG apparatus drift, or both. Time series analysis (Ref 6) shows that a nonstationary process which has sample functions which are "locally stationary" or homogenous (Figure 5) can be represented by a process model which calls for the d'th difference of the process to be stationary. The proof of this assertion is given by Box and Jenkins (Ref 6:85-125).

If the nonstationary process sample sequence (time series) exhibits a random level as illustrated in Figure 5a, then a first difference operation will remove this "moving average" and force the resulting difference sequence to be centered around zero. If the waveform exhibits a random slope, as shown in Figure 5b, then a second difference operation will remove this quadratic "bias" with a corresponding reduction in amplitude distribution variance (Figure 7.) Comparison of Figures 5 and 6 shows that an EKG trace can "look" similar to the example in Figure 5b.

From the theory of stochastic processes (Ref 9:330-331), it is known that the expected value of a sample mean obtained by sampling a wide-sense stationary random process along a sample function in time is equal to the constant mean value of that random process.



Sample function with quadratic bias components.



Sample function after second difference operation.

Fig 7. Simulated variance reduction via second differencing.

In addition, the variance of the sample mean is inversely proportional to the number of samples taken when the samples are pairwise uncorrelated. Another fact about wide-sense stationary processes is that their autocorrelation functions are dependent only on the time difference between observation of the sample sequence and that  $R(t_1 - t_2) \leq R(0)$  for  $t_1 - t_2 \neq 0$ . Since the second difference is the first "difference" which can be modelled as coming from a "stationary" process, and the maximum correlation occurs for zero time difference, then the correlation must be reduced for adjacent sequence values in the second difference operation.

This heuristic argument is far from complete and does not explain why the third difference exhibits worse behavior than the second difference. The EKG waveform is a complex function from a complex source and higher order effects could begin to dominate with the third difference operation. Further analysis of this anomaly is left for future study.

The remainder of this chapter is organized in the following manner. First the Tolan algorithm is described, followed by a description of the Dower procedure. Next a short synopsis will be made of three other EKG data compression algorithms uncovered during the research of this thesis.

### Tolan EKG Data Compressor

The Tolan EKG data compression algorithm is a redundancy reduction procedure which processes a three lead EKG (VCG) and produces a compressed, digital output. This digital output sequence could subsequently be channel encoded for "errorless" transmission (Ref 27 and Appendix B) or stored for later retrieval and reconstruction.

As was the case for the RR techniques described in chapter 2, the Tolan compressor is subdivided into a data decorrelator and an entropy encoder. The decorrelator is discussed first.

Tolan Decorrelator. The Tolan decorrelator is a second order difference reduction operation which utilizes a time compression approach to form a decorrelated output sequence. A second order system was chosen based on the experimental results of Dower and Berghofer (Ref 12) and Cox and Ripley (Ref 7) which showed maximum compression gain with a second order difference operation. The Tolan decorrelator works on a three lead EKG (VCG) signal set, assumed to be sampled at a constant rate. The data compression is achieved in real time between successive samples.

The Tolan decorrelator algorithm is defined in Figure 8. Close examination of the algorithm in Figure 8 reveals that the second difference data is stored only if : 1) any of the three  $\Delta^2$  values are nonzero or ; 2) if the At counter

records more than 127 repetitive cases where  $\Delta^2 x = \Delta^2 y = \Delta^2 z = 0$ . Rule 1 was chosen due to the high degree of correlation between the leads of an EKG (i.e. when one lead is changing, so are the others). The capacity of the time counter in rule 2 was arbitrarily chosen to be sufficient to record the long quiescent periods which occur in the EKG (see Figure 1) and short enough to be efficiently stored.

Step 12 of the Tolan algorithm in Figure 8 calls the variable length encoding subroutine (Figure 10) which encodes the  $\Delta^2$  terms calculated by the second difference decorrelator. This encoding procedure is the next subject to be discussed.

Tolan Entropy Encoder. The Tolan code is an uniquely decodable variable length code which stores the  $\Delta^2$  values as a contiguous sequence of binary 1's. The length of this "run" of binary 1's is equal to the magnitude of the  $\Delta^2$  term. To delineate between the "runs", binary 0's are used as codeword delimiters. Since the second difference has both negative and positive values, a sign bit (0 for positive, 1 for negative) immediately follows the 0 bit delimiter. The three values  $\Delta^2 x, \Delta^2 y, \Delta^2 z$  are encoded and stored sequentially followed by a delimited, 7 bit, uncoded  $\Delta t$  value. A  $\Delta^2$  value of zero is indicated by 3 successive 0 bits.

```

1:  $\Delta x = 0$ ,  $\Delta y = 0$ ,  $\Delta z = 0$ ,  $\Delta t = 1$ 

2:  $x(i) = (a/d \text{ ch } 0)$ ,  $y(i) = (a/d \text{ ch } 1)$ ,  $z(i) = (a/d \text{ ch } 2)$ 

3:  $mem(0) = x(i)$ ,  $mem(1) = y(i)$ ,  $mem(2) = z(i)$ 

4:  $\tilde{x}(i+1) = x(i) + \Delta x$ ,  $\tilde{y}(i+1) = y(i) + \Delta y$ ,  $\tilde{z}(i+1) = z(i) + \Delta z$ 

5: if ready for next sample then GOTO 6 else GOTO 5

6:  $i = i + 1$ ,  $x(i) = (a/d \text{ ch } 0)$ ,  $y(i) = (a/d \text{ ch } 1)$ ,  $z(i) = (a/d \text{ ch } 2)$ 

7:  $\Delta^2 x = x(i) - \tilde{x}(i)$ ,  $\Delta^2 y = y(i) - \tilde{y}(i)$ ,  $\Delta^2 z = z(i) - \tilde{z}(i)$ 

8: if  $\Delta^2 x \neq 0$  or  $\Delta^2 y \neq 0$  or  $\Delta^2 z \neq 0$  then GOTO 11 else GOTO 9

9:  $\Delta t = \Delta t + 1$ 

10: if  $\Delta t \leq 127$  then GOTO 4 else GOTO 12

11:  $\Delta x = \Delta x + \Delta^2 x$ ,  $\Delta y = \Delta y + \Delta^2 y$ ,  $\Delta z = \Delta z + \Delta^2 z$ 

12: go subroutine coder {  $\Delta^2 x, \Delta^2 y, \Delta^2 z, \Delta t$  }

13: if memory is full then STOP else GOTO 14

14:  $\Delta t = 1$ 

15: GOTO 4

```

where

$x(i), y(i), z(i)$  = sampled, 8 bit precision, EKG data  
 $\Delta x, \Delta y, \Delta z$  = first difference { $\Delta x(n) = x(n) - x(n-1)$ }  
 $\Delta^2 x, \Delta^2 y, \Delta^2 z$  = second difference { $\Delta^2 x(n) = \Delta x(n) - \Delta x(n-1)$ }  
 $\Delta t$  = time difference between nonzero  $\Delta^2$  values  
 $\tilde{x}(i+1), \tilde{y}(i+1), \tilde{z}(i+1)$  = next predicted data points

Fig 8. Tolan Collection and Decorrelation Algorithm.

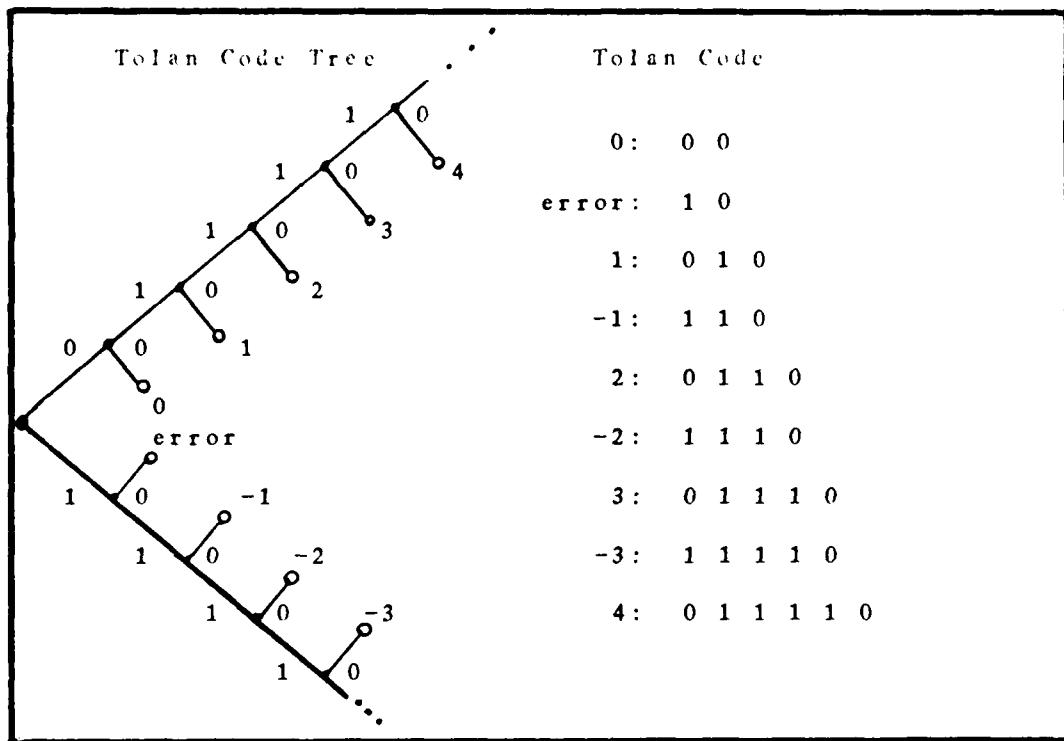


Fig. 9. Tolan Code and Code Tree (8 smallest codewords).

As an example of Tolan encoding, let  $\Delta^2 x = 3$ ,  $\Delta^2 y = -1$ ,  $\Delta^2 z = 0$ ,  $\Delta t = 13$ . With these values, then the code string generated is:

D S V V V D S V D S D T T T T T T T T D  
0 0 1 1 1 0 1 1 0 0 0 0 0 0 0 1 1 0 1 0

where

D = Delimiter bit

S = Sign bit

v = Value bit

T = Time bit

Examination of the above sequence shows that the  $\Lambda^2$

codewords are prefix codes as shown in the tree diagram of Figure 9. To decode the sequence correctly, however, it is imperative that codeword synchronization be maintained in order to determine when the coded  $\Delta^2$  values end and the uncoded, 7 bit  $\Delta t$  variable starts. This loss of synchronization is detected by the "error" state shown in the code tree of Figure 8. The decoder which implements the code tree in Figure 9 will not be discussed here but is listed in Appendix C.

In contrast with the Huffman code (see chapter 2 and Ref 16) and the Dower code to be discussed, the Tolan code is not constructed using the apriori knowledge of the source word (i.e.  $\Delta^2 x, \Delta^2 y, \Delta^2 z$ ) relative frequency of occurrence. This means that the Tolan code will only produce bit compression (bits out/bits in < 1) if the  $\Delta^2$  values are sharply peaked around a mean of zero such that few code words exceed the 8 bit, fixed length value of the  $\Delta^2$  terms. As will be shown in chapter 5, the second order difference decorrelator does produce such a sharply peaked relative distribution.

The algorithm which implements the Tolan variable length coder is shown in Figure 10. This algorithm works in conjunction with the Tolan decorrelator in Figure 8.

```

1: entry ( $\Delta^2 x$ ,  $\Delta^2 y$ ,  $\Delta^2 z$ ,  $\Delta t$ )
2: cnt  $\leftarrow$  0
3: reset current memory bit, increment bit pointers
4: if end of memory, set eom flag and RETURN else GOTO 5
5: cnt  $\leftarrow$  cnt + 1
6: if cnt=1 then tvar  $\leftarrow$   $\Delta^2 x$ 
7: if cnt=2 then tvar  $\leftarrow$   $\Delta^2 y$ 
8: if cnt=3 then tvar  $\leftarrow$   $\Delta^2 z$ 
9: if cnt  $\geq$  4 then GOTO 18 else GOTO 10
10: if tvar  $\geq$  0 then GOTO 11 else GOTO 12
11: reset memory bit, increment bit counters and GOTO 13
12: set memory bit, increment bit counters
13: if end of memory, set eom flag and RETURN else GOTO 14
14: if tvar=0 then GOTO 5 else GOTO 15
15: set memory bit, increment bit counters
16: if end of memory, set eom flag and RETURN else GOTO 17
17: tvar  $\leftarrow$  tvar - 1 and GOTO 14
17: store 7 bit  $\Delta t$  counter to memory, update bit pointers
18: if end of memory, set eom flag and RETURN else RETURN

```

Fig 10. Tolan Variable Length Encoder Algorithm.

```

0: START OF POWER DATA COLLECTION AND DECORRELATION

1: FLAG=TTT=DT=DX1=DY1=DZ1=DX2=DY2=DZ2=DDX=DDY=DDZ=CNT=0

2: If sample interrupt detected then GOTO 3 else GOTO 2

3: X(n) ← A/D CH 0, Y(n) ← A/D CH 1, Z(n) ← A/D CH 2

4: CNT=CNT + 1

5: If CNT =1 then GOTO 6 else GOTO 7

6: X(n-1) ← X(n), Y(n-1) ← Y(n), Z(n-1) ← Z(n) and GOTO 2

7: If CNT =2 then GOTO 8 else GOTO 13

8: DX1 ← X(n)-X(n-1),DY1 ← Y(n)-Y(n-1),DZ1 ← Z(n)-Z(n-1),DT ← DT+1

9: If DX1≠0 OR DY1≠0 OR DZ1≠0 OR DT > 127 then GOTO 11 else GOTO 10

10: CNT ← 1 and GOTO 2

11: TDT ← DT-63, X(n-2) ← X(n-1), Y(n-2) ← Y(n-1), Z(n-2) ← Z(n-1)

12: X(n-1) ← X(n), Y(n-1) ← Y(n), Z(n-1) ← Z(n) DT ←0 and GOTO 2

13: DX2 ← X(n)-X(n-1),DY2 ← Y(n)-Y(n-1),DZ2 ← Z(n)-Z(n-1),DT ← DT+1

14: If DX2≠0 OR DY2≠0 OR DZ2≠0 OR DT > 127 then GOTO 16 else GOTO 15

15: CNT ← 2 and GOTO 2

16: DDX ← DX2-DX1, DDY ← DY2-DY1, DDZ ← DZ2-DZ1

17: If |DDX| > 63 then GOTO 18 else GOTO 21

18: DUMX(n)=.25X(n-1)+.75X(n)

    DUMX(n-1)=.25X(n-2)+.5X(n-1)+.25X(n)

    DUMX(n-2)=.75X(n-2)+.25X(n-1)

19: X(n) ← DUMX(n), X(n-1) ← DUMX(n-1), X(n-2) ← DUMX(n-2)

20: DX1=X(n-1)-X(n-2), DX2=X(n)-X(n-1) and GOTO 16

```

Fig 11-a. Power Collection and Decorrelation Algorithm.

```

21: If |DDY| > 63 then GOTO 22 else GOTO 25

22: DUMY(n)=.25Y(n-1)+.75Y(n)

    DUMY(n-1)=.25Y(n-2)+.5Y(n-1)+.25Y(n)

    DUMY(n-2)=.75Y(n-2)+.25Y(n-1)

23: Y(n) ← DUMY(n), Y(n-1) ← DUMY(n-1), Y(n-2) ← DUMY(n-2)

24: DY1=Y(n-1)-Y(n-2), DY2=Y(n)-Y(n-1) and GOTO 16

25: If |DDZ| > 63 then GOTO 26 else GOTO 29

26: DUMZ(n)=.25Z(n-1)+.75Z(n)

    DUMZ(n-1)=.25Z(n-2)+.5Z(n-1)+.25Z(n)

    DUMZ(n-2)=.75Z(n-2)+.25Z(n-1)

27: Z(n) ← DUMZ(n), Z(n-1) ← DUMZ(n-1), Z(n-2) ← DUMZ(n-2)

28: DZ1=Z(n-1)-Z(n-2), DZ2=Z(n)-Z(n-1) and GOTO 16

29: If FLAG=0 then GOTO 30 else GOTO 31

30: VLC ← (0:DX1,DY1,DZ1)

31: VLC ← (TDT:DDX,DDY,DDZ)

32: If MEMORY FULL then STOP else GOTO 33

33: X(n-2) ← X(n-1), Y(n-2) ← Y(n-1), Z(n-2) ← Z(n-1)

34: X(n-1) ← X(n), Y(n-1) ← Y(n), Z(n-1) ← Z(n)

35: TDT ← DT-63, FLAG=1 and GOTO 2

```

Fig 11-b. Power Collection and Decorrrelation Algorithm.

### Dower EKG Data Compressor

The Dower EKG compression technique, like the Tolan method, is a redundancy reduction procedure. The Dower compressor combines a zero order time compression operation with a second order difference reduction transformation to produce a decorrelated residual frame sequence. This residual frame sequence is then compressed by a specially tailored variable length code whose average codeword bit length approaches the entropy bound without the buffer overflow problem encountered with the "optimal" Huffman code.

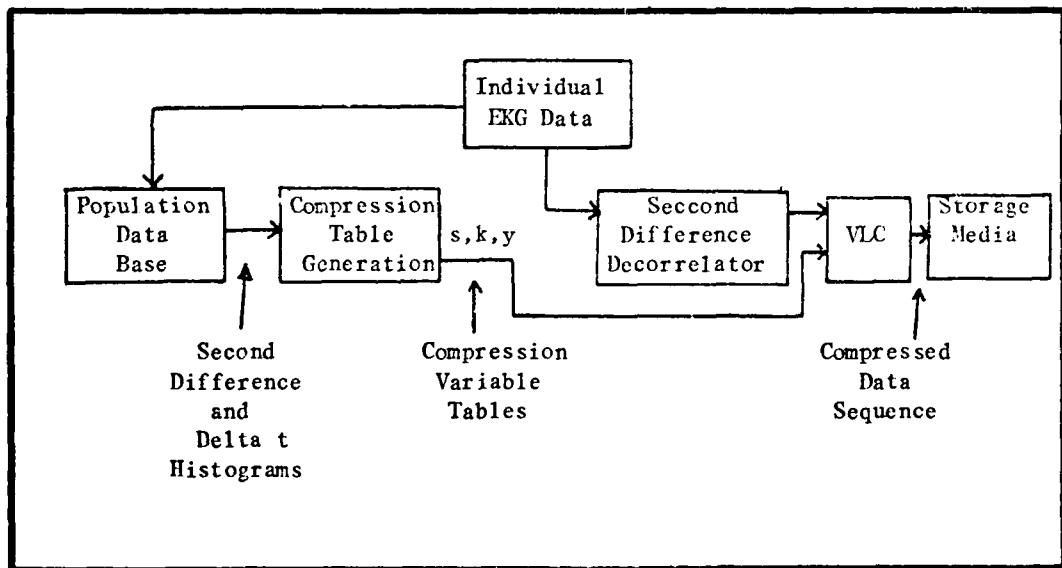


Fig 12. Dower EKG Data Compression System.

Dower Decorr elator. The Dower data collection and decorr elation operation (Figure 11) begins with a zero order time compression process where only those EKG sample values ( $x, y, z$ ) which differ from the previous sample ( $\Delta x$  or  $\Delta y$  or  $\Delta z \neq 0$ ) are saved along with a run length counter  $\Delta t$ . This operation produces a sequence of data frames ( $\Delta t: x, y, z$ ). Following the creation of the data frames, a second difference with respect to frame number is performed generating a sequence of second order difference frames ( $\Delta t: \Delta_f^2 x, \Delta_f^2 y, \Delta_f^2 z$ ).

The Dower algorithm operates on 8 bit data, hence there are 256 potential source symbols for each lead. When second differences are taken, the range increases to 1024 potential  $\Delta^2$  symbols. By experimental evidence, Dower and Berghofer (Ref 12) have found that  $0 \pm 63\Delta^2$  values are sufficient to reproduce all but the fastest EKG artifacts (e.g., pacemaker spikes). The variable length encoder, therefore, is designed to expect 127 source symbols (values) and no more. To insure the  $\Delta^2$  dynamic range of  $0 \pm 63$  is not exceeded, the  $\Delta^2$  values are limited by a preprocessor shown in steps 18-26 of Figure 11. This preprocessor is iterative, and irreversably modifies the three sample points which produced the  $|\Delta^2| > 63$  until the  $\Delta^2$  value falls within the encoder range. With care taken to record the necessary initial conditions (first  $x, y, z$  and first  $\Delta x, \Delta y, \Delta z$ ), this second difference frame sequence now contains all the significant information in the original sample sequence with

reduced interframe correlation.

The mechanics of the Power decorrelator are best illustrated by an example. Figure 13 shows three hypothetical sample sequences. By applying the rules in the preceding paragraph, a set of zero order data frames ( $\Delta t : x, y, z$ ) is formed. That is:

(0:0,-1,2),(0:1,1,1),(0:3,2,3),(0:4,3,1),(3:3,1,-3),(0:1,0-1),  
(0:-1,-1,-2),(0:-2,-2,-1),(4:-3,1,-1)

To complete the correlation reduction process, the second difference with respect to frame number must now be performed. This results in the set of second difference frames ( $\Delta t : \Delta_f^2 x, \Delta_f^2 y, \Delta_f^2 z$ ) shown below.

(0:0,-1,2),(0:1,2,-1) | (0:1,-1,3),(0:-1,0,0),(3:-2,-3,-2),  
(0:-1,1,6),(0:0,0,-3),(0:1,0,2),(4:0,4,-1)

The data sets preceding the verticle bar are the initial conditions necessary for reconstruction. The first data set is the first zero order time compression sample frame. The second frame is the first difference between the first two zero order data frames.

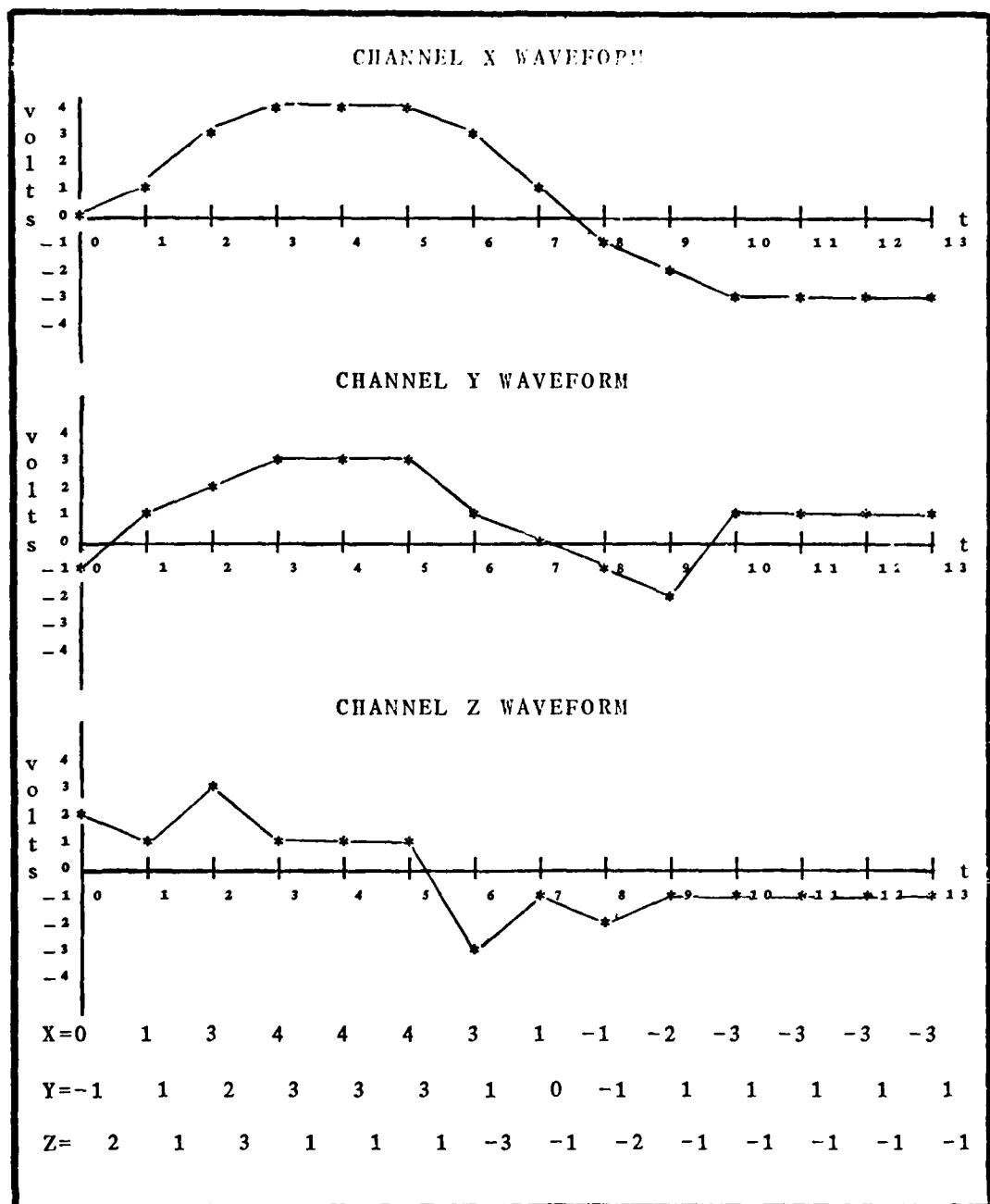


Fig 13. Three hypothetical sample sequences.

Dower Entropy Encoder. The Dower entropy encoder also uses variable length codes for the second stage of the EKG data compressor. The Dower code, however, is significantly different than the prefix code used in the Tolan compression method.

As was mentioned previously, the Dower VLC is configured for 127 source symbols (8 bit  $\Delta^2$  values). Unlike the Tolan code, however, the Dower VLC does not perform a 1:1 mapping between a single source symbol and a single code symbol.

In the Dower coder, the source symbols ( $\Delta^2_i$ ) are mapped onto the state space of a 14 bit accumulator (A). This state space is partitioned into symbol regions  $R_i$  which are assigned according to the probability of occurrence of the source symbols. The size ( $||R_i||$ ) of the symbol region  $R_i$  is given by the relation

$$||R_i|| = ds_i = s_{i+1} - s_i = \Pr(\Delta^2_i) 16384 \quad (24)$$

where  $s_i$  is the initial state of  $R_i$  and  $s_{i+1}$  is the initial state of  $R_{i+1}$ . The initial state,  $s_i$ , is also determined by symbol probability. The symbol regions  $R_i$  corresponding to those symbols which occur least often are assigned to the low end of the state space range such that if:

$$\Pr(\Delta^2_k) > \Pr(\Delta^2_n) > \Pr(\Delta^2_m) > \Pr(\Delta^2_j)$$

then

$$||R_k|| > ||R_n|| > ||R_m|| > ||R_j||$$

and

$$s_k > s_n > s_m > s_j$$

This relationship is illustrated in Figure 14. For those potential symbols  $\Delta_i^2$  which do no occur, 1 state is always assigned.

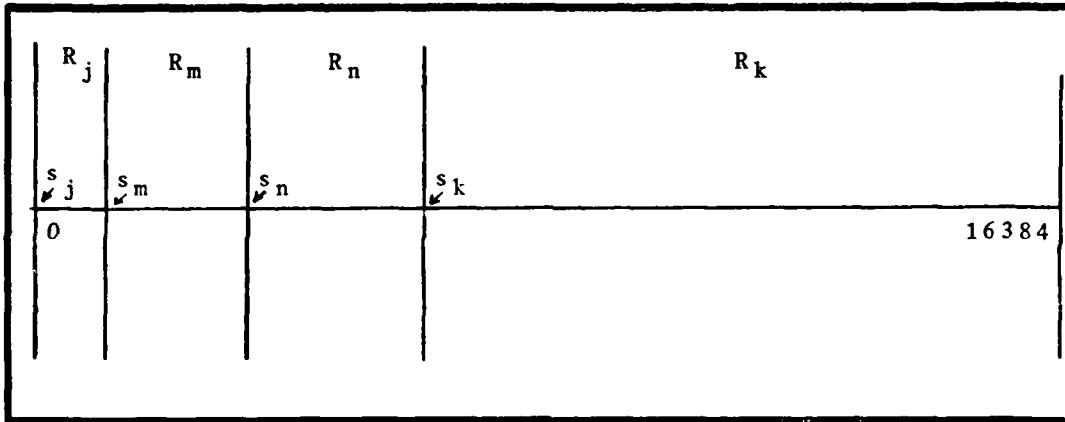


Fig 14. Example of State Space Partition for 4 Symbols.

The physical configuration of the Dower encoder is shown in Figure 15. As can be seen from Figure 15, the memory storage buffer for the encoded data is contiguous with accumulator and all accumulator bit shifts (right or left) also shift the entire memory buffer. Right shifting operations always shift a binary 0 into the most significant bit (MSB) of A.

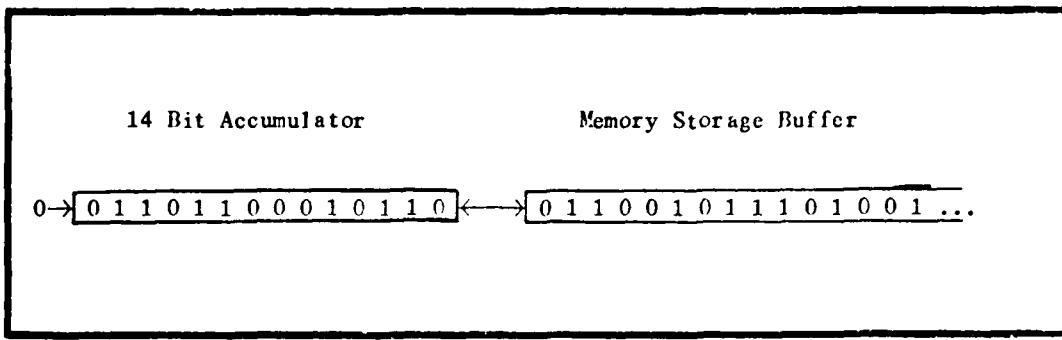


Fig 15. Dower Encoder Accumulator-Memory Buffer Interface.

A source symbol  $\Delta_i^2$  is encoded by adding the value  $s_i$  to the current contents of the accumulator. If the value in the accumulator is such that  $s_i + A \geq 2^{14}$ , then an arithmetic overflow condition would occur with a corresponding loss of data. To insure this overflow does not occur, the accumulator is first right shifted  $k_i$  times into memory. The value  $k_i$  is the maximum number of right shifts which would ever be necessary and can be found from the inequality:

$$2^{14} \leq 2^{k_i(s_{i+1}-s_i)} < 2^{15}. \quad (25)$$

The  $k_i$  values are tabulated in Table I for various values of  $ds = s_{i+1} - s_i$ . This right shifting diminishes the value in the accumulator (i.e., divides A by 2 per right shift) until the value in A can be safely mapped into a region  $R_i$ . The initial state,  $s_i$  can now be added to A without overflow occurring.

In some instances, the maximum shift  $k_i$  would not be

necessary (i.e.,  $k_i - 1$  right shifts would still prevent overflow). To determine when this situation occurs, a third variable  $y_i$  is defined where

$$y_i = ds - 2^{14-k_i} \quad (26)$$

represents the excess states by which  $ds$  exceeds the next lowest integral power of 2.

Table I  
Dower Code Variables Per Accumulator Partition Size

Partition Size	k	y
$ds=1$	14	0
$1 < ds < 4$	13	$ds - 2$
$4 < ds < 8$	12	$ds - 4$
$8 < ds < 12$	11	$ds - 8$
$16 < ds < 32$	10	$ds - 16$
$32 < ds < 64$	9	$ds - 32$
$64 < ds < 128$	8	$ds - 64$
$128 < ds < 256$	7	$ds - 128$
$256 < ds < 512$	6	$ds - 256$
$512 < ds < 1024$	5	$ds - 512$
$1024 < ds < 2048$	4	$ds - 1024$
$2048 < ds < 4096$	3	$ds - 2048$
$4096 < ds < 8192$	2	$ds - 4096$
$8192 < ds < 16384$	1	$ds - 8192$

If the value in  $A < y_i$  after the initial  $k_i$  shifts, then one too many right shifts occurred.  $A$  is then left shifted once and  $s_i$  added to  $A$  without fear of overflow. If  $A \geq y_i$  then a full  $k_i$  right shifts were required. To insure maximum efficiency in the next encoding, however, the position of the accumulator in state space map should be

right justified within the current symbol region. This right justification is obtained by adding  $y_i + s_i$  to A whenever  $A \geq y_i$ .

By use of the variable  $y_i$  maximum efficiency in the number of right shifts (i.e. information bits stored to memory) can be obtained. Dower and Berghofer (Ref 12) assert that, on the average, the number of right shifts per source symbol ( $\Delta_i^2$ ) will approach quite closely to the value  $-\log_2 \Pr(\Delta_i^2)$  which is the self information "content" of the  $\Delta_i^2$  value (see appendix B).

The algorithm which implements the Dower encoder is shown in Figure 15. This algorithm assumes that the variable tables for  $\Delta_f^2 x, \Delta_f^2 y, \Delta_f^2 z$  and  $\Delta t$  have been calculated from the second difference histograms.

The Dower decoding operation is just the inverse of the encoding process. For decoding, the current state of the accumulator is mapped into the state space table of the current frame variable ( $\Delta t, \Delta_f^2 x, \Delta_f^2 y, \Delta_f^2 z$ ). Once the correct region  $R_i$  is determined, then  $s_i$  is subtracted from A. If  $A < 2y_i$  then one right shift is performed followed by  $k_i$  left shifts. If  $A \geq 2y_i$ , then  $y_i$  is subtracted from A followed by  $k_i$  left shifts. For a more detailed example of the Dower encoding/decoding operation, the reader is referred to Dower and Berghofer (Ref 12).

Although conceptually more difficult than the Tolan VLC, the algorithmic structure of the Dower entropy encoder is only slightly more complex. As was shown above, the

Dower VLC can be implemented solely with shifting and table look up operations. The Dower method does, however, require that code tables ( $s_i, k_i, y_i$ ) be constructed prior to the encoding process.

```

1: Entry ( $\Delta t, \Delta_f^2 X, \Delta_f^2 Y, \Delta_f^2 Z$ )
2: IF this is the first frame stored then Acc $\leftarrow 0$  else GOTO 3
3: DDVAR  $\leftarrow \Delta t, k \leftarrow k(\Delta t), s \leftarrow s(\Delta t), y \leftarrow (\Delta t)$ 
4: RIGHT SHIFT Acc k times
5: IF Acc < y then GOTO 8 else GOTO 6
6: LEFT SHIFT Acc once
7: Acc  $\leftarrow$  Acc + s and GOTO 9
8: Acc  $\leftarrow$  Acc + s + y
9: IF memory is full, then SET eom flag and RETURN else GOTO 10
10: IF DDVAR= $\Delta t$  then GOTO 11 else GOTO 12
11: DDVAR  $\leftarrow \Delta_f^2 X, k \leftarrow k(\Delta_f^2 X), s \leftarrow s(\Delta_f^2 X), y \leftarrow y(\Delta_f^2 X)$  and GOTO 4
12: IF DDVAR= $\Delta_f^2 Y$  then GOTO 13 else GOTO 14
13: DDVAR  $\leftarrow \Delta_f^2 Y, k \leftarrow k(\Delta_f^2 Y), s \leftarrow s(\Delta_f^2 Y), y \leftarrow y(\Delta_f^2 Y)$  and GOTO 4
14: IF DDVAR= $\Delta_f^2 Z$  then GOTO 15 else RETURN
15: DDVAR  $\leftarrow \Delta_f^2 Z, k \leftarrow k(\Delta_f^2 Z), s \leftarrow s(\Delta_f^2 Z), y \leftarrow y(\Delta_f^2 Z)$  and GOTO 4

```

Fig 16. Dower VLC algorithm.

### Tolan versus Dower

The similarities and differences between the Tolan and Dower EKG compression techniques are now discussed. The decorrelators are compared first.

Both the Dower and the Tolan decorrelators use a second difference technique where time compression is used to eliminate the storing of the most common  $\Delta^2$  value of zero. The Tolan algorithm, however, continuously calculates the second difference and applies these differences to the Tolan VLC. If a long "run" of  $\Delta^2 = 0$  values occur, the Tolan encoder overflows at  $\Delta t = 128$ . This run counter overflow forces a storage "dump" to the VLC with a resultant loss in efficiency.

The Dower decorrelator approaches time compression slightly differently. In the Dower algorithm, only those sample points where the EKG data was changing (i.e.,  $\Delta x \text{ or } \Delta y \text{ or } \Delta z \neq 0$ ) are saved forming data frames. A second difference with respect to frame number is performed and the  $\Delta_f^2$ 's and  $\Delta t$  are fed to the VLC. Although this frame methodology appears more efficient than the Tolan technique, the Dower decorrelator is still constrained by a maximum  $\Delta t$  of 127 (i.e., maximum Dower VLC code range). On the basis of this analysis, it appears that both the Tolan and the Dower decorrelators have similar performance.

Since the Tolan and Dower decorrelators appear about equally efficient, the real compression payoff is in the variable length encoders. In the Tolan VLC, code words as

long as  $1024+2$  bits are feasible. These long codewords would only occur if an extremely large and fast "spike" (e.g. a pacemaker pulse) appeared in the data. In the population as a whole, spikes of this magnitude occur very infrequently. Nevertheless, the Tolan compression efficiency will degrade seriously in an environment where "impulsive" artifacts appear.

The Dower VLC is designed to minimize sensitivity to impulses. In the Dower system, the VLC is configured such that the longest codeword (i.e., the number of right shifts to memory) is 14 bits. Extremely large signal spikes are numerically filtered to reduce their second difference within the  $0 \pm 63$  range. Such drastic limiting action would occur infrequently, however, and not seriously affect the reproduction of the EKG.

As an example of their performance, let a second difference of 45 be encountered by both the Tolan and Dower encoders. The Tolan VLC would require 47 bits to encode this data. The exact codeword size of the Dower routine is dependent on the  $\Delta^2 = 45$  probability of occurrence. Nonetheless, the "codeword" is always  $\leq 14$  bits long; an obvious increase in efficiency over 47 bits.

A calculation of the performance differential between the Tolan and Dower routines will be estimated in chapter 5. For now it is sufficient to say that the Dower EKG compressor outperforms the Tolan method. A short synopsis will now be made of two other EKG compression techniques.

## CHAPTER 4 Microcomputer EKG Compressors.

Two other EKG compression techniques were discovered by this author. The first is referred to as the Turning Point Method and is discussed in reference 32. The second method is currently in use by Marquette Electronics for data compression in a commercial EKG "cart". The Turning Point technique is presented first.

Turning Point Algorithm. The turning point algorithm is, by definition, a 2:1 data compressor where one of two consecutive sample points is discarded. The algorithm which determines which sample point is discarded is as follows. The first sample point is stored and assigned as the reference point ( $X_0$ ). The next two consecutive points become  $X_1$  and  $X_2$ . With 3 sample points there are 8 "patterns" or combinations which reflect the "trends" in data (see Figure 17). The Turning Point algorithm stores the circled point (Figure 17) which becomes the new reference point  $X_0$ . The point not circled ( $X_1$  or  $X_2$ ) is discarded. The next two points are sampled, their values are assigned to  $X_1$  and  $X_2$ , and the process repeated.

PATTERN	$x_0$	$x_1$	$x_2$
1		•	○
2		○	•
3		•	○
4	○	•	
5	○	•	○
6	•	•	○
7			○
8	•	•	
9	•	•	○

Fig 17. Turning Point Patterns (from Ref 32:6,61).

It can be shown (Ref 32:6.59-6.65) that :

if  $(X_2 - X_1) * (X_1 - X_0) < 0$        $X_0 = X_1$

if  $(X_2 - X_1) * (X_1 - X_0) \geq 0$        $X_0 = X_2$

hence the sign of the product of consecutive first differences determines the "significant" point to be saved.

The advantage of the Turning Point compressor is speed of execution. Execution of this algorithm on even the slowest microprocessor would pose no problem at all. The disadvantages of the turning point routine are numerous, however. First the routine discards data so that reconstruction of the original sample sequence is not possible. In this regard, the Turning Point algorithm falls in the class of Entropy Reducing (ER) techniques as was discussed in chapter 2. Second no attempt is made to use the probabilistic distribution of the EKG to enhance data compression as is done with the Tolan and Dower routines. This obviously leads to inefficiency. Finally, this routine only produces a 2:1 compression ratio; incredibly poor in relation to the other techniques already discussed in this thesis. Only where the simplest technique is necessary, would the Turning Point technique be beneficial.

Marquette Algorithm. The Marquette algorithm was developed by Marquette Electronics, Milwaukee, Wisconsin. This similar to the Dower and Tolan algorithms in that variable length encoding is used to compress the output of a

"difference" decorrelator.

In contrast to the Dower and Telan techniques, the Marquette compressor only calculates the first difference for input into the VLC. The Marquette variable length encoder stores the first difference data as 1,3,5, or 7 nibbles where a nibble is defined as 4 bits. These nibble codes are arranged as follows:

<u>Range of Difference</u>	<u>Code Length in Nibbles</u>
(-7,7)	1
(-127,127)	3
(-2047,2047)	5
(-32767,32767)	7

To encode the first differences, and delineate between code words, the following rules apply:

- 1) Differences must be coded on the smallest possible range, and attempts to encode a difference of +5, for example, using more than 1 nibble will result in a decoding error.
- 2) Single nibble codes are difference plus 8. A nibble value of zero does not occur.
- 3) Three nibble codes start with a single zero nibble. The remaining two nibbles are obtained as follows:
  - (1) Positive Differences +8
  - (2) Negative Differences +7
- 4) Five nibble codes start with two zero nibbles. The remaining three nibbles are obtained as follows:
  - (1) Positive Differences +128
  - (2) Negative Differences +127
- 5) Seven nibble codes start with three zero nibbles. The remaining four nibbles are obtained as follows:

(1) Positive Differences +2048  
(2) Negative Differences +2047

6) The sequence of nibbles in a code starts with the zero flag nibbles (if any) followed by the most significant through least significant nibble.

The above set of encoding rules were obtained from Mr. Tom Divers, Marquette Electronics project engineer (Ref 11).

From their own analysis, Marquette has shown (Ref 11) that at a 250 Hz sampling rate, 89.1 percent of the first differences fall within the  $\pm 7$  range with 99.8 percent falling within a  $\pm 127$  range. Marquette reports that at an A/D precision of 10 bits, an average of 4.89 bits/sample (across the total EKG population) is obtained with their compression routine.

The Marquette EKG compression appears to work well, even with a first order difference correlation reducer. The Marquette variable length encoder, however, is tailored to the 8 bit ASCII data communications environment and is not "optimum" in any sense. The Marquette VLC does perform "exact" redundancy reduction entropy compression with sufficient efficiency to make this encoding scheme commercially viable.

### Chapter III Summary.

This chapter has looked in detail at two EKG data compression techniques which perform "exact" redundancy reduction. The Tolan routine, which was implemented by this author (see chapters 4 and 5) decorrelated the sample

sequence data by performing a second order time compression operation. The residual sequence resulting from the decorrelator was "compressed" by a uniquely decodable variable length code which was shown to be "suboptimal" with respect to the "optimal" Huffman code.

The Dower compression technique also performed a second difference operation, but preceded the second order "differencer" by a zero order time compressor. The Dower zero order time compressor produced a sequence of "data frames" which in turn were converted to a sequence of second order difference frames with respect to frame number. Since the Dower VLC limited the  $\Delta t$  time compression counter to a maximum of 127 (as did the Tolan VLC), input symbols, the Dower and Tolan decorrelators were considered to perform equally well. The Dower entropy encoder was also shown to be a variable length coding operation but not a "prefix" code as was the Tolan VLC. The Dower VLC maps decorrelator "symbols" into the "state space" of 14 bit accumulator which encodes data by adding the initial address of a symbol's state space region to the accumulator. To prevent accumulator overflow, the accumulator data is shifted out to a memory storage buffer. The "number of shifts" necessary to prevent accumulator overflow represent the codeword size and it was shown that the Dower VLC approached the "entropy" bound of the decorrelated input sequence.

The chapter was concluded by a synopsis of two other

EKG compression techniques. The first of these two, the Turning Point algorithm, was shown to be of marginal use because it is not an "exact" technique and produces compression of only 2:1. The second technique, however, was the Marquette compression system and it was shown to produce acceptable compression worthy of commercial application.

From the available algorithms, the Dower compression technique has the capacity to produce the best "exact" compression of any of the techniques studied in this chapter. An interesting experiment would be the combining of the second order interpolator (discussed in chapter 2) used by Ruttiman and Pipberger (Ref 28) and the variable length encoder used by Dower (Ref 12). This combination should prove to be very powerful and effective and is left for further study.

The next chapter in this thesis discusses the configuration of the EKG Data Acquisition and Analysis System assembled by this author to test the Tolan EKG compression algorithm.

#### IV. EKG-Data Acquisition and Analysis System

The EKG-Data Acquisition and Analysis System (EKG-DAAS) was assembled for this thesis as the testbed on which experimental EKG data could be acquired, compressed, analyzed, stored, and reconstructed. The EKG-DAAS design can be separated into the categories of hardware and software. The hardware is discussed first.

##### EKG-DAAS Hardware.

The EKG-DAAS was constructed around the Motorola Exorciser microcomputer (appendix E) which uses a 6800 microprocessor for its central processing unit (CPU). In the EKG-DAAS, the Exorciser is configured with 32 kilobytes (K) of read/write (RAM) memory and 16 K of read only memory (ROM). In addition, the Exorciser was equipped with the EXBUG debugging module which allowed interactive program debugging with preselectable software breakpoints, execution tracing modes, and CPU register display.

To provide extended memory, a Midwest Scientific Instruments (MSI) FD-8 Disk Memory unit was interfaced to the Exorciser and provides approximately 290 K of online user memory. The FD-8 is accessed by a MSI Disk Operating System (DOS) and communicates with the CPU via a MEX6820 Input/Output Module installed in the Exorciser chassis.

Terminal input/output (I/O) is accomplished by means of a standard RS-232 serial interface which has switch selectable baud rate from 110 to 9600 baud. In the EKG-DAAS configuration, the Exorciser serial "port" was connected in parallel with a Heathkit H-14 dot matrix line printer for program listings and data printouts.

Data I/O was accomplished by means of a Sinetrac ST-6800 Analog/Digital-Digital/Analog (A/D-D-A) converter module which samples and digitizes analog data to 12 bit precision. The ST-6800 has the capability of sampling 32 distinct analog channels (A/D) as well as output 2 channels (D/A) with simple memory addressed LDA (load) and STA (store) instructions. For the EKG-DAAS, the ST-6800 was setup for  $\pm$  5 volt, 2's compliment data and was addressed (A/D ch 0) at E400 Hexadecimal (Hex). The internal configuration of the Exorciser as used in the EKG-DAAS is illustrated in Figure 18.

To uniformly sample the EKG input data via the ST-6800 required the use of an external interrupt timer as is shown in Figure 19. This timer allowed data sampling rates between 300 and 700 Hertz but for the duration of this research was set, and calibrated, at 500 Hertz. The interrupt was interfaced to the Exorciser via an interrupt line on the ST-6800.

Considerable problems arose in this thesis due to hardware problems associated with the FD-8 Disk Memory. The original configuration of the EKG-DAAS used two FD-8 systems

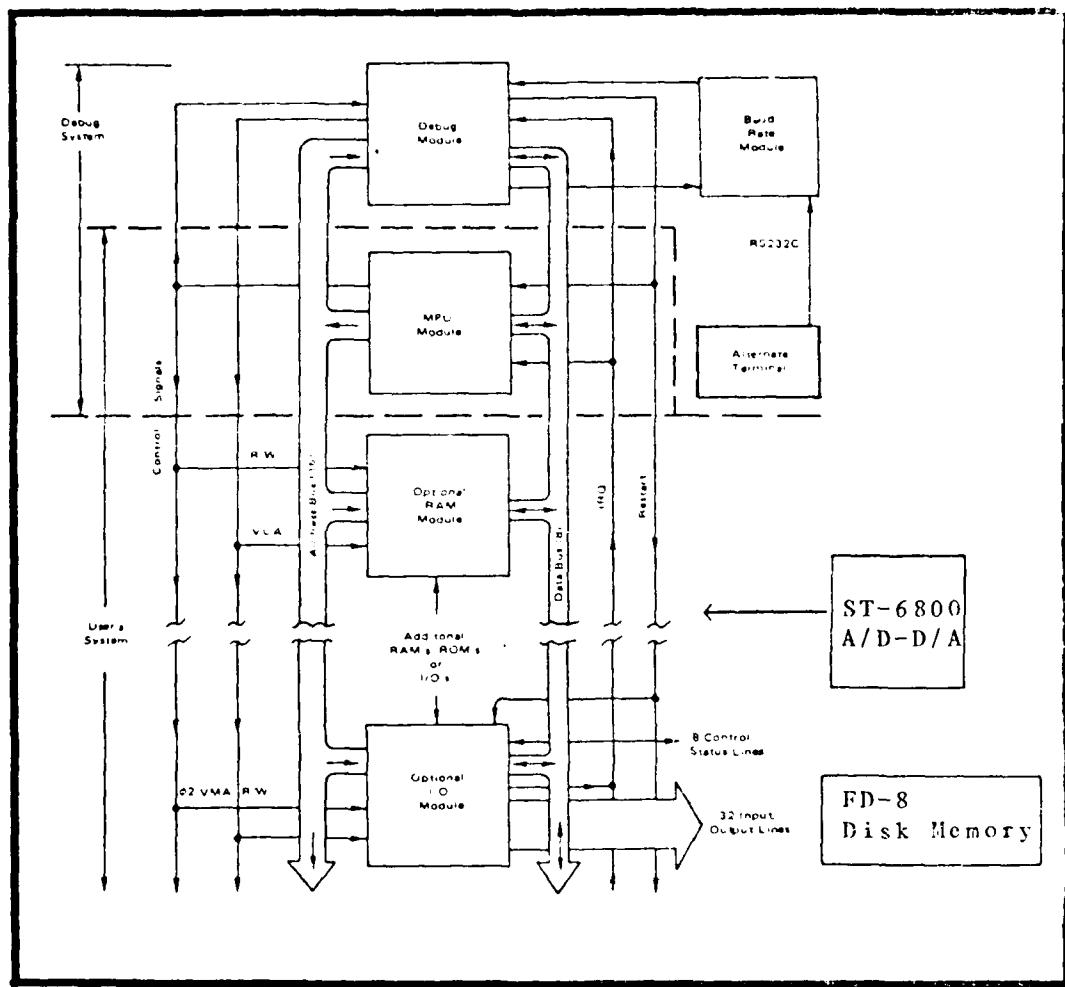


Fig 18. EXORCISER Component Module Layout (From Ref 22).

but one failed about midway through the software development. This failure caused a major rewrite of the thesis software and destroyed several weeks of work. Although the EKG-DAAS can now operate in a one disk environment, considerable "manhandling" of the data diskettes is necessary.

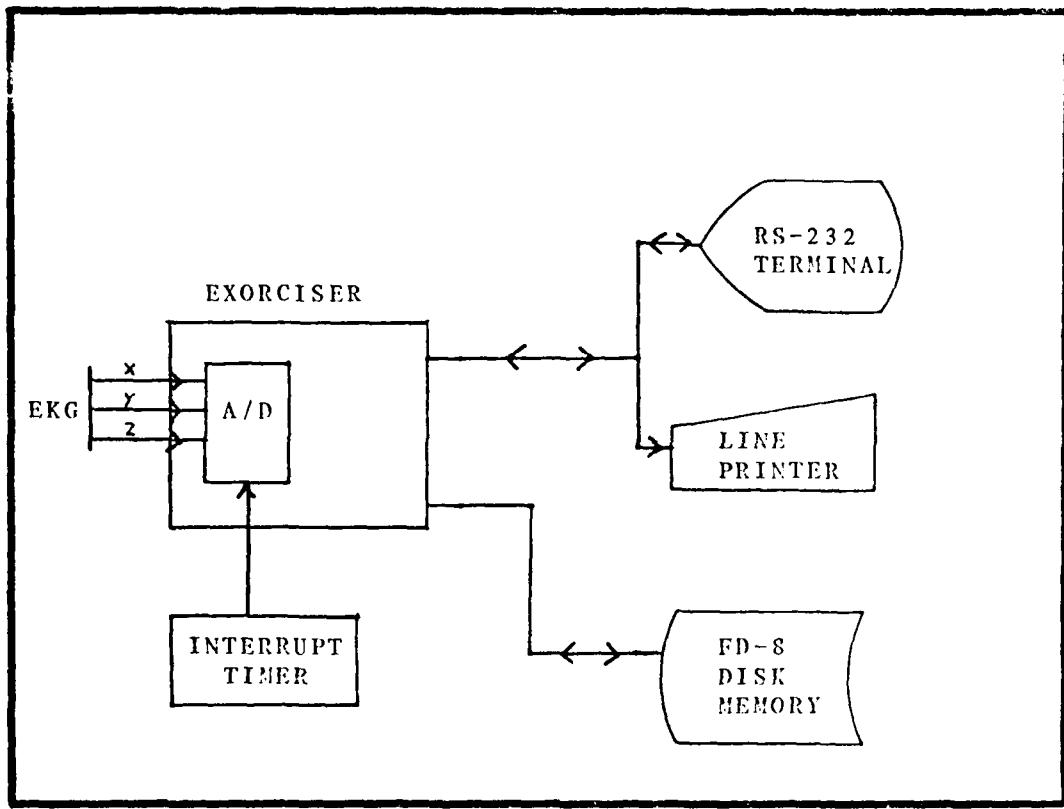


Fig 19. EKG-DAAS Hardware Configuration

Further information concerning the Exorciser's hardware capability can be obtained in appendix E and reference 22. Attention is now turned to the EKG-DAAS software.

#### EKG-DAAS Software.

The EKG-DAAS software was written in 6800 assembly language and controls all aspects of terminal, disk memory, and A/D-D/A operation. The EKG-DAAS programs consist of approximately 4300 lines of assembly language and are listed in appendix C.

The software used in the EKG-DAAS was written in 6800

assembly language for two reasons: 1) no high order language (e.g., PASCAL, FORTRAN) is available on the Exerciser and; 2) speed limitations imposed by the thesis requirement for online, real time EKG data compression made it imperative that the compression programs run as fast as possible. The laborious task of writing and testing assembly language slowed software development to the point where only one EKG data compression/reconstruction routine (Tolan) was completed.

The EKG-DAAS software is integrally tied to the MSI DOS (Ref 23). All disk I/O operations initiated by the EKG-DAAS routines flow through the MSI-DOS and hence the DOS must be "live" somewhere in memory. To insure that the DOS routines are always available, the DOS was disassembled and relocated in high memory ROM (C400 Hex).

The basic flow of EKG-DAAS program control is illustrated in Figure 20. The EKG-DAAS software is broken into overlaid modules which are called into memory and executed by EKG-EXEC and DISPLAY. The basic memory map and overlay structure is illustrated in Figure 21.

As can be seen in Figure 21, extensive memory management was required in order to allow a sufficiently large memory buffer for the EKG data. As configured in Figure 21, the EKG-DAAS could collect 11.6 seconds of uncompressed (3 leads, 8 bits/lead, 500 samples/sec) EKG data. For compression with the Tolan algorithm, a maximum of 26.2 seconds of data (TA1359PA, appendix D) was

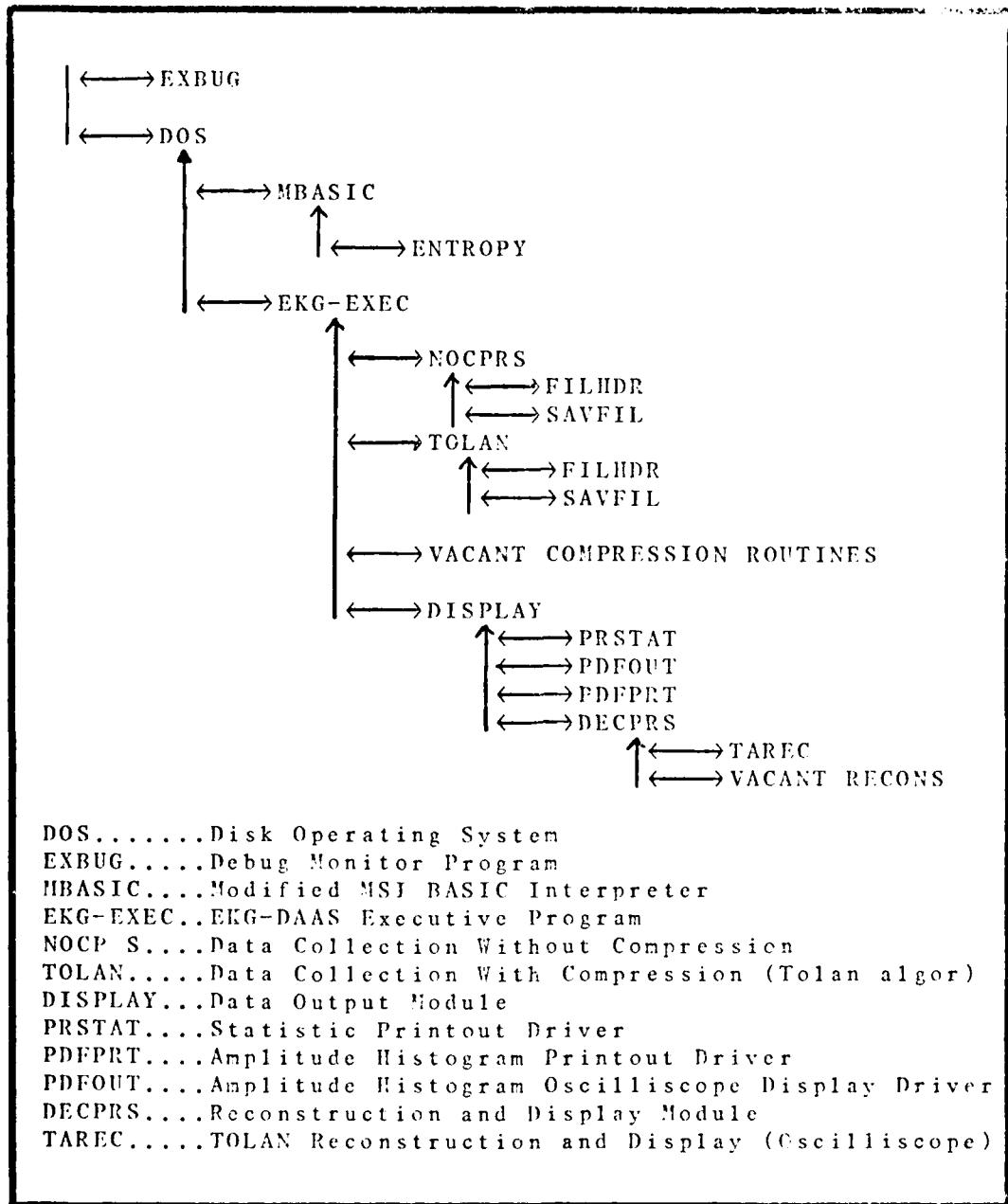


Fig 20 EKG-DAAS Software Control Flowgraph.

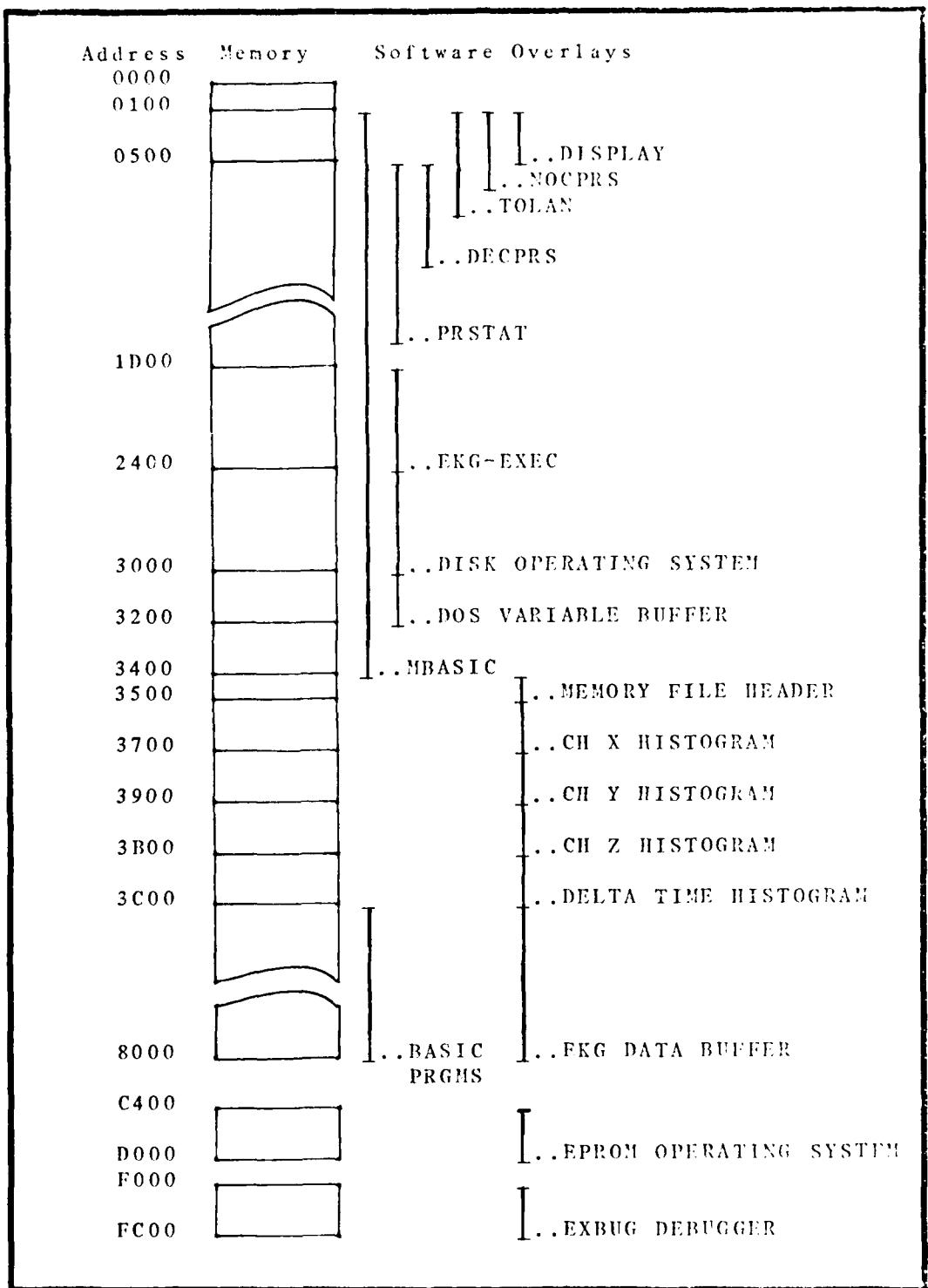


Fig. 21. EKG-DAAS Overlay Structure and Memory Map.

collected.

With the basic structure of the EKG-DAAS software defined, this chapter will now examine the components of the EKG-DAAS software in more detail. In the next section of this chapter, a software module will be described along with a simplified flowchart of that module's operation.

EKG-EXEC. EKG-EXEC is the executive command module which controls the execution flow of the EKG-DATA system. Upon input of a command number, EKG-EXEC loads the appropriate overlay routine into the program work buffer (0100-1D00 Hex) and then passes control to that overlay. The above command and control operation is illustrated in Figure 22.

In addition to the command "handler", EKG-EXEC contains the utility subroutines FILHDR, SAVFIL, HXASC, OVRLAY, and PDFPRT. These subroutines are described as follows:

FILHDR. FILHDR clears the memory data buffer, initializes the statistics buffer variables, and queries the console for data such as FILENAME, SUBJECT, DATE, etc.

SAVFIL. SAVFIL reads the filename in the memory buffer header and then writes the memory file to disk. Disk I/O is passed through EOS subroutines.

HXASC. This subroutine converts hexadecimal data to ASCII for display on the terminal and printer devices.

OVRLAY. OVRLAY is the routine which actually performs the overlay function. After an overlay is loaded into memory, OVRLAY jumps program control to the overlay program.

PDFPRT. PDFPRT prints the amplitude distribution to the terminal device (printer). Although resident in EKG-EXEC, PDFPRT is called only by the DISPLAY module.

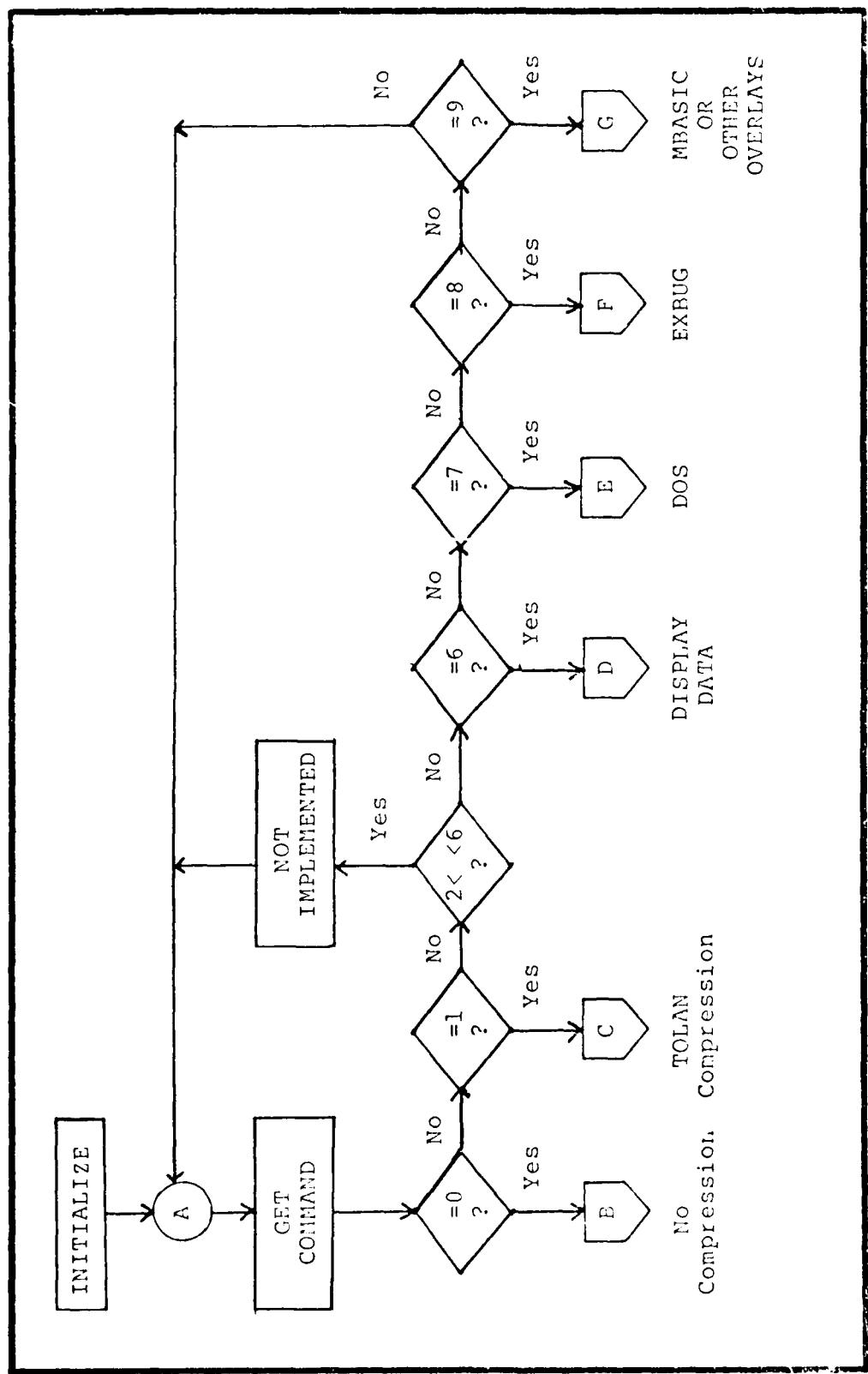


Fig 22. EKG-EXEC Functional Flowchart

NOCPRS. NOCPRS is a data acquisition module in which the EKG waveform is sampled and stored without data compression. The data is rounded to 8 bits from 12 bits. The reason for this rounding is explained in chapter 5.

This module was constructed for two reasons : 1) uncompressed data was considered useful for doing experimental studies on potential data compression techniques implemented after the original data collection session and ; 2) this module was the structure around which the Tolan (and potentially other) compression routines were built. The basic operation of NOCPRS is illustrated in Figure 23.

As is seen in Figure 23, NOCPRS does more than just sample the EKG. Statistical parameters are collected and updated during an EKG data collection. These parameters are used to measure the real time performance of the compression (no compression) software. A detailed description of these measurement parameters is described in chapter 5.

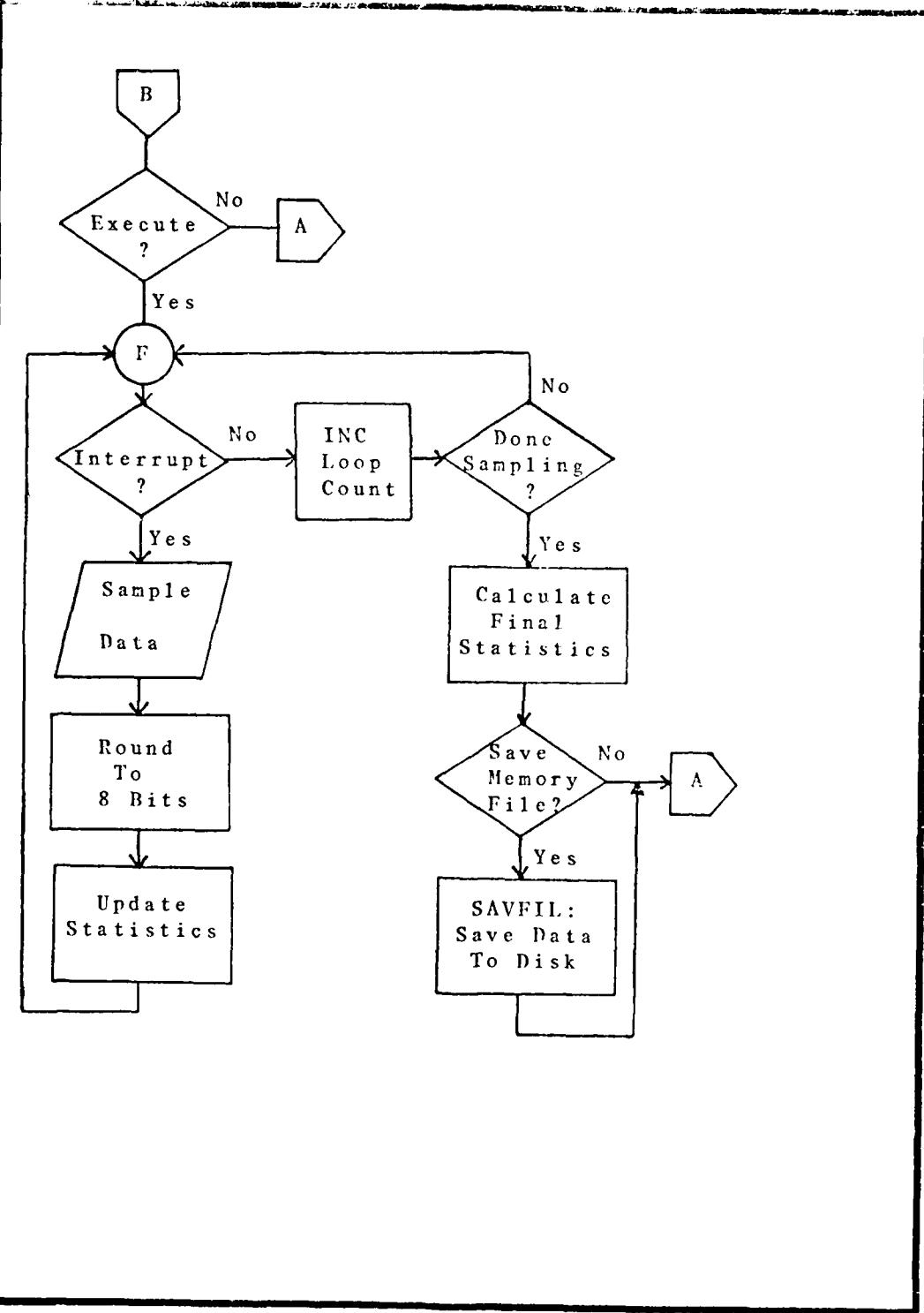


Fig 23. NOCPRS Functional Flowchart.

DISPLAY. DISPLAY is a command module similar in construction to EKG-EXEC. DISPLAY controls the data output modules which display the sampled EKG data and data statistics to two output devices. The first device is the terminal (printer) and the second device is an oscilloscope. The basic command structure of DISPLAY is shown in Figure 24. DISPLAY is broken into 5 working submodules. These modules are described as follows:

PRSTAT. This submodule reads the memory file header (3C00-3D00 Hex), formats the statistical data found there, and prints this data to the terminal (printer). The statistical data in appendix D was generated by PRSTAT.

PDFPRT. This submodule prints the memory file histogram tables to the terminal (printer). PDFPRT output is also listed in appendix D.

PDFOUT. PDFOUT scans a user selected lead histogram (X,Y,Z) and formats the data for display to an oscilloscope. The data is output via D/A ch 0. An example of PDFOUT output is found in Figure 31 in chapter 5.

DECPRS. The DECPRS module scans the memory file header and identifies the compression technique which was used to encode the data in memory. The appropriate decompression algorithm is then selected and the data decoded and output on D/A channel 0.

LOAD. LOAD initiates a data file load from disk memory to RAM. This load is performed by FOS routines called by load.

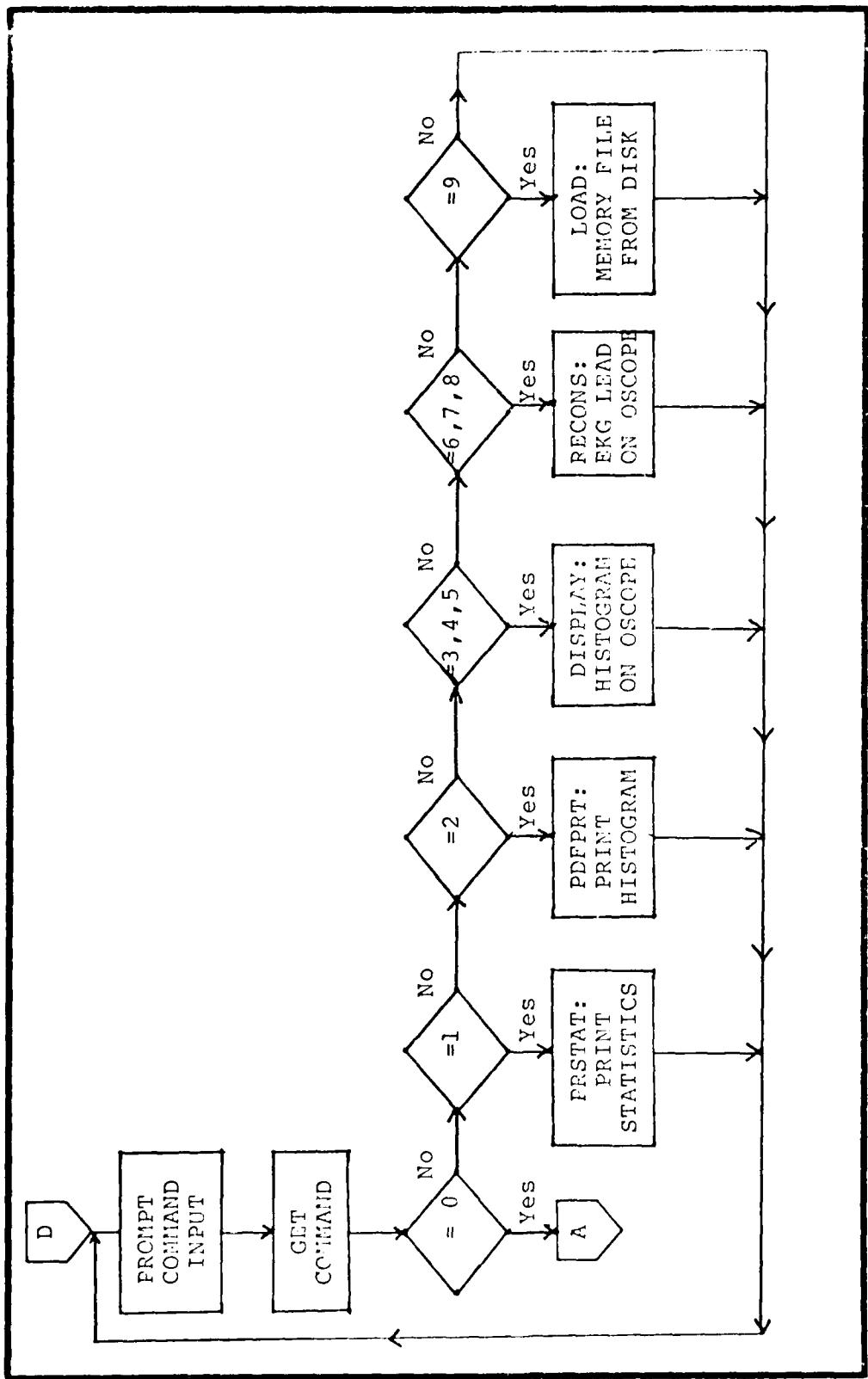


Fig 24. DISPLAY Functional Flowchart.

TOLAN. The TOLAN module is the "heart" of the EKG-DAAS and performs the actual data compression on the sampled EKG. As was described in chapter 3, the TOLAN algorithm first performs a second difference time compression operation followed by a variable length encoder. The operation of the TOLAN compression module is shown in Figure 25.

To detect sample clock (or CPU clock) drift, time calibration operations are performed prior to and after the data collection run. Other statistical data parameters are also collected allowing post collection measurement of compression performance.

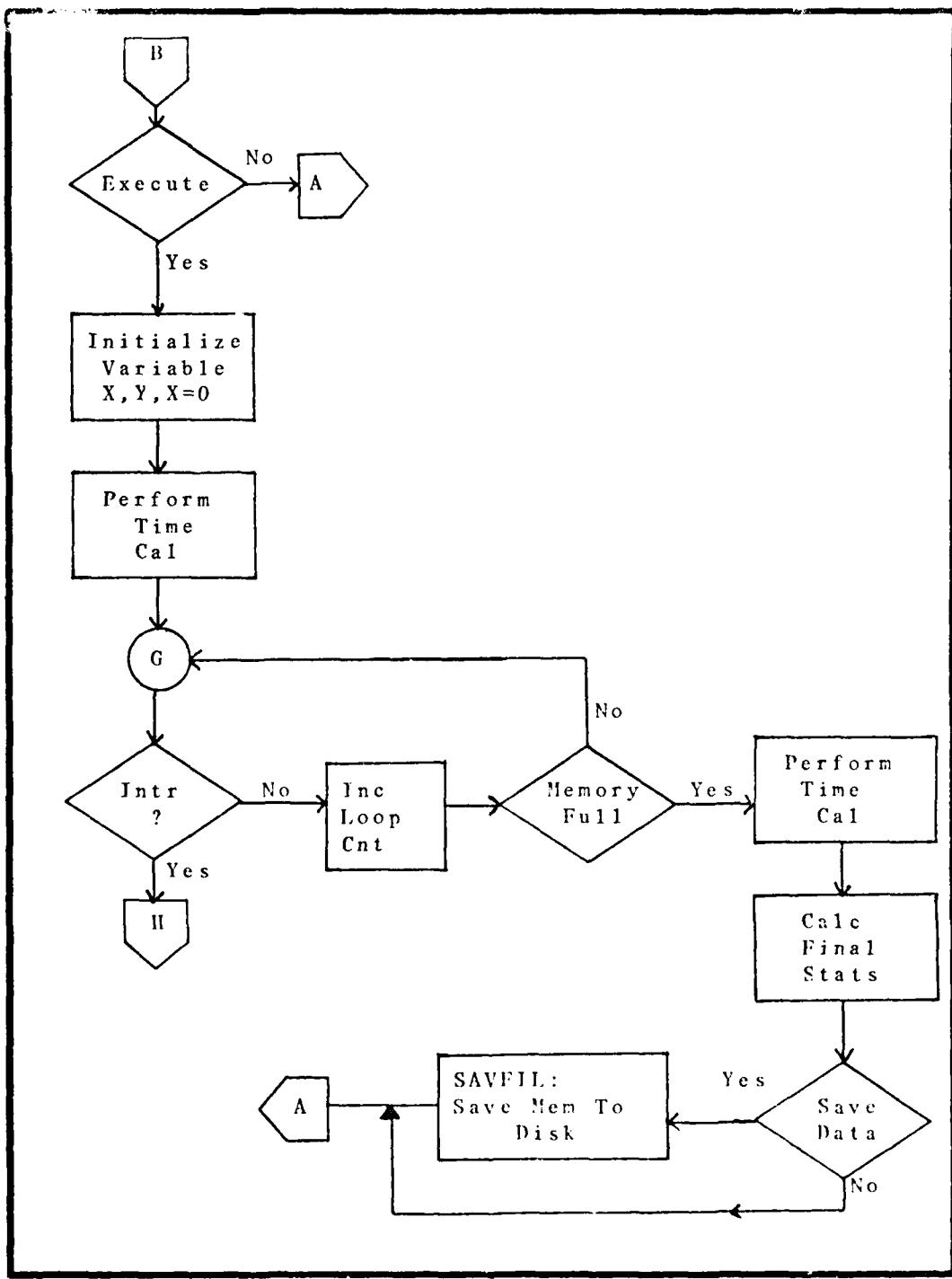


Fig 25a. TOLAN Functional Flowchart.

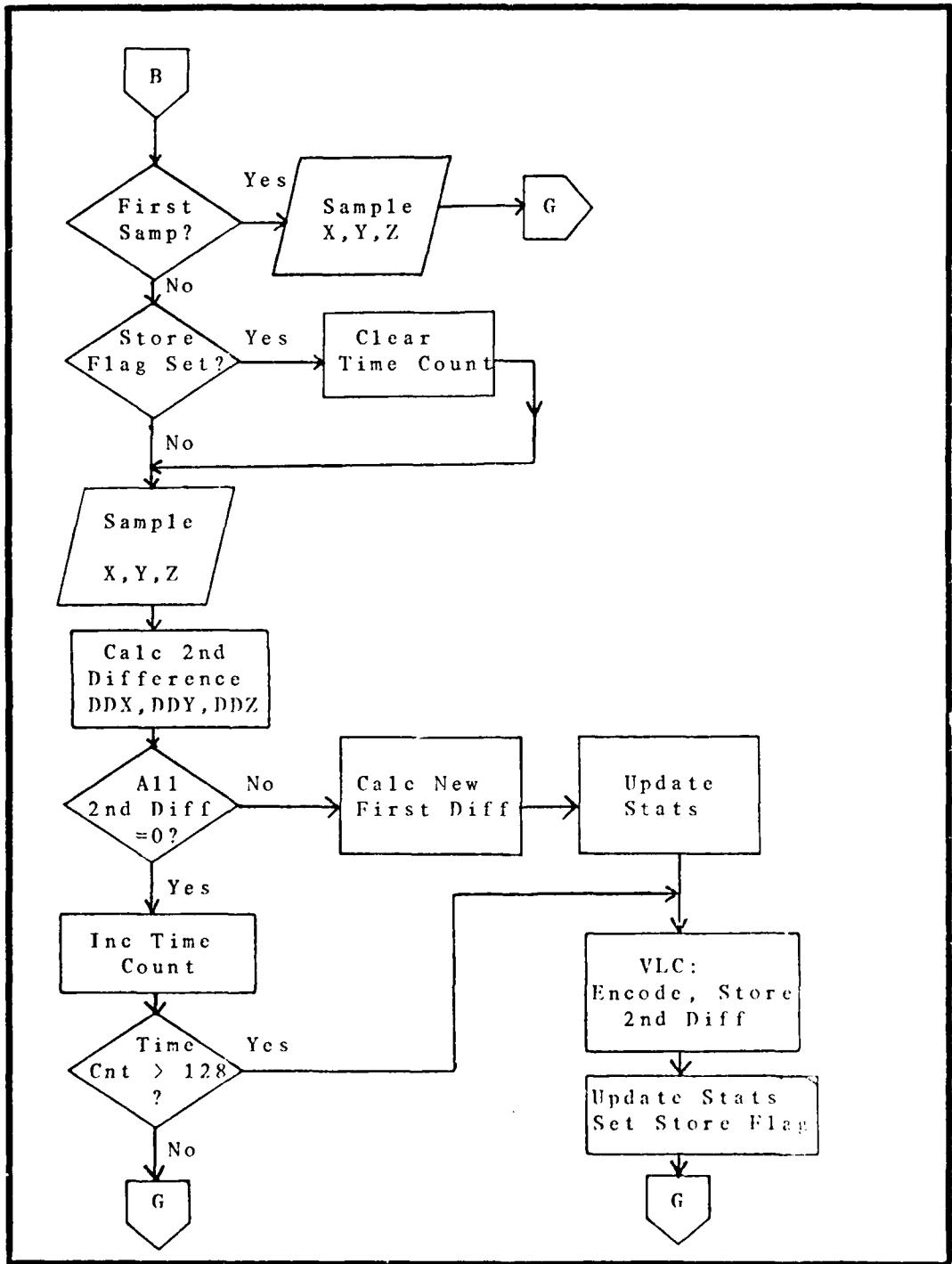


Fig 25b. TOLAN Functional Flowchart

### Summary

This chapter has presented an overview of the complex software which makes up the EKG-DAAS. The reader is referred to the program listings in appendix C for a more thorough description of the program operation.

The software structure in the EKG-DAAS was written in a "Top Down" manner and all attempts have been made to document the operation of each routine. Since assembly language is difficult to read, this chapter was written to assist the reader in understanding the basic structure of the EKG-DAAS. The next chapter presents the results of the EKG experiment where data was collected, analyzed, and compressed by the EKG-DAAS.

## V. Experimental Procedure, Data Analysis, and Results

This chapter presents the results of an EKG collection experiment where "in vivo" EKG data was taken from test subjects in real time. Data was taken and stored in both compressed and uncompressed formats for later analysis and reconstruction.

Chapter 5 is organized as follows. First the "experimental" setup is described along with a description of the EKG equipment, collection environment, and subject personnel. Next the parameters used to determine compression performance are defined followed by the analysis and results of the experimental data. Finally, the chapter concludes with a comparison between the results obtained using the Tolan compression algorithm and the estimated performance of the Dower technique. Discussion now turns to the experimental procedure.

### Experimental Procedure

The data was taken from a set of nine fellow students during a laboratory course on electrocardiograms. The equipment used to produce the EKG was the model DR-12 research recorder built by Electronics for Medicine, Inc. (see appendix E). The DR-12 is a vintage medical recording system built in the late 1950's and is constructed with vacuum tube amplification circuitry.

The personnel used for test subject were all Air Force

officers in good physical health. No test subjects with obvious heart disease were used although significant EKG variations between subjects was noted (Fig 27-29). Output of the DR-12 was limited to one EKG signal which could be switched to any of the 6 "limb leads" (Ref 13:29-34) by controls on the DR-12. The electrodes of the EKG were applied to the wrists of the test subjects and in some cases, not all, an electrode jelly was applied to reduce skin-electrode resistance.

Since the EKG-Data Acquisition and Analysis System (EKG-DAAS) was configured for a 3 lead system, the X,Y,Z inputs were connected in common and the single signal available from the DR-11 applied to this connection. A Brush Instruments Mark II recorder and a Tektronics Model 465M Oscilloscope were used as analog output devices (Appendix E). The display instruments were connected in common with the A/D inputs as is shown in Figure 26.

Prior to the data recording session, the A/D was calibrated in accordance with the operating manual (Ref 8). The A/D was configured for a dynamic range of -5.000 volts to +4.9976 volts with 2's complement binary representation.

The actual data collection proceeded as follows. First the subject was connected to the DR-12 and the amplitude of the resulting EKG signal adjusted to fall within the A/D dynamic range. A test run of all six limb leads (I, II, III, AVL, AVR, AVF) was then taken (without storage by

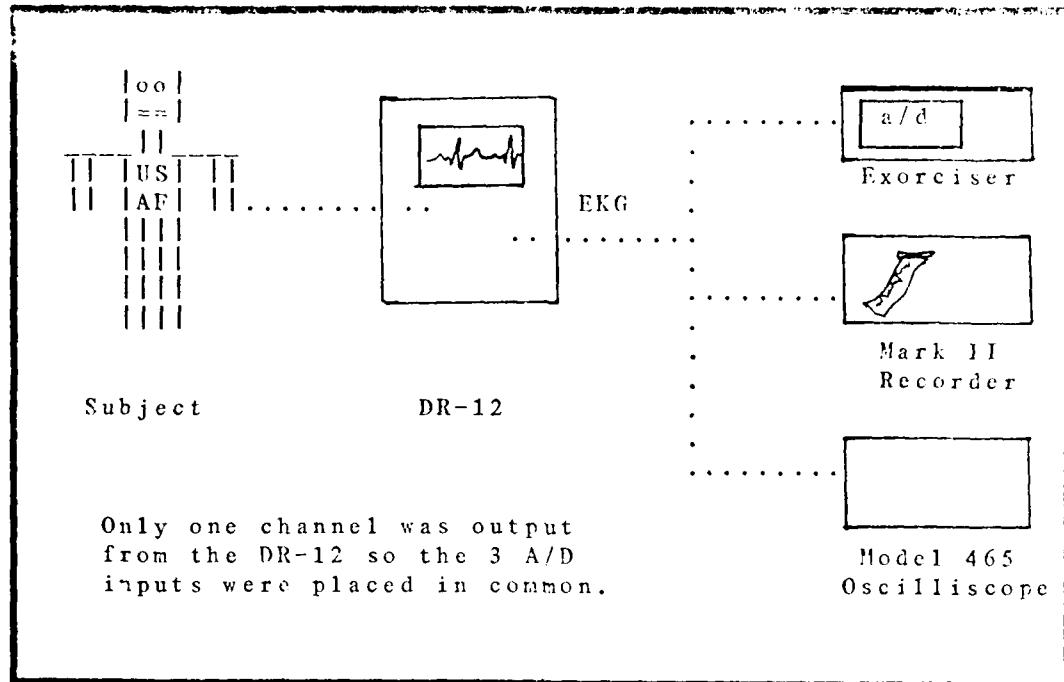


Fig 26. Experimental Data Collection Setup.

EKG-DAAS) and the lead with the least "noisy" signal was selected. EKG-DAAS was then executed and a data collection made without compression. This uncompressed data collection was followed as soon as possible with another "run" in which the Tolan compression algorithm was enabled. Following both data collections (uncompressed and compressed), the raw data traces from the Brush recorder were annotated with the time and subject and filed for later data comparison with the reconstructed waveform.

The data was stored on "floppy" diskettes and processed post collection for the entropy and maximum compression statistics. Before the results of these data

collections are presented, the compression measurement parameters will be defined.

#### EKG Compression Measurement Parameters

To permit determination of compression performance, a set of statistical parameters was calculated and saved during each data collection. These statistical parameters are described as follows:

Number of Samples. This statistic was saved to determine the total number of bits that were input to the compression stages. The total number of bits were calculated by  $(8 \text{ bits/sample}) * (3 \text{ leads}) * (\text{num of samples/lead})$ .

2nd Difference Frequency of Occurrence. Four frequencies of occurrence tables were kept with double precision binary counters. Following the data collections, these  $\Delta^2 x, \Delta^2 y, \Delta^2 z, \Delta t$  histogram tables were input to a BASIC program (ENTROPY) where the entropy of second difference "source" was calculated.

Total Waiting Loop Counts During Collection. A counting loop was established in the TOLAN compression module which allowed determination of the percent of the sampling period used for the compression and statistics calculations. One circuit of this counting loop takes 46 machine cycles of the 6800 microprocessor. A count of the total number of loop cycles completed following the sampling/compression interrupts is kept in the collection statistics buffer.

Maximum Loop Count Per Interrupt. To offset the inaccuracy which would develop if the interrupt clock period changed between runs (or if the master clock in the Exerciser drifted), a loop count calibration was performed immediately before and after each collection run. This calibration was accomplished by performing 256 sequential interrupts with no interrupt processing except return-from-interrupt. The before and after calibration counts were then averaged and the maximum loop counts per interrupt calculated. Time Efficiency of the compression operation was

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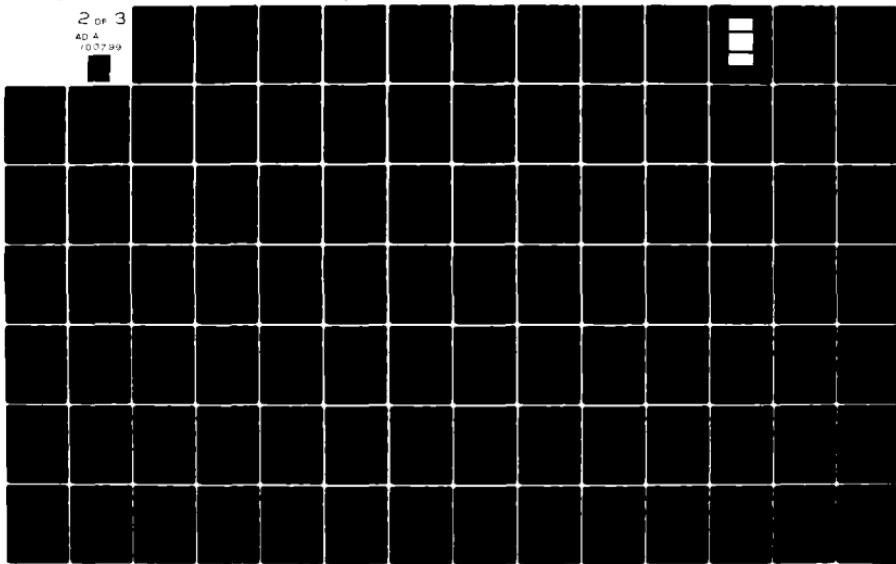
AIR FORCE INST OF TECH WRIGHT-PATTERSON AFB OH SCHO0—ETC F/6 6/5  
ANALYSIS AND PERFORMANCE EVALUATION OF ELECTROCARDIOGRAM DATA C—ETC(U)  
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then calculated from the equation:

$$T.E. = (1 - (\text{Total Loop Count} / ((\text{Num of Samples}) * (\text{Max Cnt})))) * 100$$

Channel Maximums and Minimums. The channel maximums and minimums were retained to allow determination if the analog inputs exceeded the A/D dynamic range.

Number of Memory Bits Available. This number was constant and was determined by the amount of Read/Write (RAM) memory available for data storage. For the current configuration of the Exorciser and the EKG-EXEC program this was 139248 bits (17406 bytes).

Number of Bits Available to Variable Length Coder. This counter measured the number of bits out of the data decorrelator and allowed calculation of the decorrelator's compression ratio (bits out/bits in).

Number of Bits Used to Store Channel X,Y,Z. These counters measured the number of bits used to store the data from the three input leads. This count is the number of code bits out of the variable length encoder.

Number of Bits Used to Store Time. This counter was identical to the channel counters above but measured the number of code bits used to store the  $\Delta t$  run counts.

Total Compression Ratio Achieved. This figure was calculated post collection by dividing the total code bits stored by the total R/W memory bits available.

The statistical data defined above was compiled by the EKG-EXEC program and is printed by the DISPLAY software module as illustrated in appendix D.

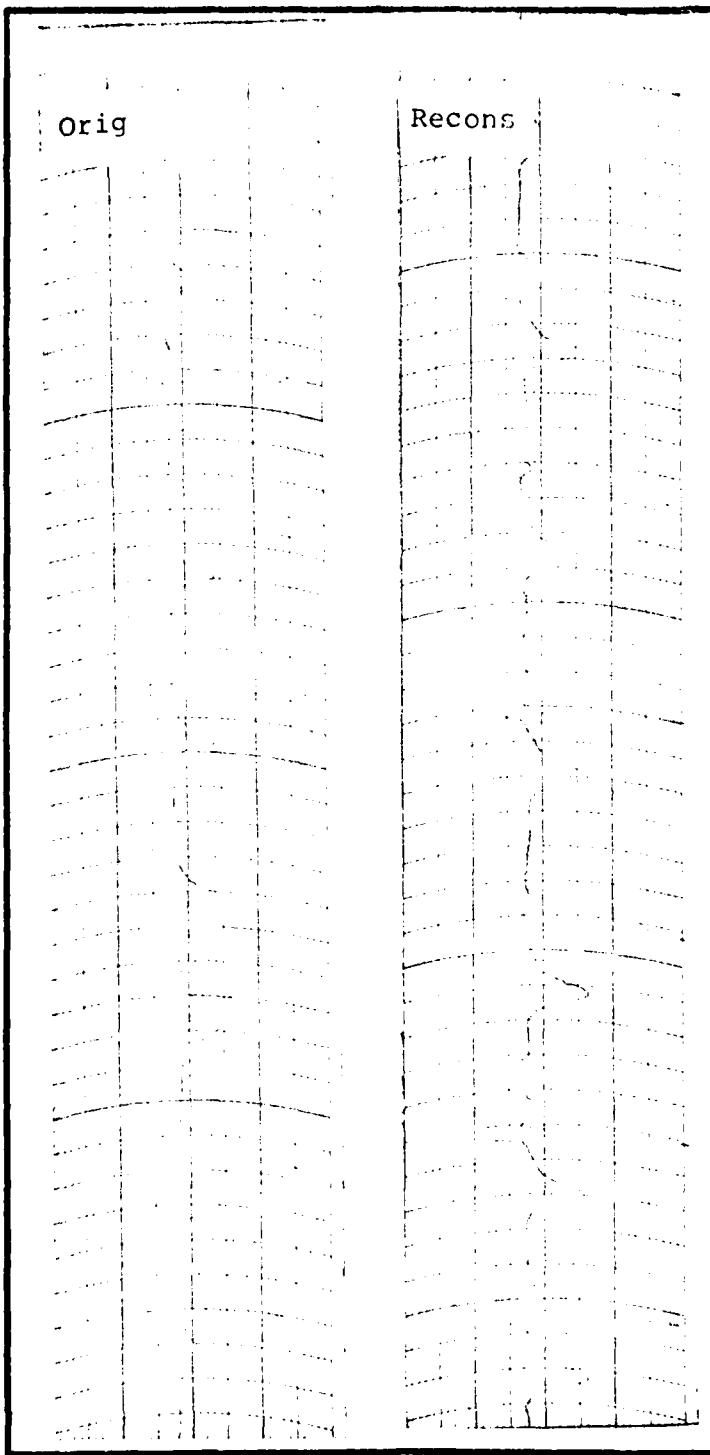


Fig 27. Original and Reconstructed EKG with Best Compression Ratio (2.260 : 1)

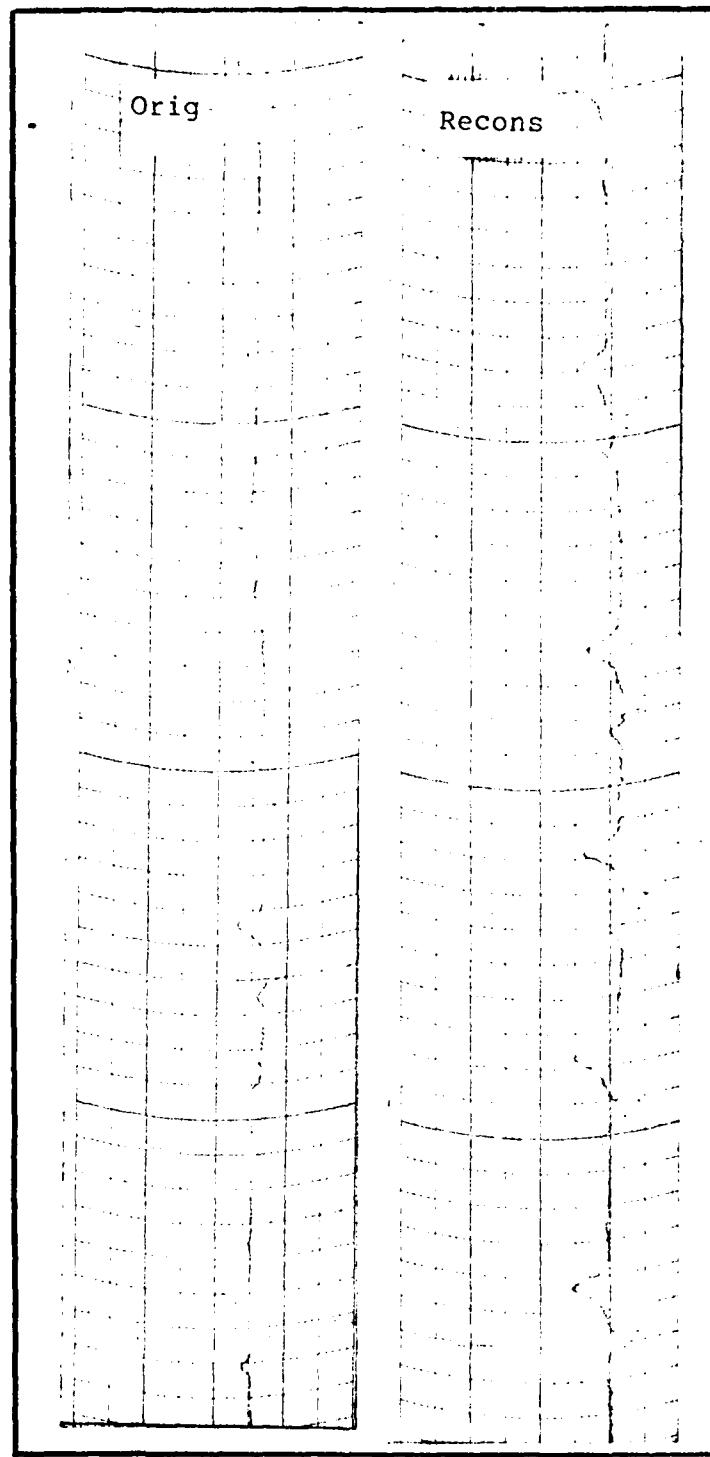
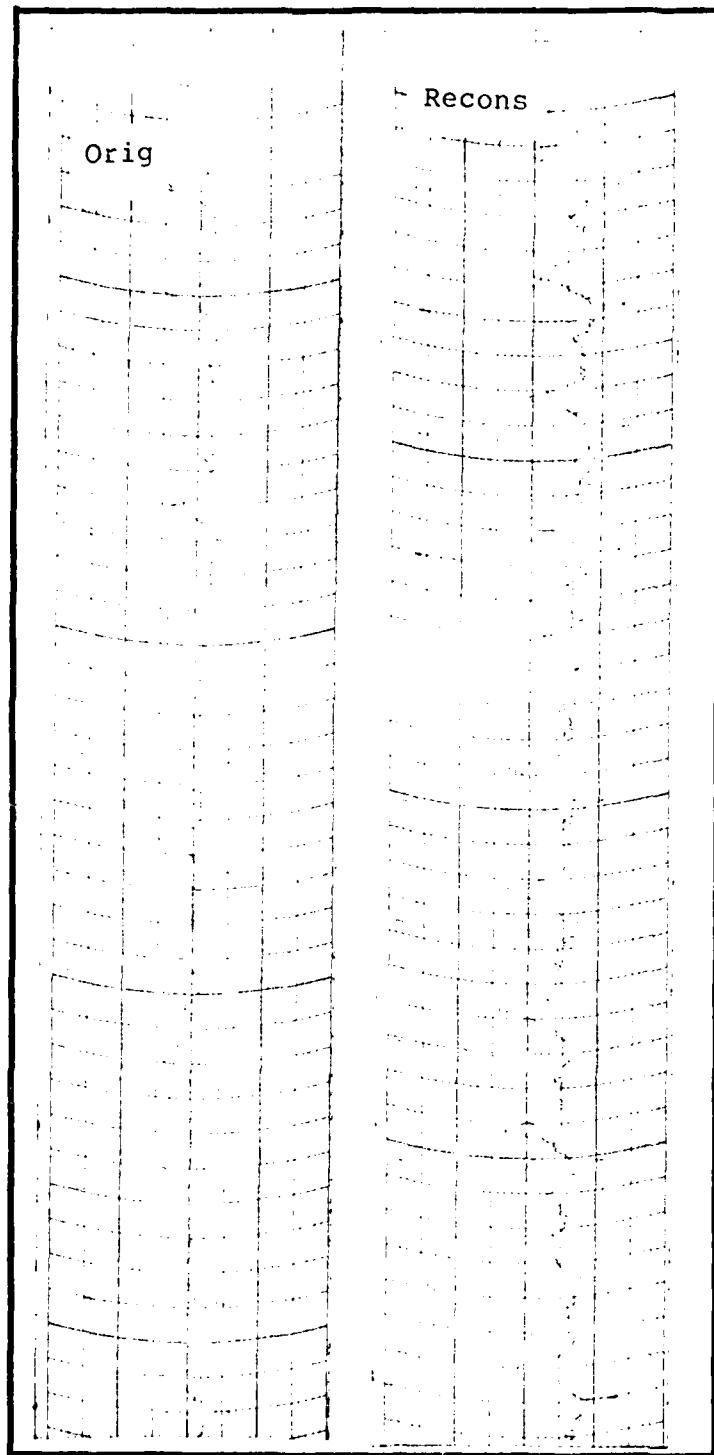


Fig 28. Original and Reconstructed EKG with Average Compression Ratio (1.60 : 1)



**Fig 29 . Original and Reconstructed EKG with Worst Compression Ratio (1.252 : 1)**

## Data Analysis and Results

To determine the compression efficiency of a redundancy reduction EKG compression algorithm is, in general, a difficult job to accomplish. The reasons for this difficulty are: 1) the sampled EKG data sequence contains significant correlation (memory) between sample values making calculation of the absolute bound of the entropy extremely difficult (Ref 34:479-489) and ; 2) the techniques used for both the data decorrelator and entropy encoder vary significantly from algorithm to algorithm.

In this thesis, as is done in the literature (Ref 7,12,28), the assumption is made that the output of the decorrelator is ''almost decorrelated''. Decorrelated is not ''independent'' (unless the source was statistically gaussian), but it is assumed that true entropy of the 2nd difference sequence approaches the value which would be calculated by Eq.(1) (reproduced below).

$$\Delta^2 \text{ entropy} \approx - \sum -p_i \log_2 p_i \quad (27)$$

( $p_i$  is the probability of the i'th second difference).

This second difference entropy can then be used as an upper bound on the potential entropy encoding compression of the Tolan redundancy reduction technique.

The  $\Delta^2$  entropy values tabulated in Table II were calculated by Eq.(27). To calculate the approximate limit on the entropy encoder's compression ratio, the uncompressed data word length of 8 bits was divided by the ''lowest''

average code word (i.e., the 2nd difference entropy). A look at Table II shows that the entropy encoder compression ratio varied between 58% and 71% of this entropy bound. Since the Tolan variable length code is suboptimal, a lower efficiency is expected. Nevertheless, the Tolan entropy encoder performed more-or-less consistently across the data set. This last observation would imply that the data decorrelator influences the overall compression ratio more

Table II  
Experimental Data Summary

Subject Id	2'nd Difference Entropy	Maximum Compression Possible (Approx)	Achieved 2nd Difference Compression	Percent of Max Encoder Comprs	Percent of Sample Interval	Achieved Total Comprs
TA1545T	3.2801	2.42 : 1	1.55 : 1	64.0%	93.9%	1.36 : 1
TA1548T	3.0171	2.63 : 1	1.61 : 1	61.2%	91.4%	1.50 : 1
TA1559B	3.016	2.64 : 1	1.72 : 1	65.1%	93.0%	1.53 : 1
TA1511S	3.323	2.39 : 1	1.65 : 1	69.0%	94.0%	1.43 : 1
TA1520B	2.601	3.06 : 1	1.81 : 1	59.1%	90.9%	1.72 : 1
TA1448L	2.930	2.72 : 1	1.73 : 1	63.6%	90.6%	1.60 : 1
TA1439S	3.267	2.43 : 1	1.63 : 1	67.1%	93.4%	1.43 : 1
TA1359P	2.487	3.22 : 1	1.88 : 1	58.3%	77.8%	2.26 : 1
TA1413L	3.783	2.10 : 1	1.50 : 1	71.4%	95.4%	1.25 : 1

than the entropy encoder.

As was described in chapter 3, the Tolan data decorrelator uses time compression in conjunction with a 2nd

difference operation. The Tolan time compression technique makes the assumption that the second difference value of zero occurs so frequently that the encoding of 0 would be less efficient than the storing of a zero value run counter. In the experimental situation in this thesis, signal noise was quite evident in the EKG traces (Fig 29). The sharp "spikes" induced by noise are accentuated by the second difference operation, hence a  $\Delta^2 \neq 0$  was a common occurrence. This forced the storage of a lot of 7 bit time counters.

In 8 out of the 9 compression runs made, the frequent storage of time counts actually caused the Tolan data decorrelator to produce negative compression (i.e. more bits out than went in). Since the entropy encoder was producing a larger positive compression ratio, the overall compression figure remained positive. The effect of this operation is graphically illustrated in Figure 30.

A look at the original and reconstructed EKG traces (Fig 27-29) in conjunction with the data in Table II, shows that as the "noise" level increased on the signal the compression became progressively worse. Since it was concluded that the entropy encoder performed approximately the same across the data set, the degradation in total compression must be due, in large degree, to the degradation in the Tolan time compression data decorrelator performance.

A clear effect of the noise is illustrated in Figure

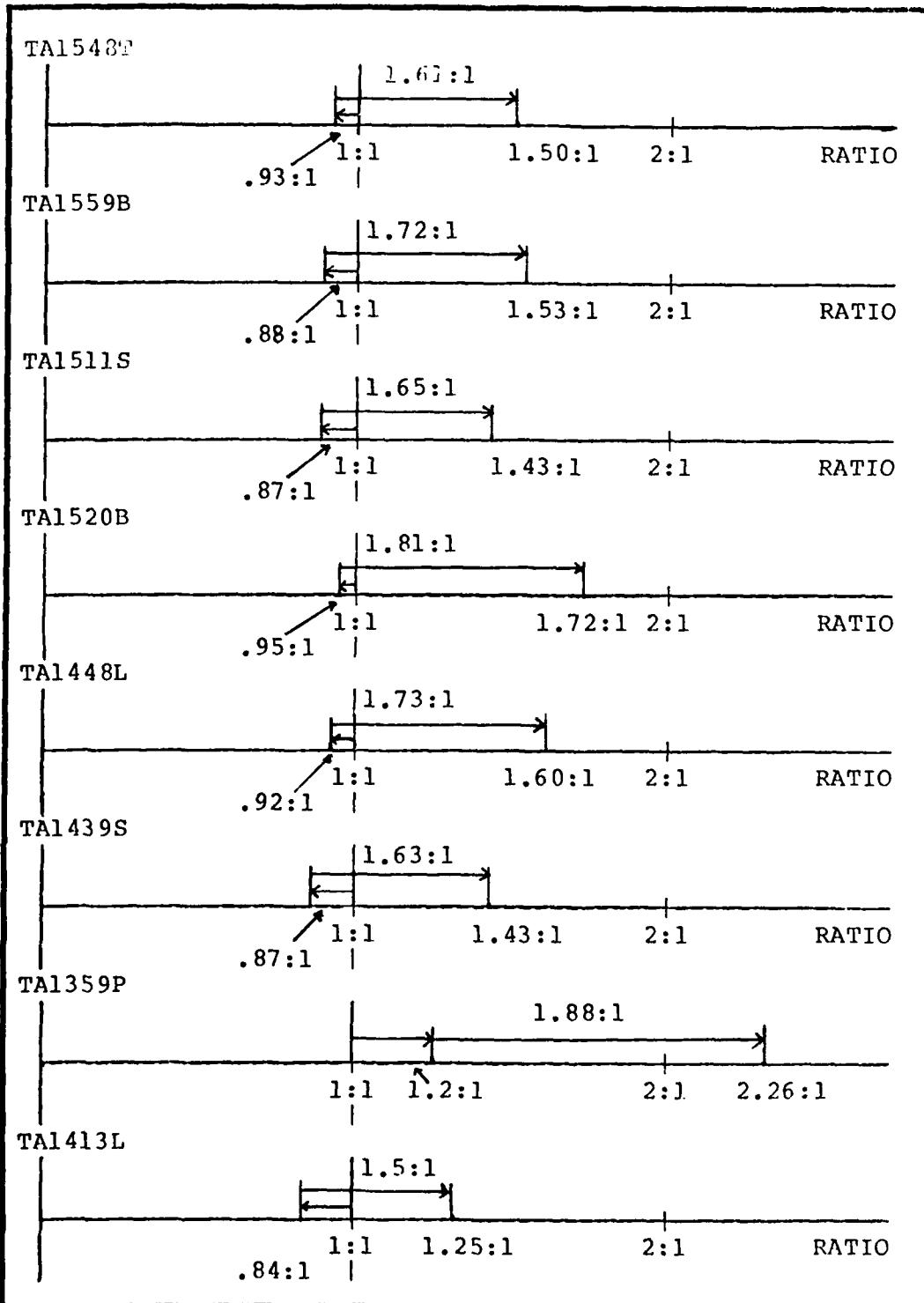


Fig 30 . Tolan Compression Ratio Breakdown.

\* 31. From Figure 31, it can be seen that as the "noise" level increased the variance of the 2nd difference distribution became increasingly larger. As the distribution became less peaked, the efficiency of the variable length encoder decreased with a subsequent loss in achieved compression ratio.

The nonzero component at  $\Delta^2 = 0$  in Figure 31 can be explained as follows. The decision was made early in the design of the EKG-DAAS that 8 bit data would be used versus 9 or 10 as recommended by the American Heart Association (Ref 3). This decision was made to simplify the software (i.e. single precision could be used). Since the A/D converter has 12 bit resolution, the sample was rounded to 8 bits for uniform error distribution (Ref 24:424-432). Unfortunately because of this rounding action, small differences in the least significant bits of the A/D converter may have affected the rounding operation. Since the A/D cannot sample all three channel simultaneously, the probability that a "noisy" signal will change the least significant bit (or bits) is high. With a rounding operation, these changes may ripple to affect the least significant bit of the 8 bit data.

At most this effect would only cause a change in the least significant bit of the 8 bit rounded values. This would cause, however, the second differences between of channel X,Y,Z to be different even though they were connected in common. As was described in chapter 3, any of

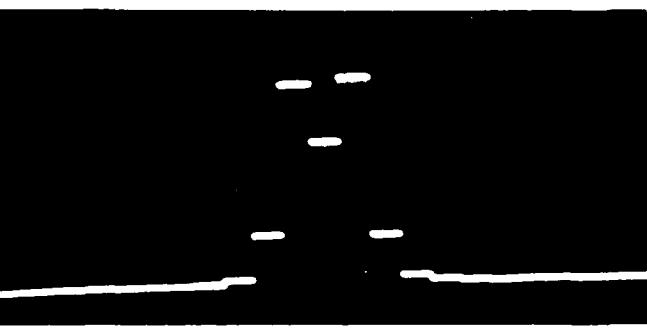
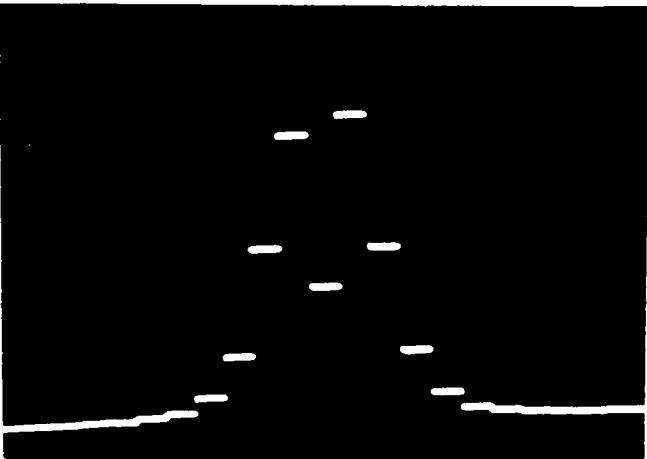
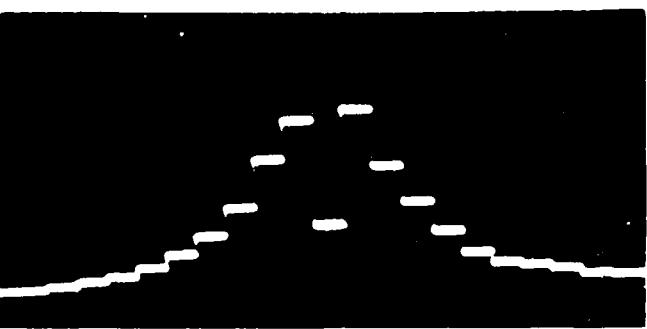


Figure 1. Three photographs of the same area of a concrete wall showing the effect of different methods of applying a thin film of polyurethane sealant.

(a) A thin film applied by a brush.

(b) A thin film applied by a roller.

(c) A thin film applied by a spray gun.

the three leads with a non zero second difference forced the storage of all three data points. Hence many zeros data points were stored because of 2nd difference asymmetry which would not have occurred had the 12 bit A/D data been truncated instead of rounded.

This fact undoubtably affected the overall compression efficiency of the Tolan algorithm. Nonetheless, the distribution in Figure 31 would not have changed (except the zero value count) significantly and noise would still have broadened the 2nd difference distribution.

#### Tolan and Dower Performance Comparison

As has been reiterated several times in the text of this thesis, the author's original intention was to implement both the Dower and the Tolan algorithms for experimental test and comparison. Since time did not permit the Dower implementation, an experimental comparison was not possible.

As an attempt to compare the results of the Dower and Tolan compression routines, the results quoted from the papers by Dower, Berghofer, and Stewart (Ref 12,29) will be used. Dower states that his state space variable length encoder approaches the entropy bound of the second difference source "with about 1.65% wastage" (Ref 12:3). This value is significantly higher than the approximately 30% "wastage" observed with the Tolan variable length encoder.

Assuming similar operations of the second difference time compression (in this author's "noisy" environment), the Dower decorrelator is expected to have negative compression. With the Tolan and Dower decorrelators assumed "equal", then the real gain of the Dower technique over the Tolan approach is in the VLC.

Assuming the Dower entropy encoder approached closely to the entropy bound, it can be extrapolated that the Dower algorithm would have achieved a maximum compression of approximately  $(1.2)*(3.22) = 3.86:1$  for the "best" EKG in Table II and Figure 30. The worst compression ratio would have been  $(.84)*(2.1) = 1.76:1$  for the worst (noisest) EKG.

### Chapter 5 Summary

This chapter began with a description of the EKG collection experiment. Although the EKG apparatus was limited to one channel, successful collection and compression of EKG data was performed. Analysis of the data revealed that time compression is inefficient in a "noisy" environment and that "rounding" of the 12 bit A/D samples in conjunction with placing all three sample lead in common accentuated the degradation caused by the time compression data decorrelator. Nonetheless, the Tolan algorithm did achieve an overall positive data compression figure but significantly lower than the 9:1 ratio achieved by Ruttiman and Pipberger (Ref 28) or the average value of 7.3:1 reported by Stewart, Berghofer, and Dower (Ref 29). Finally

it was heuristically shown that if the Dower compression algorithm lived up to the statements by Dower, then a compression ratio gain of 3.86:1 to 2.26:1 could have been achieved with the Dower EKG data compression technique. This thesis will now proceed to provide conclusions and recommendations

## VI. Summary, Conclusions, and Recommendations

### Summary and Conclusions

This thesis has investigated the field of electrocardiogram data compression with the objective of evaluating compression algorithms on a 6800 microprocessor based computer system. Accomplishment of this goal required the construction of the EKG-Data Acquisition and Analysis System utilizing the Motorola Exorciser microcomputer.

To determine those EKG compression algorithms which had potential for Exorciser implementation, a literature search was made to locate EKG data compression techniques. In addition to the literature search, personal correspondence (Ref 11,31) yielded several EKG compression algorithms. The results of this research is presented in chapter 2.

Since thesis requirements dictated the need for an online, real time data compression algorithm, only the fastest EKG compressors could be considered. Two routines were selected for detailed analysis and implementation. These two compression algorithms (Tolan and Dower) were discussed at length in chapter 3.

It was deduced from the methodology of the two data compression techniques that the Dower algorithm would perform better than the Tolan procedure. To test this

hypothesis, the EKG-DAAS was constructed and the Tolan algorithm implemented. The implementation of the EKG-DAAS is documented in chapter 4.

Time constraints prohibited completion of the Dower compression algorithm, but data was successfully compressed, analyzed, and decompressed with the Tolan algorithm. The results of this analysis are presented in chapter 5.

The conclusions of this research effort are as follows. First, EKG data compression can be accomplished in real time by a microprocessor based computer system. The Exorciser is a slow microcomputer (1 MHz cycle time) yet it was still possible to implement the Tolan algorithm with a 500 Hz sample rate. The second conclusion is that signal noise can dramatically affect the efficiency of the EKG routines in the same class as the Tolan algorithm. The expense and implementation difficulties of prefiltering, low electromagnetic noise environment, and proper EKG lead attachment are well worth the gain in compression achieved. Finally, a software project of this magnitude should not be attempted totally in assembly language. Although assembly language offers the greatest flexibility and speed, algorithm implementation and debugging efforts are enormous. A high order language would have allowed this author to complete his original thesis objectives.

### Recommendations

Several recommendations are offered for further study in microcomputer based EKG data compression. First, implementation of the EKG compression algorithms using a 16 bit microprocessor (e.g. 6 MHz Intel 8086) would permit an order of magnitude improvement in speed of execution. Hardware multiply and divide along with 16 bit arithmetic registers would permit easy implementation of the Dower algorithm and would even make use of the Transform compressors (i.e. FFT) feasible. Second, further study is needed on determination of decorrelator inefficiency on the overall data compression. A large study of different decorrelators such as 1, 2, 3 difference operations with and without time compression is needed. Next, an EKG compression algorithm implemented using the Ruttiman and Pipberger 2nd order interpolator (Ref 28) for the decorrelator along with the Dower variable length encoder (Ref 12) should be built. This combination should prove to be very efficient. Finally, programming and experimental testing should be done on a full scale microcomputer development system, complete with a high order language, A/D-D/A capability, and flexible disk file manipulation software. Such a system is the Zilog MCZ1/25 microcomputer resident here at A.F.I.T.

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## Appendix A

### The Electrocardiogram

#### Introduction

The electrocardiogram (or EKG) is a record of the electrical activity of the heart as measured from the body surface. The magnitude, shape, and timing of the electrical potentials generated by the heart reveal a great deal of information concerning the health of the cardiac system. This appendix will describe how the electrical signals from the heart are generated, how the cardiac cycle is coordinated and controlled, and finally how the EKG can be used as a diagnostic tool.

#### The Physiology and Electrical Characteristics of the Heart

The heart (Fig. A1) is an organ about the size of a fist with four main pumping chambers and a specialized electrical conduction system. At the top are two thin walled pumps called the atrium which prime the main pumps of the heart, the ventricles. The ventricles are separated by a thick wall of muscle tissue called the septum. Blood from the right ventricle goes to the lungs and blood from the left ventricle goes to the rest of the body.

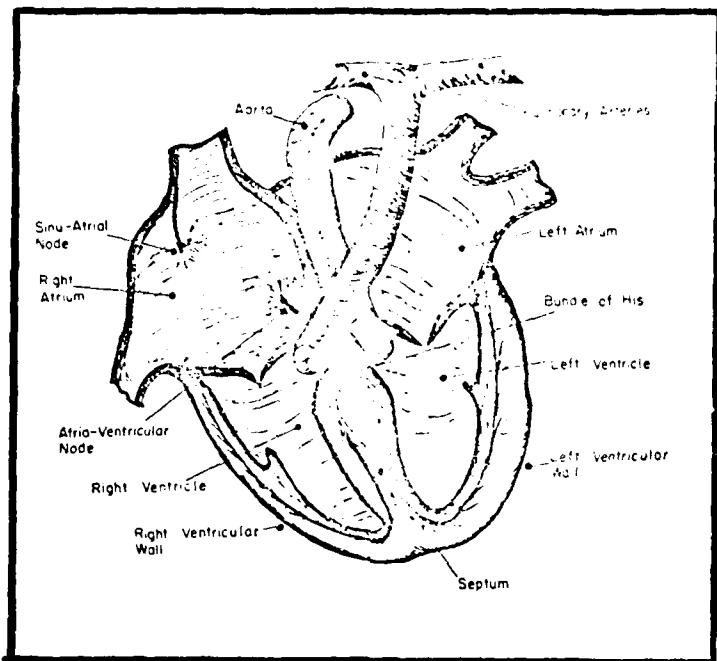


Fig A1. The Heart Cross-Section (From Ref 20:291).

During the resting state between contractions, the cells of the heart are "polarized" with an electrical potential existing between the inside (-) and the outside (+) of the cell. This potential is generated by an ionic gradient across the cellular membrane and is normally maintained for approximately .2 to .4 of a second before spontaneous "depolarization" occurs. Depolarization (caused by an inrush of sodium ions into the cell) induces the cell to contract for approximately 1/4 of a second. Because all heart cells contain specialized conducting fibers, the depolarization of one cell initiates the depolarization of neighboring cells and a "wave of excitation" sweeps across the myocardium (heart) at a rate

of about 1 meter per second.

The voltages measured at the body surface are the superposition of thousands of heart cells depolarizing (or polarizing) as the wave of excitement flows through the myocardium. Early work by Wilson (Ref 20:292) showed that the heart could be represented by an equivalent electrical dipole whose vector orientation sweeps through a closed loop during one cardiac cycle. In simple terms, as the wave of excitement flows toward a positive skin electrode, a positive slope is generated on the EKG record representing the projection of the heart vector onto the axis of the EKG lead.

To observe this sweeping dipole vector, Wilson developed the 12 lead EKG system in almost universal use today. This system (Fig. A2) attempts to measure the heart vector from a variety of vantage points in hope of determining the actual direction of propagation of the wave of excitement. Unfortunately, the leads of the Wilson system are not orthogonal and reconstruction of the actual heart vector orientation and amplitude is difficult. To overcome this problem, Frank (Ref 14:737-749) developed the vector cardiogram which combines 7 leads in a summing network to produce three orthogonal components of the heart vector. Frank VCG systems are popular in heart diagnosis and research because all of the information is contained in three leads of data versus 12.

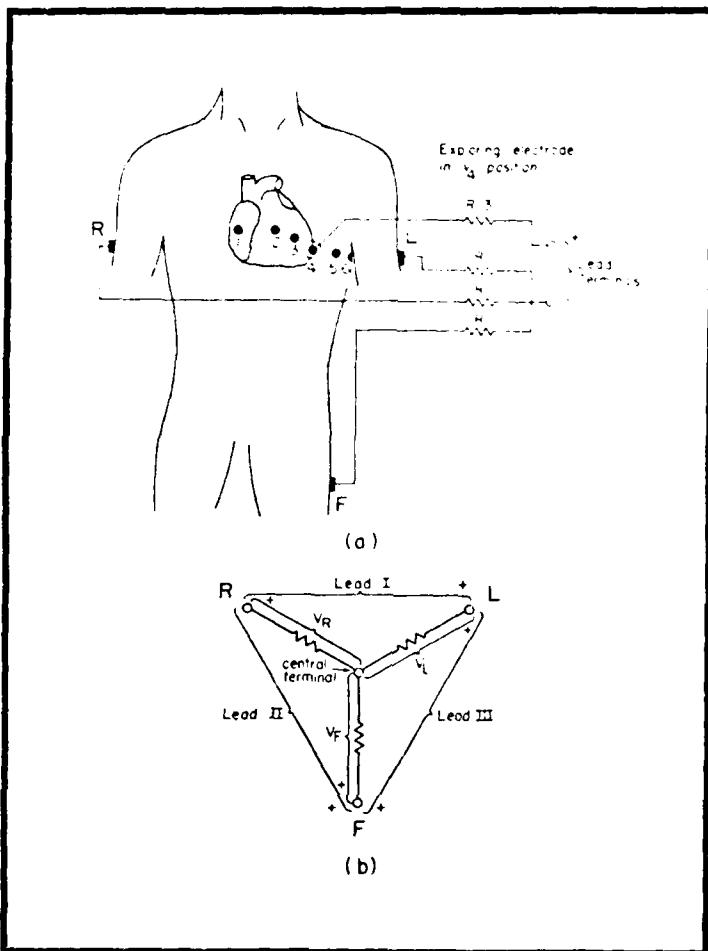


Fig A2. Wilson EKG Electrode System (From Ref 20)

### The Cardiac Cycle

For effective pumping action, the heart muscle must contract in a controlled, coordinated way. This means that the stimulating wave of excitation must follow a well defined conduction path to allow the heart to contract in the most efficient manner. In "normal" hearts, the contraction sequence begins in the right atrium where

special cells in an area known as the Sino-Atrium (SA) node spontaneously depolarize faster than the rest of the heart tissue. The SA node hence initiates a wave of excitation which covers the atrium in about 80 milliseconds. This atrial contraction produces the electrical signal called the P wave (Fig. A3) on a typical EKG record.

The wave of excitation started by the SA node then reaches another specialized receptor known as the Atrial-Ventricular (AV) node. Here connecting fibers delay the excitation impulse for about 50 milliseconds to allow the ventricles to fill with blood. After the 50 millisecond delay, the excitation signal is relayed to special conducting fibers in the septum known as the Bundle of His. These conducting fibers rapidly (30 ms) transmit the excitation wave to the interior (endocardium) wall of the ventricles where the wavefront propagates radially to the outer wall (epicardium) in another 30 milliseconds. The ventricular wave of excitation produces the QRS waveform complex seen on the EKG.

Following the ventricular contraction, the muscle cells of the ventricles repolarize over a period of 100 milliseconds. No muscular action is occurring but the EKG responds to this electrical activity and the T wave is noted on the EKG record. Atrial repolarization occurs during the QRS hence it is generally invisible on the EKG.

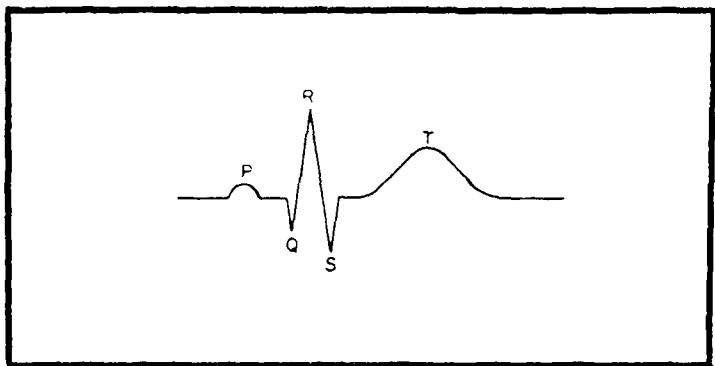


Fig A3. Typical EKG Waveform (From Ref 20)

Finally the heart rests for approximately .2 of a second and the cycle starts again. The above cardiac cycle is typical of a healthy heart. Disease, however, can effect this sequence dramatically.

#### Heart Disease and the EKG

The variety of ailments which plague the human heart are numerous and no attempt will made to describe the spectrum of diseases possible. There are, however, several common heart defects which routine EKG analysis usually detects. These include premature contractions (Atrial and Ventricle), bundle branch blocks, hypertrophy, and infarction.

Premature Contractions. Premature contractions are caused by the spontaneous depolarization of heart tissue outside of the SA node. This depolarization initiates a wave of excitement causing the atrium, or more notably, the

ventricles to contract out of rhythm with their normal cycle. Points where this "unscheduled" depolarization occurs are called ectopic foci and can arise because of infarct damage (to be discussed later), coronary heart disease producing oxygen starvation and a number of other causes. Premature Ventricular Contractions (PVC's) are highly visible on the EKG record. Normally the ventricles contract simultaneously and the voltage vectors generated tend to cancel keeping the QRS amplitude relatively small. A PVC, however, causes depolarization of one ventricle before the other generating an unbalanced, hence larger, voltage output.

Bundle Branch Blocks. A bundle branch block is caused by a block of the impulse of the right or left Bundle Branch. This causes a delay in the transmission of the stimulation impulse to ventricle blocked and forces the two ventricles to contract at slightly different times. This difference in ventricular contraction time shows up on the EKG as a double humped QRS.

Hypertrophy. Hypertrophy is an enlargement of one section of the heart muscle tissue. This enlargement affects the duration and strength of the wave of excitement and is visible on the EKG record as a diphasic trace if atrial hypertrophy is present. If ventricular hypertrophy exists, the QRS amplitudes are much larger than normal due to the fact that more tissue is depolarizing.

Infarction. Myocardial infarction is an injury to the heart tissue caused by an occlusion of a coronary artery. An area of the heart is then without a blood supply and often permanent damage occurs. This heart disease is often the one most commonly called a "heart attack" and, as is well known, is many times fatal. If a person survives the original "attack", then this permanent damage shows up in the EKG as a change in the QRS and T waves. This change occurs because the infarcted tissue no longer responds to the excitation wave and the wave moves around, not through, the damaged tissue.

Summary.

This appendix has briefly examined the physiology of the heart and discussed how the myocardial tissue generates the electrical fields measured by the electrocardiogram. Though EKG analysis has been practiced for fifty years, intense research continues in improving EKG diagnosis. Dramatic improvement in computer aided EKG analysis and better understanding of the electro-physiology of the heart is leading to improved cardiac health care worldwide.

The heart diseases discussed above are only a small subset of the problems which can afflict the human heart. Should the reader desire a more thorough background on heart physiology and cardiac disease, Dubin's book (Ref 13) is highly recommended. This programmed text carefully leads the reader through EKG analysis and is easily read. For a

more firm background on the electro-physiology of the heart,  
the article by McFee and Baule (Ref. 20) provides a good  
tutorial review on EKG history and research.

## Appendix B

### Fundamentals of Information Theory

#### Introduction

In 1948, Claude Shannon published a classic paper (Ref 30:379-423) titled "A Mathematical Theory of Communication" in which he laid the foundation of modern information theory. Shannon used his "information theory" to describe, mathematically, the interrelationships between the components of a "typical" communication system as illustrated in Figure B-1.

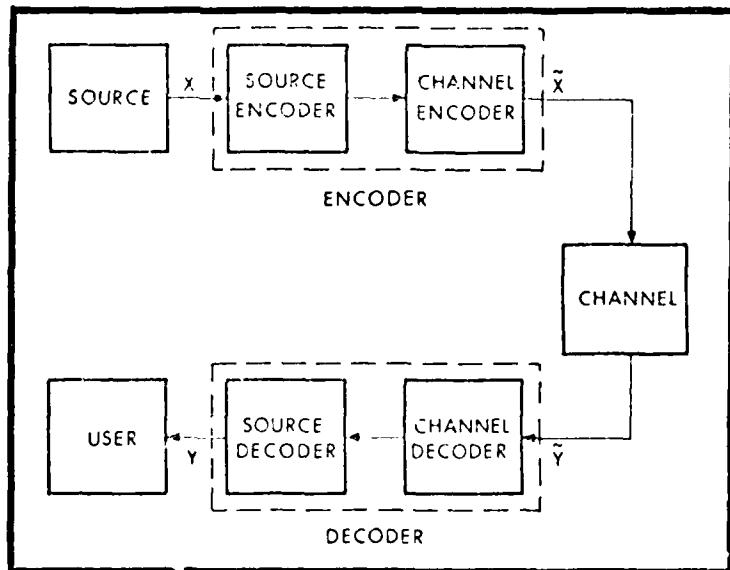


Fig B-1. The communication system model (From Ref 4).

This appendix is written as a basic tutorial on information theory, and is intended to acquaint the reader

with the terminology used in describing the data compression techniques in chapters 2 and 3. The appendix begins with a review of the information source, proceeds to a discussion of the transmission channel and Shannon's rate distortion theory, and concludes with a description of the system encoder/decoder.

### Information Source

For his initial analysis, Shannon proposed modelling the information source as a discrete stochastic process whose output is governed by known statistics. As pointed out by Davisson and Gray (Ref 10: 2-4), the discrete-time model was commonly used for any or all of the following reasons: (1) digital communication links have become commonplace; (2) a continuous time process can be modelled as discrete by sampling, orthogonal function expansion, or waveform segmentation; (3) greater simplicity.

The source is characterized by a finite set of possible outcomes known as its alphabet A. The occurrence of a particular alphabet symbol is governed by probabilistic descriptors (i.e. probability density functions) and it is assumed that the source produces only one symbol from the alphabet every  $T_s$  seconds. Hence the information source has a symbol rate of  $R_s = 1/T_s$  symbols per second.

The next question of interest is how much information is conveyed by the occurrence of a given source symbol? If X is a discrete random variable occurring at time t, and x is

an element of A, then the random variable self information can be defined. That is

$$I(x) = -\log p\{X=x\} \quad (B.1)$$

where  $p\{X=x\}$  is the probability that  $X=x$ .

According to this description, the less probable an event is, the more information is conveyed when it occurs. The base of the logarithm is unspecified, but in this thesis it is assumed to be base 2. Hence the occurrence of symbol x reveals  $I(x)$  bits of information.

The amount of information received per observation is of interest, but one would like a measure of the "uncertainty" or "randomness" of the source. If the stochastic process defining the source is considered stationary (a pretentious assumption but one generally made) then the output of the source is a sequence of random variables with identical probabilistic descriptors. The probabilistic descriptor will be defined as  $\{u\}$  and could represent the moments of the random variable or its probability density function (PDF). If, in addition, the source is considered ergodic, then statistical averages equal time averages and calculation of the set  $\{u\}$  is greatly simplified.

Given that the discrete process is ergodic, or at least stationary, then the source output at any time is described by the random variable X with range  $A=\{x(1), x(2), \dots, x(n)\}$ . The measure of the "uncertainty" or "randomness" is

defined as the entropy of the source and is given by the equation

$$H(X) = - \sum_{i=1}^n p_i \log_2 p_i \quad (B.2)$$

where  $p_i$  is the probability of occurrence of the discrete value  $x_i$ . If  $p_i = 0$ , then the term  $\log_2 1/p_i$  is defined equal to 0 (i.e. no contribution to the entropy). Should the range A be of infinite extent (i.e. a continuous source), then the above series may not converge nor would  $p(x)$  necessarily be defined. In this case,  $H(X)$  is defined as positive infinity.

As example (Ref 19:15), let X represent the outcome of a single roll of a fair die. Then  $A=\{1,2,3,4,5,6\}$  and  $p_i = 1/6$  for each i. Here  $H(X) = \sum_6 1/6 \log_2 6 = 2.58$  bits.

In the above example, the statistics governing the occurrence of a given outcome were uniform. This represents the "most random" case with a resultant maximum of the entropy function. Should the die be "loaded", then the predicted outcome is "less random" and hence the value of  $H(X)$  would be reduced.

The next important component in Shannon's communication model is the transmission channel.

### Transmission Channel

The transmission channel is also assumed to be a discrete time, finite alphabet device which accepts and transmits to the receiver one symbol in a finite alphabet  $B$  each  $T_c$  seconds. The alphabet  $B$  is often binary, hence the dimension of the symbol space (defined as  $\|B\|$ ) is 2 and  $B=\{0,1\}$ . The transmission rate of the channel is defined as  $R_c=1/T_c$  channel symbols (bits if binary) per second. If the output of the source process is defined as  $\{X\}$ , and the channel is "noisy", then the received symbol  $\{Y\}$  might not equal the transmitted symbol  $\{X\}$ . A noiseless channel transmits symbols with no error in which case  $\{X\}=\{Y\}$ .

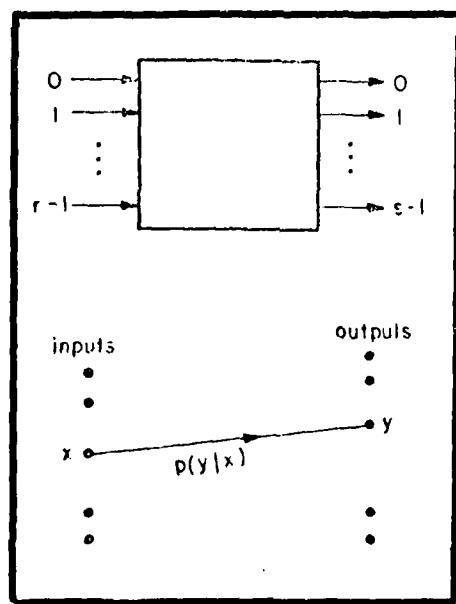


Fig B-2. A discrete memoryless channel (From Ref 19:19)

For simplicity, let the channel be modelled as a

discrete memoryless channel (DMC) as depicted in Figure P-2. The DMC is described mathematically by a conditional probability which relates the chance that a given output  $y$  was the result of a given input  $x$ . If the input to the channel is a random variable  $X$ , and the output is a random variable  $Y$ , then a quantity called conditional entropy can be defined. The equation defining conditional entropy is:

$$H(X|Y) = \sum_{x=1}^n \sum_{y=1}^n p(x,y) \log_2 \frac{1}{p(x|y)} \quad (B.3)$$

For a given pair  $X, Y$  of random variables,  $H(X|Y)$  represents the amount of uncertainty remaining about  $X$  after  $Y$  has been observed.

Now since  $H(X)$  represents the uncertainty about  $X$  before  $X$  is known and  $H(X|Y)$  represents the uncertainty after, the difference  $H(X) - H(X|Y)$  must represent the amount of information provided about  $X$  by  $Y$ . This quantity is called the mutual information between  $X$  and  $Y$ , and is denoted by :

$$I(X;Y) = H(X) - H(X|Y) \quad (B.4)$$

With the above definitions in hand, the most important quantity of a communications channel can be described; that quantity is the channel capacity. Channel capacity is

defined as the maximum amount of information, per unit of time, which can be "reliably" transmitted over the channel. That is :

$$C = \max\{I(X;Y)\} \quad (B.6)$$

Channel capacity is closely related to another important parameter of a communication system known as the rate-distortion function.

Rate-Distortion. Shannon postulated the existence of a mathematical distortion measure,  $d(X,Y)$ , to measure the distortion or loss resulting if a source symbol X is reproduced as Y. Unfortunately this abstract distortion measure is difficult to quantify. As stated by Berger (Ref 4:6), "the unavailability of a distortion measure that is both physically meaningful and analytically tractable constitutes one of the major obstacles to progress in (communication) system design."

Assuming one has such a distortion measure, then associated with most source-user pairs is a function  $R(D)$  called the rate distortion function. The rate distortion is important in that it gives the designer a mathematical tool to measure the amount of distortion that can be expected for a given transmission rate. A communication system can achieve a given fidelity D if and only if the capacity C exceeds  $R(D)$ . Hence  $R(D)$  is the effective rate at which the

source produces information subject to the constraint that the user can tolerate an average distortion of  $D$ .

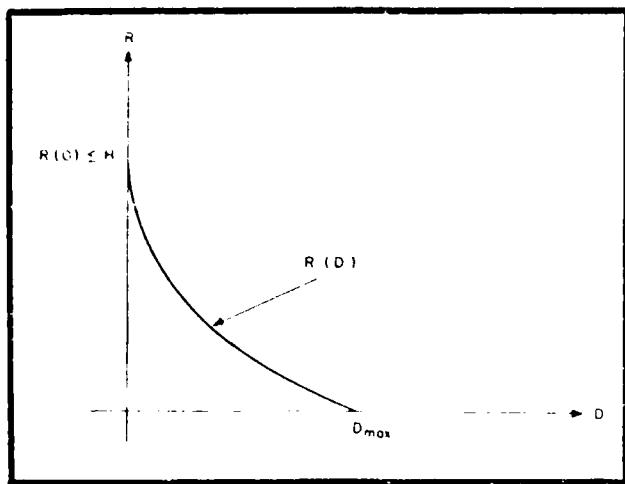


Fig B-3. A typical rate distortion function (From Ref 4:7).

The simplest source-user pair is a discrete memoryless source (DMS) and a single letter fidelity criteria. The DMS produces statistically independent, identically distributed, discrete random variables. Assuming a single letter fidelity criterion, every time the system presents the letter  $y$  to the user when the source output was actually  $x$ , a nonnegative penalty  $p(x, y)$  is determined. The rate-distortion function for the above case is plotted in Figure B-3. As can be seen from the Figure,  $R(0)$  is equal to  $H(X)$ . This last result leads to Shannon's channel coding theorem which states that if the symbol rate is less than the channel capacity, it is possible to transmit with perfect fidelity. That is, if the entropy of the source,

$B(X)$ , is less than or equal to the channel capacity  $C$ , then transmission with zero error can be obtained. This amazing result is not without cost, however. To reduce the source entropy (as seen by the channel) requires coding. Encoder

In most circumstances, the output of the information source is not suitable for direct input into the channel. To match the source to the channel, much like matching impedances in circuit theory, is the job of the encoder.

The encoder incorporates all of the functions which process the source data for transmission over the communication channel. This includes coding, analog-to-digital conversion, and modulation. In order to transmit without loss, integers  $K$  and  $L$  must exist such that  $KT_s = LT_c$ . This guarantees that the received sequence  $\{Y\}$  has the same symbol rate as the transmitted sequence  $\{X\}$ .

As illustrated in Figure B-1, the encoder is divided into two functional subunits. The first of these subunits is the source encoder.

Source Encoder. Source encoding is the operation by which the source output is mapped into an alternate symbol set with the goal of reducing source sequence dependence (i.e. reduce redundancy). Source coding is the transformation in which data compression occurs.

The first source codes, as conceived by Shannon, were block codes where blocks of source symbols were mapped into a single representative channel symbol. This type of

encoding is simple, but efficient only if the source repeats a given, constant length, sequence of symbols on a frequent basis. Another type of source coding is variable length coding.

In variable length codes, those source symbols which occur most frequently are assigned the shortest channel codewords.

One of the most popular variable length codes in use was originally proposed by Huffman (Ref 16:31-34) in the early fifties. This code is in a class known as uniquely decodable (UD) codes. UD codes imply that the codewords, regardless of length, are unique sequences. Hence decoding is instantaneous upon codeword reception(i.e. the decoding does not depend upon reception of the next codeword). McEliece has shown (Ref 19:244-245) that the Huffman code is optimal in the class of UD source codes.

Many techniques for source coding are available. Chapter 2 of this thesis discusses some approaches used in "compressing" electrocardiogram waveforms.

The output of the source coder is generally a sequence of discrete symbols (i.e. binary 1|0) suitable for further processing by the channel coder.

Channel Encoder. Shannon proved that if the rate of the source, as seen by the channel, is less than the channel capacity C, then "noiseless" or error free transmission is theoretically possible. The goal of channel encoder is to combat channel noise to achieve a probability of error ( $P_e$ )

which approaches zero.

Channel encoders allow a reduction of  $P_e$  by selectively reinserting redundancy which has been removed by the source coder. This redundancy allows error detection, and with the proper codes, error correction by the decoder. How close the  $P_e$  approaches zero is purely a function of the effort (and money!) spent on channel coding.

A common channel coding technique is the use of "parity" bits in digital communication. With parity checking, a bit (or bits) is added to the source word which represents the number (odd or even) of "ones" in the word. For example, if odd parity is defined and the source word is 1100101 then a single parity bit of 0 implies an even number of ones (1100101|0).

Parity checking in the example above will detect a single bit error in transmission. A more powerful technique is Hamming codes.

The Hamming code is in a large class of codes known as linear codes. Hamming not only detects errors in transmission, but will correct errors to a certain level. This correction is accomplished by multiplying the received codeword by a matrix known as the syndrome. In the case of Hamming codes, the output from this transformation is the bit position in error.

#### Decoder

As would be expected, the decoder is the inverse

operation which outputs an estimate of the data input to the encoder. The decoder is composed of the channel decoder and the source decoder. The channel decoder uses the redundancy added by the channel encoder to perform error checking and/or correction. The source decoder takes the "correct" data output from the channel decoder and "decompresses" the data to produce an estimate of the information source. Information theory has shown that if enough time, complexity, and money is spent on channel encoding, and the information rate is below the channel capacity, then "error free" transmission is possible.

#### Conclusion

This appendix has been a very brief summary of a very large field of study. The key words underlined throughout this text are terminology which appear in the theory chapters (2 and 3) of the thesis. For a more thorough, mathematical treatment of the fields of Information Theory and Coding, the reader is encouraged to refer to the textbooks by McEliece (Ref 19) and Berger (Ref 4). The IEEE Press and Benchmark book (Ref 10) is an excellent source for a survey of the key papers in the field of Data Compression.

### Appendix C

This appendix contains the 6800 assembly language source programs of the EKG-DAAS.

```

00030      *:
00040      *
00050      * PROGRAM NAME: EKG-IXEC
00060      * AUTHOR : CAPT. RNL TOLAN
00070      * VERSION : 1.8
00080      * VERSION DATE 2 OCT 80
00090      *
00100      *
00110      * PROGRAM DESCRIPTION
00120      *
00130      * THIS PROGRAM IS THE EXECUTIVE ROUTINE WHICH
00140      * CONTROLS THE EXORCISER EKG DATA ACQUISITION
00150      * SYSTEM. THIS ROUTINE CALLS OVERLAYERED PROGRAMS
00160      * WHICH PERFORM DATA COLLECTION, COMPRESSION, STORAGE
00170      * AND RECONSTRUCTION.
00180      * THIS SOFTWARE IS IN SUPPORT OF THESIS RESEARCH
00190      * TO IDENTIFY THE MOST EFFICIENT EKG DATA COMPRESS-
00200      *ION ALGORITHM IN THE TEST SET.
00210      *
00220      * COMMAND OPTIONS
00230      *
00240      * 0=STORAGE WITHOUT COMPRESSION (10 BIT)
00250      *
00260      * 1=COMPRESS WITH ALGOR TOLAN-A
00270      *
00280      * 2=COMPRESS WITH ALGOR TOLAN-B
00290      *
00300      * 3=COMPRESS WITH ALGOR DOWER
00310      *
00320      * 4=COMPRESS WITH ALGOR 2ND ORDER INTERPOL
00330      *
00340      * 5=COMPRESS WITH ALGOR TURNPT
00350      *
00360      * 6=DISPLAY COLLECTED DATA & STATS
00370      *
00380      * 7=JUMP PROGRAM CONTROL TO DOS
00390      *
00400      * 8=JUMP PROGRAM CONTROL TO EXBUG
00410      *
00420      * 9=LOAD & EXECUTE OTHER OVERLAYS
00430      *
00440      * S=SAVE CUR MEM FILE TO DISK
00450      *
00460      * START OF PROGRAM
00470      *
00480      *
00490 1D00      ORG      $1D00      PROGRAM START LOCATION
00500      *
00510      OPT      O      ASSB OPT TO CREATE Q11 FILE
00520      OPT      NOG      ASSB OPT TO SUP FCC LIST
00530      *
00540      *
00550      * LABEL DECLARATIONS
00560      *

```

* SUPPORT SUBROUTINE ADDRESSES				
*				
00570	C75B	CLRPAS EQU	\$C75B	EOS. CLR PASSWD BUFFER
00580	CE52	KEYRD0 EQU	\$CE52	EOS. KEYRD INPUT ROUTINE
	CE53	KEYRD1 EQU	\$CE53	EOS. KEYRD INPUT ROUTINE
00610	CA36	KEYRD0 EQU	\$CA36	EOS. KEYRD INPUT ROUTINE
00620	F000	EXBUG EQU	\$F000	EXBUG. EXBUG ENTRY PT
00630	2800	DOS EQU	\$2800	DOS. DOS ENTRY PT
00640	C75B	WRITEO EQU	\$C75B	EOS. CLR PASSWD BUFFER
00650	CC87	CLRPAS EQU	\$CC87	EOS. CLR PASSWD BUFFER
00660	C7C8	DRIVE EQU	\$C7C8	EOSIO. DSK DRIVE SELECT
00670	C803	RLIB EQU	\$C803	EOSIO.
00680	C807	WLIB EQU	\$C807	EOSIO.
00690	CE52	LOAD2 EQU	\$CE52	EOS. LOAD PRGM ROUTINE
00700	CBDA	DOSTR2 EQU	\$CBDA	EOS. ALT FCS ENTRY LOC
00710	E055	BYTE EQU	\$E055	MIKEUG. GET TO HEX DIG PRM TE
00720		*		
00730		*		
00740		*		
00750	3066	NAME EQU	\$3066	NAME BUFR FOR FILE I/O
00760	3060	TEMPX1 EQU	\$3060	TEMP 2 BYTE STOR BUFR
00770	3062	TEMPX2 EQU	\$3062	TEMP 2 BYTE STOR BUFR
00780	3058	STARTX EQU	\$3058	LOWEST ADDR USED IN PRGM
00790	305A	ENDX EQU	\$305A	HIGHEST ADDR USED IN PRGM
00800	305C	GOX EQU	\$305C	START EXECUTE ADDR
00810	001B	PROGX EQU	\$001B	DISK I/O ERR VEC ADDR
00820	0008	EMEMH EQU	\$0008	
00830	3002	ENDBUF EQU	\$3002	BUF WITH ADDR OF LAST CHAR IN
00840	3400	HDRSTR EQU	\$3400	START OF MEM FILE HDR SECTOR
00850	3004	BUFFER EQU	\$3004	KEYED INPUT BUFR START
00860	3057	TYPE EQU	\$3057	TYPE OF FILE FOR I/O
00870	1C96	STKSAV EQU	\$1C96	TEMP STACK SAVE BUFR
00880	1C98	CPRTYP EQU	\$1C98	COMPRESSION TYPE BUFR
00890	1C9A	COUNT EQU	\$1C9A	GENERAL 8 PIT COUNTER
00900	1C9B	NAMPTR EQU	\$1C9B	NAME POINTER FOR CONSOLE 1/0
00910	1C9D	VECSAV EQU	\$1C9D	I/O VEC SAVE BUFR
00920	1C9F	OLAYGO EQU	\$1C9F	OVERLAY EXECUTE FLAG
00930	1CA0	LGOFIG EQU	\$1CA0	PRSTAT LOAD VS EXECUTE FIG
00940	1CA1	FILHLC EQU	\$1CA1	FILER SUB ADDR PASS BUF
00950	1CA3	SAVEFL EQU	\$1CA3	SAVEFL SUB ADDR PASS BUF
00960	1CA5	HXASILC EQU	\$1CA5	HXASC SUB ADDR PASS BUF
00970	1CA7	HXBUP EQU	\$1CA7	HXASC PARAMETER BUFR
00980	1CA9	PDFPLC EQU	\$1CA9	PDFPRT SUB ADDR PASS BUF
00990	1CAB	OVRLLC EQU	\$1CAB	OVRLLC SUB ADDR PASS BUF
01000	1CAD	OVRBUF EQU	\$1CAD	OVRBUF PARAMETER PASS BUFR
01010	3490	LOCPC1 EQU	\$3490	TOTL CAL LOOPS ERRC IN CH1
01020	3494	SAMPNO EQU	\$3494	NUM OF SAMPLES TAKEN
01030	3496	LPCAL EQU	\$3496	NUM OF CAL LOOPS ERRC 1 INTR
01040	3497	MAXZ EQU	\$3497	MAX VLU IN CH Z
01050	3498	MAXZLO EQU	\$3498	LOC OF MAX VLU IN CH Z
01060	349A	MINZ EQU	\$349A	MIN VLU IN CH Z
01070	349B	MINZLO EQU	\$349B	LOC OF MIN VLU IN CH Z
01080	349D	MAXY EQU	\$349D	MAX VLU IN CH Y
01090	349E	MAXYLO EQU	\$349E	LOC OF MAX VLU IN CH Y
01100	34A0	MINY EQU	\$34A0	MIN VLU IN CH Y

01110	34A1	MINXLO EQU	\$34A1	LOC OF MIN VLU IN CH Y
01120	34A3	MAXX EQU	\$34A3	MAX VLU IN CH X
01130	34A4	MAXYLO EQU	\$34A4	LOC OF MAX VLU IN CH X
01140	34A6	HIX X EQU	\$34A6	HIN VLU IN CH X
01150	34A7	MINXLO EQU	\$34A7	LOC OF MIN VLU IN CH X
01160	34A9	MEMBIT EQU	\$34A9	NUM OF BITS AVAIL FOR STO
01170	34AC	DTABIT EQU	\$34AC	NUM OF BITS USED TO STO DTA
01180	34B0	XBITS EQU	\$34B0	NUM OF BITS USED TO STO X
01190	34B3	YBITS EQU	\$34B3	NUM OF BITS USED TO STO Y
01200	34B6	ZBITS EQU	\$34B6	NUM OF BITS USED TO STO Z
01210	34B9	TBITS EQU	\$34B9	NUM OF BITS USED TO STO TIME
01220	34BC	ACELCT EQU	\$34BC	# BITS FED TO VAR LEN CODER
01230	34C2	BASSAV EQU	\$34C2	SAVE LOC FOR VLU \$3620 PRM
01240	3460	MAXMIN EQU	\$3460	START OF MAX,MIN ASCII PUFFER
01250	34C3	ENTRPY EQU	\$34C3	START OF ENTROPY TABLE IN ASC
01260	3C00	SECZRO EQU	\$3C00	SEC 0 OF MEM FILE
01270	0019	SAVEX EQU	\$0019	TEMP LOC TO SAVE INDEX REG
01280	3500	XPDF EQU	\$3500	LOC OF X PDF BUFFER
01290	3700	YPDF EQU	\$3700	LOC OF Y PDF BUFFER
01300	3900	ZPDF EQU	\$3900	LOC OF Z PDF BUFFER
01310	3B00	TPDF EQU	\$3B00	LOC OF TIME PDF BUFFER
01320	8000	BUFEND EQU	\$8000	END OF MEM BUF
01330		*		
01340		* HARDWARE ADDRESSES		
01350		*		
01360	E500	DACZRO EQU	SE500	DAC 0 ADDRESS
01370	E502	DAONE EQU	SE502	D :THIS ROUTINE IS THE EXECUTIVE
01460		* CONTROLLER OF THE EKG DATA ACQ SYS		
01470		*		
01480		*		
01490	1D00 OF	START	SEI	STOP POSSIBLE INTR ON RESET
01500	1D01 CE 4000	LDX	#\$4000	
01510	1D04 FF E500	STX	DACZRO	CLR DACS & SET SEL 1 HIGH
01520	1D07 FF E502	STX	DACONE	
01530	1D0A CE 214A	LDX	#FILEHDR	
01540	1D0D FF 1CA1	STX	FILHLC	PUT FILEHDR ADDR IN PASS BUFF
01550	1D10 B6 3620	LDA A	\$3620	
01560	1D13 R7 34C2	STA A	BASSAV	SAVE CURRENT #\$3620 IN BUF FO
01570	1D16 CE 208E	LDX	#SAVEFL	PUT SAVEFL ADDR IN PASS BUF
01580	1D19 FF 1CA3	STX	SAVELC	
01590	1D1C CE 23DB	LDX	#HXASC	PUT HXASC ADDR IN PASS BUF
01600	1D1F FF 1CA5	STX	HXASLC	
01610	1D22 CE 22C2	LDX	#PDFPRT	PUT PDFPRT ADDR IN PASS BUF
01620	1D25 FF 1CA9	STX	PDFPLC	
01630	1D28 CE 2067	LDX	#OVRLAY	PUT OVRLAY ADDR IN PASS BUF
01640	1D2B FF 1CAB	STX	OVTLIC	
01650	1D2E CE 1D00	LDX	#START	PUT START ADDR IN PASS BUF
01660	1D31 FF 2801	STX	DOS+1	
01670	1D34 7F 1C9F	CLR	OLAYGO	CLR OVLAY PCN FLAG
01680	1D37 7F 1CA0	CLR	LGOFIG	CLR ERSTAT LOAD FIG
01690	1D3A 4F	CLR A		
01700	1D3B FD C7C8	JSR	DRIVE	INSURE EXEC FILE TO DRIVE 0
01710	1D3E BD 1F6F	JSR	CMDIN	PROMPT & GET IN CH7
01720	1D41 C1 30	CMP B	'0	IS CMD ASCII 0?

01730 1D43 26 09		BNE	EXCM1	NO. CHECK OTHER CIDS
01740 1D45 CE 1DC7		LDX	#NOCPR	LOAD OVERLAY NAME PTR
01750 1D45 FF 1CND		STX	OUTPUT	
01760 1D45 73 2C67		CMP	OVRLAY	LOAD & EXECUTE OVRLAY
01770 1D4E C1 31	EXCM1	CMP B	#'1	
01780 1D50 26 09		BNE	EXCM2	
01790 1D52 CE 1DCF		LDX	#TOLAN1	
01800 1D55 FF 1CAD		STX	OVRLBUF	
01810 1D58 7E 2067		JMP	OVRLAY	
01820 1D5B C1 32	EXCM2	CMP B	#'2	
01830 1D5D 26 09		BNE	EXCM3	
01840 1D5F CE 1DD7		LDX	#TOLAN2	
01850 1D62 FF 1CAD		STX	OVRLBUF	
01860 1D65 7E 2067		JMP	OVRLAY	
01870 1D68 C1 33	EXCM3	CMP B	#'3	
01880 1D6A 26 09		BNE	EXCM4	
01890 1D6C CE 1EDF		LDX	#DOWNR	
01900 1D6F FF 1CAD		STX	OVRLBUF	
01910 1D72 7E 2067		JMP	OVRLAY	
01920 1D75 C1 34	EXCM4	CMP B	#'4	
01930 1D77 26 09		BNE	EXCM5	
01940 1D79 CE 1DE6		LDX	#INTER	
01950 1D7C FF 1CAD		STX	OVRLBUF	
01960 1D7F 7E 2067		JMP	OVRLAY	
01970 1D82 C1 35	EXCM5	CMP B	#'5	
01980 1D84 26 09		BNE	EXCM6	
01990 1D86 CE 1DEE		LDX	#TURNPT	
02000 1D89 FF 1CAD		STX	OVRLBUF	
02010 1D8C 7E 2067		JMP	OVRLAY	
02020 1D8F C1 36	EXCM6	CMP B	'6	
02030 1D91 26 09		BNE	EXCM7	
02040 1D93 CE 1DF6		LDX	#DISPI	
02050 1D96 FF 1CAD		STX	OVRLBUF	
02060 1D99 7E 2067		JMP	OVRLAY	
02070 1D9C C1 37	EXCM7	CMP B	#'7	
02080 1D9E 26 09		BNE	EXCM8	
02090 1DA0 CE 2FC0		LDX	#\$2BCD	
02100 1DA3 FF 2801		STX	DOS+1	
02110 1DA6 7E 2800		JMP	DOS	
02120 1DA9 C1 38	EXCM8	CMP B	#'8	
02130 1DAB 26 03		BNE	EXCM9	
02140 1DAD 7E F000		JMP	EXBUG	
02150 1DB0 C1 39	EXCM9	CMP B	#'9	
02160 1DB2 26 09		BNE	EXCM10	
02170 1DB4 CE 1DPE		LDX	#OVLJSC	
02180 1DB7 FD CA87		JSR	OUTPUT	
02190 1DRA 7E C0DA		JMP	DOSTR2	
02200 1DPD C1 53	EXCM10	CMP B	#'S	
02210 1DPF 26 03		BNE	EXCM11	
02220 1DC1 FD 208E		JSR	SAVFL	
02230 1DC4 7E 1D00	EXCM11	JMP	START	
02240	*			
02250 1DC7 4E	NOCPR	FCC	/NOCPRS*	/
02260 1DCF 54	TOLAN1	FCC	/TOLAN-A*/	

0210 1DD7 54	TOLAN2	FCC	/TOLAN-B*/
02280 1DDE 44	DOWER	FCC	/DOWER* /
0230 1E49 49	INTER	FCC	/INTERPL* /
0230 1E4A 44	TURNPT	FCC	/TURNPT* /
02310 1DF6 44	DISPL	FCC	/DISPLAY*/
02320 1DFE 1A	OVLMSG	FCB	\$1A
02330 1DFF 45		FCC	/ENTER "RUN FILENAME" TO LOAD & EXECUT
02340 1E25 4F		FCC	/OVERLAY "FILENAME"./
02350 1E38 0D0A		FDB	\$0D0A
02360 1E3A 28		FCC	/(WARNING: OVERLAY MUST FIT BETWEEN /
02370 1E5D 30		FCC	/0100-1400)/
02380 1E67 0D0A		FDB	\$0D0A,\$0D0A,\$003F,\$2004
02390	*		
02400	*		* END EKG-EXEC
02410	*		
02420	*		*FUNCTION :CMDIN
02430	*		*INPUTS (REG) :NONE
02440	*		*OUTPUTS (REG):B
02450	*		*CALLS :OUTPUT,KEYRD0
02460	*		*DESTROYS :ALL REGS
02470	*		*PURPOSE :TO PRINT PROMPT MESSAGE AND READ COMMAND
02480	*		* FROM CONSOLE.
02490	*		
02500 1E6F CE 1E7F	CMDIN	LDX	#CMDMSG LOAD OUT PTR
02510 1E72 BD CA8F		JSR	OUTMCR
02520 1E75 BD CA36		JSR	KEYEDO
02530 1E78 FE 3002		LDX	ENDBUF SET PRT TO LAST CHAR ENTERED
02540 1E7B 09		DEX	
02550 1E7C E6 00		LDA B	0,X
02560 1E7E 39		RTS	
02570	*		
02580 1E7F 1A07	CMDMSG	FDB	\$1A07,\$0D0A,\$0A0A,\$0A0A,\$0A0A
02590 1E89 45		FCC	/EKG DATA ACQUISITION SYSTEM/
02600 1EA4 0D0A		FDB	\$0D0A,\$0D0A
02610 1EA8 43		FCC	/COMMAND OPTIONS:/
02620 1EB8 0D0A		FDB	\$0D0A,\$0D0A
02630 1EC4 44		FCC	/DATA ACQUISITION/
02640 1ECC 0D0A		FDB	\$0D0A,\$0D0A
02650 1ED0 20		FCC	/ 0=STORAGE WITH NO COMPRESSION/
02660 1EEE 0D0A		FDB	\$0D0A
02670 1EF0 20		FCC	/ 1=COMPRESSION VIA TOLAN ALGORITHM A
02680 1F14 0D0A		FDB	\$0D0A
02690 1F16 20		FCC	/ 2=COMPRESSION VIA TOLAN ALGORITHM B
02700 1F3A 0D0A		FDB	\$0D0A
02710 1F3C 20		FCC	/ 3=COMPRESSION VIA DOWER ALGORITHM/
02720 1F5E 0D0A		FDB	\$0D0A
02730 1F60 20		FCC	/ 4=COMPRESSION VIA 2ND ORDER INTERPL
02740 1F89 0D0A		FDB	\$0D0A
02750 1F8B 20		FCC	/ 5=COMPRESSION VIA TURNING POINT ALG
02760 1FB1 0D0A		FDB	\$0D0A,\$CD0A
02770 1FB5 44		FCC	/DATA DISPLAY/
02780 1FC1 0D0A		FDB	\$0D0A,\$0D0A
02790 1FC5 20		FCC	/ 6=PRINT DATA STATISTICS OR /
02800 1FE1 44		FCC	/DISPLAY DATA/

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02810 1FED CD0A      FDB    $0D0A,$0D0A
02820 1FF1 50         FCC    /PROGRAM CONTROL/
02830 2000 CD0A      FDB    $CD0A,$0D0A
02840 2004 20         FCC    / 7=RETURN TO DOS/
02850 2014 0D0A      FDB    $0D0A
02860 2016 20         FCC    / 8=RETURN TO EXBUG/
02870 2028 0D0A      FDB    $0D0A
02880 202A 20         FCC    / 9=LOAD & EXECUTE USER ENTERED OVERL
02890 2050 0D0A      FDB    $0D0A,$0D0A
02900 2054 45         FCC    /ENTER COMMAND NOW=/
02910 2066 04         FCB    4
02920 *
02930 * END CMDIN
02940 *
02950 *
02960 *FUNCTION :OVERLAY
02970 *INPUTS (REG) :X
02980 *OUTPUTS (REG) :NAME
02990 *CALLS :CLPNAM,PUTNAM,OVRBUF
03000 *DESTROYS :ALL REGISTERS
03010 *PURPOSE :TO LOAD AND THEN EXECUTE OVERLAYS
03020 * ASSOCIATED WITH EKG-EXEC &
03030 * SPECIFIED BY THE INPUT COMMAND.
03040 *
03050 *
03060 *
03070 2067 BD 20E1 OVERLAY JSR    CLPNAM CLR NAME BUFFER
03080 206A BD CC87      JSR    CLRPA$ CLEAR PASSWORD BUFFER
03090 206D CE 3066      LDX    #NAME PUT DOS FILENAME BUF IN NAMPT
03100 2070 FF 1C9B      STX    NAMPTR
03110 2073 FE 1CAD     LDX    OVRBUF POINT X TO OVERLAY NAME
03120 2076 ED 20EF     JSR    PUTNAM PUT OVERLAY NAME IN NAME
03130 2079 86 22       LDA A #$22 DEFINE FILE TYPE
03140 207B B7 3057     STA A TYPE
03150 207E BD CE52     JSR    LOAD2 LOAD FILE POINTED TO BY STK
03160 2081 86 AA       LDA A #$AA SET UP COMPARE FOR OVERLAY FCN
03170 2083 B1 1C9F     CMP A OLVGO IS FLAG TRUE?
03180 2086 27 05       BEQ    OVRRTS YES. TREAT OVERLAY AS SUBR
03190 2088 8E A049     LDS    #$A049 NO. TREAT OVERLAY AS ABS JUMP
03200 208B 6E 00       JMP    0,X JUMP TO OVERLAY
03210 208D 39         OVRRTS RTS
03220 *
03230 * END OVERLAY
03240 *
03250 *
03260 *FUNCTION :SAVEFL
03270 *INPUTS (REG) :NONE
03280 *OUTPUTS (REG) :NONE
03290 *CALLS :CLRPA$,DRIVE,WLIB,WRITED0,OUTPUT
03300 *DESTROYS :ALL REGISTERS
03310 *PURPOSE :TO SAVE EKG DATA FILES RESIDING IN
03320 *IN MEMORY LOCATIONS 3A00-7FFF ON DISK. FILE NAME IS
03330 * PICKED OFF OF FILE HEADER AT LOC 3A01-3A02.
03340 *

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03350	*				
03360 208H ED 51	SAVFILE ISR	CLRNAM	CLEAR NAME BUFFER		
03370 206D C7 0141	LDA	PUTNAME	PUT FILE NAME IN BUFFER		
03380 2093 ED 4F	PTR	CLNAMO	FILL IN NAME		
03390 2095 R6 3620	IDA A	\$3620	GET ITEM LOC FOR SAVE FROM PWS		
03400 2098 R7 34C2	STA A	BASSAV	SAVE IN BUFFER		
03410 209B CE 2141	IDK	#NAME\$G	SET UP PTR TO FRT NAME		
03420 209E FF 1C9B	STX	NAMPTR			
03430 20A1 CE 3402	IDX	#HDRSTR+2	POINT X TO FILENM IN DATA		
03440 20A4 ED 49	BSR	PUTNAM	PUT FILENM IN "NAME" BUF		
03450 20A6 86 22	IDA A	#\$22	SET TYPE OF FILE		
03460 20A8 B7 3057	STA A	TYPE			
03470 20AB CE 3400	LDX	#HDRSTR	LOAD BUFFER START		
03480 20AE FF 3058	STX	STARTX			
03490 20B1 CE 7FFF	LDX	#BUFIEND-1	LOAD BUFFER END		
03500 20B4 FF 305A	STX	ENDX			
03510 20B7 CE 2800	IDX	#\$2800	LOAD LGO ADDRESS		
03520 20BA FF 305C	STX	COX			
03530 20BD DF 1B	STX	PROGX			
03540 20BF BD CC87	JSR	CLRPAS	CLR PASSWD BUFFER		
03550 20C2 CE 2119	LDX	#DRVMSG			
03560 20C5 BD CA8F	JSR	OUTNCR			
03570 20C8 BD E055	JSR	BYTE			
03580 20CB BD C7C8	JSR	DRIVE	SELECT DRIVE 01 FOR STORE		
03590 20CE BD C807	JSR	WLIB			
03600 20D1 DE 08	LDX	EMEMH			
03610 20D3 BD C75B	JSR	WRITE0	WRITE BUFFER TO DISK		
03620 20D6 4F	CLR A				
03630 20D7 BD C7C8	JSR	DRIVE	RESET DRIVE BACK TO 0		
03640 20DA CE 2123	LDX	#CATMSG	LOAD OUTPUT POINTER		
03650 20DD BD CA87	JSR	OUTPUT			
03660 20E0 39	RTS				
03670	*				
03680 20E1 CE 3066	CLRNAM	LDX	POINT TO NAME BUFFER		
03690 20E4 C6 08	CLRNAMO	LDA B	#\$8	SET UP CLEAR COUNT	
03700 20E6 86 20		LDA A	#\$20	ASCII SPACE	
03710 20E8 A7 00	CLRNAM	STA A	0,X	FILL NAME BUF WITH SPACES	
03720 20EA 08		INX			
03730 20EB 5A		DEC B			
03740 20EC 26 FA		RNE	CLRNAM		
03750 20EE 39		RTS			
03760	*				
03770 20EF BF 1C96	PUTNAM	STS	STKSAV	SAVE STACK POINTER	
03780 20F2 7F 1C9A	CLR	COUNT		CLEAR COUNTER	
03790 20F5 35	TXS			TRANSFER X TO PTR	
03800 20F6 CE 3066	IDX	#NAME		PICK UP POS FILE NAME IN PTR	
03810 20F9 C6 08	LDA B	#8		SET UP NAME LENGTH IN PTR	
03820 20FB 32	PUTNAM	HUL A		PULL CHAR FROM PTR	
03830 20FC 81 04	CMP A	#\$04		IS IT 4?	
03840 20FE 27 15	BPL	PUTNA2		YES. NAME ENDED	
03850 2100 A7 00	STA A	0,X		NO. STORE CHAR IN "NAME"	
03860 2102 08	INX				
03870 2103 FF 3060	STX	TEMUX1		SAVE INDEX PTR	
03880 2106 FF 1C9B	IDX	NAMPTR		LOAD CURRENT NAME PTR	

03890 2109 A7 00	STA A	0,X	STORE IN NAMEHS	
03900 210B 08	INX		INCR POINTER	
03910 210C 00 1C93	STX	#NAMEP		
03916 210F F0 3060	LDX	TEMPX1	RESTORE INDEX REG	
03930 2112 5A	DEC B			
03940 2113 26 E6	RNE	PUTNA1	8 CHARS YET?	
03950 2115 FE 1C96	LDS	STKSAV	YES. RESTORE STACK	
03960 2118 39	RTS			
03970 *				
03980 2119 0D0A	DRVMSG	FDB	\$0D0A	
03990 211B 44		FCC	/DRIVE? /	
04000 2122 04		FCB	4	
04010 2123 1A07	CATMSG	FDB	\$1A07,\$0D0A,\$0D0A	
04020 2129 44		FCC	/DATA SAVED ON DISK FILE /	
04030 2141 0008	NAMMSG	RMB	8	
04040 2149 04		FCB	4	
04050 *				
04060 *	END SAVEFL			
04070 *				
04080 *				
04090 *FUNCTION :FILEHDR				
04100 *INPUTS (REG) :NONE				
04110 *OUTPUTS (REG) :NONE				
04120 *CALLS :OUTPUT,KEYBD0				
04130 *DESTROYS :ALL REGISTERS				
04140 *PURPOSE :TO INIT DATA BUFFER WITH ACCOUNT				
04150 * DATA (NAME,DATE,TIME,ETC) AND CLEAR FOR NEW				
04160 * DATA STORAGE.				
04170 *				
04180 214A BD 21D7	FILERX	JSR	CLRBUF	FILL DATA BUF WITH NULS
04190 214D BD 21C9		JSR	CLRSEC	FILL MEM FILE SEC 3400 WITH A
04200 2150 BD 2288		JSR	PDFCLR	INITIALIZE COIL PARAM
04210 2153 FE 1C98	LDX		CPRTYP	PICK UP COMPR TYPE
04220 2156 FF 3400	STX		HDRSTR	STR IN HEADER SECTOR
04230 2159 CE 21E5	LDX	#NAMEHS		OUT FTR
04240 215C BD CA8F	JSR		OUTNCR	"NAME FILE PLEASE?"
04250 215F BD CA36	JSR		KEYBD0	
04260 2162 C6 02	LDA B	#\$02		HEADER STOR OFFSET
04270 2164 8D 42	RSR		TXTSTR	STORE FILENAME IN HEADER SEC
04280 2166 37	PSH B			SAVE NEXT AVAIL HEAD LOC
04290 2167 CF 2201	LDX	#SURJMS		
04300 216A BD CA8F	JSR		OUTNCR	"TEST SUBJECT ID?"
04310 216D BD CA36	JSR		KEYBD0	
04320 2170 33	PUL B			RETRIEVE STOR OUTPUT
04330 2171 8D 35	RSR		TXTSTR	
04340 2173 37	PSH B			
04350 2174 CE 221F	LDX	#RATMS		"COLLECTION SAMPLE DATE?"
04360 2177 BD CA8F	JSR		OUTNCR	
04370 217A BD CA36	JSR		KEYBD0	
04380 217D 33	PUL B			
04390 217E 8D 28	RSR		TXTSTR	
04400 2180 37	PSH B			
04410 2181 CE 2247	LDX	#DATMS		"DATE (05 JUL 80)?"
04420 2184 BD CA8F	JSR		OUTNCR	

04440 2187 PD CA36	JSR KEYBD0	
04440 218A 33	PUL B	
04440 218B 1B	JSR TXTSTR	
04440 218C 37	PSH B	
04470 218E CE 225E	LDX #TIMEMS	
04480 2191 BD CA8F	JSR OUTNCR "TIME (1420)?"	
04490 2194 BD CA36	JSR KEYBD0	
04500 2197 33	PUL B	
04510 2198 8D 0E	BSR TXTSTR	
04520 219A 37	PSH B	
04530 219B CE 2270	LDX #CONTMS	
04540 219E BD CA8F	JSR OUTNCR "COMENTS (MAX 80)?"	
04550 21A1 ED CA36	JSR KEYBD0	
04560 21A4 33	PUL B	
04570 21A5 8D 01	BSR TXTSTR	
04580 21A7 39	RTS	
04590 *		
04600 21A8 RF 1C96	TXTSTR STS	STKSAV SAVE STACK
04610 21AB 8E 3003	LDS #BUFFER-1	POINT TO INPUT BUFFER
04620 21AE CE 3400	LDX #HDRSTR	POINT TO HEADER
04630 21B1 4F	CLR A	SET A TO ZERO
04640 21B2 4C	TXTST1 INC A	
04650 21B3 08	INX	
04660 21B4 11	CBA	IS A = OFFSET?
04670 21B5 2D FB	PLT	TXTST1 NO. KEEP INCR'ING
04680 21B7 32	TXTST2 PUL A	YES. GET CHAR FRG! STACK
04690 21B8 81 04	CMP A #\$04	IS IT 4
04700 21C0 27 06	BREQ	TXTRTN YES. END TEXT, RETURN
04710 21C1 A7 00	STA A 0,X	NO. STORE CHAR IN HEADER
04720 21C2 5C	INC B	
04730 21C3 08	INX	
04740 21C0 20 F5	BRA	TXTST2
04750 21C2 A7 00	TXTRTN STA A 0,X	
04760 21C4 5C	INC B	
04770 21C5 BE 1C96	LDS	STKSAV RESTORE CTACK
04780 21C8 39	RTS	
04790 21C9 CE 3400	CLRSEC LDX	#HDRSTR POINT TO DATA BUF START
04800 21CC 86 20	LDA A #\$20	ASCII SPACE
04810 21CE A7 00	CLRSE1 STA A 0,X	STORE SPACE
04820 21D0 08	INX	
04830 21D1 8C 3500	CPX #XPDF	IS SECTOR SPACED?
04840 21D4 26 F8	BNE CLRSE1	COUNT > ZRO?
04850 21D6 39	RTS	NO. RTN
04860 *		
04870 21D7 CE 3400	CLRRUF LDX	#HDRSTR POINT TO DATA BUF
04880 21DA 86 00	LDA A #00	ASCII NULL
04890 21DC A7 00	CLRBUI STA A 0,X	STORE NULL
04900 21DE 08	INX	
04910 21DF 8C 8000	CPX #BUFLND	IS INDEX REG AT END?
04920 21E2 26 F8	BNE CLRBUI	NO. KEEP GOING
04930 21E4 39	RTS	YES. RETURN
04940 *		
04950 21E5 0D0A	NAMES FDB	\$0D0A,\$0D0A
04960 21E9 4E	FCC	/NAME DATA FILE PLEASE? /

04970	2200	04	PCB	4
05000	2201	0D0A	SURMS FDR	\$0D0A,\$0D0A
				/COLLECTION SAMPLING RATE (500 Hz)? /
05000	221E	04	PCB	4
05010	221F	0D0A	RATEMS FDR	\$0D0A,\$0D0A
05020	2223	43	FCC	/COLLICITION SAMPLING RATE (500 Hz)? /
05030	2246	04	PCB	4
05040	2247	0D0A	DATEMS FDB	\$0D0A,\$0D0A
05050	224B	44	FCC	/DATE (05 JUL 80)? /
05060	225D	04	PCB	4
05070	225E	0D0A	TIMEMS FDB	\$0D0A,\$0D0A
05080	2262	54	FCC	/TIME (1423)? /
05090	226F	04	PCB	4
05100	2270	0D0A	CONTMS FDB	\$0D0A,\$0D0A
05110	2274	43	FCC	/COMMENTS (MAX=80)? /
05120	2287	04	PCB	4
05130			*	
05140			*	END FILHDR
05150			*	
05160			*	FUNCTION :PDFCLR
05170			*	INPUTS (REG) :NONE
05180			*	OUTPUTS (REG) :NONE
05190			*	CALLS :NOTHING
05200			*	DESTROYS :A,B,X,CC
05210			*	PURPOSE :TO INITIALIZE STATISTIC VAR & CLEAR
05220			*	PDF BUFFERS
05230			*	
05240	2288	4F	PDFCLR CLR A	FILL PARAM BUFF WITH 0
05250	2289	CE 3490	LDX #LOOPCT	
05260	228C	A7 00	PDFCL0 STA A	0,X
05270	228E	08	INX	
05280	228F	8C 34C3	CPX #ENTRPY	FILL WITH 0 UP TO ENTROPY BUF
05290	2292	26 F8	BNE PDFCL0	
05300	2294	CE 021F	LDX #\$021F	STORE MAX POSSIBLE BITS IN ME
05310	2297	86 F0	LDA A #\$F0	
05320	2299	FF 34A9	STX MEMBIT	
05330	229C	B7 34AB	STA A MEMBIT+2	
05340	229F	4F	CLR A	
05350	22A0	CE 3500	LDX #XPDF	NOW CLR PDF BUF AREA
05360	22A3	A7 00	PDFCL1 STA A	0,X
05370	22A5	08	INX	
05380	22A6	8C 3C00	CPX #SECZRO	
05390	22A9	26 F8	BNE PDFCL1	
05400	22AB	86 F0	LDA A #\$80	NOW SET UP INITIAL MAX & MINS
05410	22AD	C6 7F	LDA B #\$7F	
05420	22AF	B7 3497	STA A MAXZ	
05430	22B2	B7 349D	STA A MAXY	
05440	22B5	B7 34A3	STA A MAXX	
05450	22B8	F7 349A	STA B MINZ	
05460	22BB	F7 34A0	STA B MINY	
05470	22BC	F7 34A6	STA B MINX	
05480	22C1	39	RTS	
05490			*	
05500			*	END PDFCLR

06050 2327 DE 19	LDX	SAVEX	PUT CUR PDF ADDR
06060 2329 A6 00	LDA A	0,X	LDA MSB BYTE OF PDF VLJU
06110 232D BC 61	LDY B	1,X	LD LS1
06160 232D CC 23CD	LDX	#PDFMS5	
06090 2330 FF 1CA7	STX	HXBUF	
06100 2333 BD 23DB	JSR	HXASC	CONV & STO MSB IN ASC STR
06110 2336 17	TBA		
06120 2337 CE 23CF	LDX	#PDFMS5+2	
06130 233A FF 1CA7	STX	HXBUF	
06140 233D BD 23DB	JSR	HXASC	CONV & STORE LSB IN ASC STR
06150 2340 CE 23BB	LDX	#PDFMS3	
06160 2343 BD CA87	JSR	OUTPUT	PRINT STRING
06170 2346 32	PUL A		RETRIEV VLJU INDEX
06180 2347 DE 19	LDX	SAVEX	RETRIEC CUR PDF ADDR
06190 2349 39	RTS		
06200	*		
06210 234A 1A0C	PDFMS1	FDB	\$1A0C,\$0D0A,\$0A0A,\$0A0A,\$0A0A
06220 2354 3A		FCC	/:FILE /
06230 235A 0008	PDFNAM	RMB	8
06240 2362 0D0A		FDB	\$0D0A,\$CD0A
06250 2366 43		FCC	/:CHANNEL /
06260 236E 0001	PDFMS2	RMB	1
06270 236F 20		FCC	/ AMPLITUDE DISTRIBUTION/
06280 2386 0D0A		FDB	\$0D0A,\$0D0A
06290 238A 44		FCC	/:DATA VALUE/
06300 2394 20		FCC	/ NUMBER OF OCCURRENCES/
06310 23D6 0D0A		FDB	\$0D0A,\$CD0A
06320 23EA 04		FCB	4
06330 23EB 20	PDFMS3	FCC	/ /
06340 23EF 0002	PDFMS4	RMB	2
06350 23C1 2E		FCC	/:...../
06360 23CD 0004	PDFMS5	RMB	4
06370 23D1 20		FCC	/ (HEX)/
06380 23D7 04		FCB	4
06390 23D8 0D23	PDFMS6	FDB	\$0D23
06400 23DA 04		FCB	4
06410	*		
06420	*		
06430	*		
06440	*		
06450	*		
06460	*		
06470	*		
06480	*		
06490	*		
06500	*		
06510	*		
06520	*		
06530	*		
06540 23DB FE 1CA7	HXASC	LDX	HXBUF
06550 23DE 36		PSH A	SAVE DTA ON STK
06560 23DF 44		LSR A	SHIFT TOP 4 BITS TO DV 4
06570 23E0 44		LSR A	
06580 23E1 44		LSR A	

```

05510      *
05520      *
05530      *FUNCTION : PRINT
05540      *INPUTS (REG) : NONE
05550      *OUTPUTS (REG) : NONE
05560      *CALLS : OUTPUT
05570      *DESTROYS (REG) : A,B,X,CC
05580      *PURPOSE : TO PRINT AMPLITUDE DISTRIBUTION
05590      * DATA COMPILED DURING DATA COLLECTION.
05600      *
05610 22C2 CE 235A PDFPRT LDX #PDFNAM PUT CUR MEM FILE NAME IN MSG
05620 22C5 FF 1C9B STX NA1PTR
05630 22C8 CE 3402 LDX #HDRSTR+2
05640 22CB BD 20EF JSR PUTNAM
05650 22CF CE 234A LDX #PDFMS1 LOAD MSG POINTER
05660 22D1 86 58 LDA A #'X
05670 22D3 B7 236E STA A PDFMS2 STOR ASCII X TO MSG
05680 22D6 FD C4A7 JSR OUTPUT "CHANNEL X AMPLITUDE .."
05690 22D9 CE 3500 LDX #XPDF POINT TO X PDF DTA
05700 22DC 86 7F LDA A #$7F INIT COUNTER
05710 22DE 8D 25 BSR PDFPR1 PRINT X PDF DATA
05720 22E0 CE 234A LDX #PDFMS1 STOR ASCII Y TO MSG
05730 22E3 86 59 LDA A #'Y
05740 22E5 B7 236E STA A PDFMS2
05750 22E8 BD CA87 JSR OUTPUT
05760 22EB CE 3700 LDX #YPDF POINT TO Y PDF DTA
05770 22EE 86 7F LDA A #$7F
05780 22F0 8D 13 BSR PDFPR1
05790 22F2 CE 234A LDX #PDFMS1
05800 22F5 86 5A LDA A #'Z
05810 22F7 B7 236E STA A PDFMS2
05820 22FA BD CA87 JSR OUTPUT
05830 22FD CE 3900 LDX #ZPDF
05840 2300 86 7F LDA A #$7F
05850 2302 8D 01 ESR PDFPR1
05860 2304 39 RTS
05870      *
05880 2305 4C PDFPRI INC A
05890 2306 81 7F CMP A #$7F 256 VLU'S PRINT'D YET?
05900 2308 27 07 BEQ PDFPR2 YES. EXIT
05910 230A FD 231B JSR PRTVLU NO. PRINT VLU TO CONSOLE
05920 230D 08 INX
05930 230E 08 INX
05940 230F 20 F4 BRA PDFPRI PRINT NEXT VLU
05950 2311 FD 231B PDFPR2 JSR PRTVLU
05960 2314 CE 23D8 LDX #PDFMS6
05970 2317 FD CA8F JSR OUTNCR
05980 231A 39 RTS
05990      *
06000 231B 36 PRTVLU PSH A SAVE VALUE INDEX IN A
06010 231C DF 19 STX SAVEX SAVE LOC FOR HX TO BY X
06020 231E CE 23BF LDX #PDFMS4 PICK UP STRING VLU
06030 2321 FF 1CA7 STX HXBUF
06040 2324 FD 23DB JSR HXASC CONV VLU INDEX TO ASCII & STR

```

06590 23E2 44	LSR A		
06600 23E3 8D 08	RSR HXTOAS	CONV SHIFT'D BITS TO ASC	
06610 23E5 A7 00	STA A 0,X	STR IN ALPH STRING	
06620 23E7 32	PUL A		
06630 23E8 8D 03	BSR HXTOAS	CONV LOW 4 TO ASC	
06640 23EA A7 01	STA A 1,X	STR IN ALPH STRING	
06650 23EC 39	RTS		
06660	*		
06670 23ED 84 0F	HXTOAS AND A #\$0F	CLR TOP 4 BITS	
06680 23EF 81 0A	CMP A #\$0A	IS NUM GE HEX A?	
06690 23F1 2C 03	EJE HXTOAL	YES. BRANCH & ADD BIAS OF 37	
06700 23F3 8B 30	ADD A #\$30	NO. ADD BIAS OF 30 & RET	
06710 23F5 39	RTS		
06720 23F6 8B 37	HXTOAL ADD A #\$37		
06730 23F8 39	RTS		
06740	*		
06750	* END HXASC		
06760	*		
06770	* END OF EKG-EXEC ROUTINES		
06780	*		
06790	END		

```

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00470
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00490
00500
00510 0100
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\*

\*      DISPLAY NAME      :DISPLAY

\*      AUTHOR            :JOHN H. SCHNEIDER

\*      VERSION           :1.5

\*      VERSION DATE     :4 NOV 80

\*

\*\*\*\*\*

\*

\*                        PROGRAM DESCRIPTION

\*

\*   THIS PROGRAM DISPLAYS THE STATISTICAL DATA

\*   CALCULATED DURING THE COLLECTION OF AN EKG

\*   RECORDING. THE USER HAS SEVERAL MODES OF DATA

\*   OUTPUT AND THEY ARE LISTED IN THE COMMAND STRING

\*   BELOW. THIS PROGRAM ALSO PRODUCES PROBABILITY

\*   DENSITY FUNCTION (PDF) PLOTS OF THE DATA COLLECTED

\*   ON CHANNEL X,Y,OR Z VIA D/A CONVERTERS. THE DATA

\*   (PDF) IS DISPLAYED ON AND OSCILLISCOPE.

\*   IN ADDITION, THE COMMAND STRING ALLOWS RECONSTRUCT

\*   AND DISPLAY OF PREVIOUSLY COLLECTED DATA ON THE

\*   OSCILLISCOPE. THE RECONSTRUCTION ROUTINES ARE

\*   OVERLAYED OVER THE PRINT STATISTICS (PRSTAT)

\*   ROUTINES.

\*

\*                        COMMAND OPTIONS

\*

\*   0= RETURN TO EKG-EXEC

\*   1= PRINT CURRENT MEMORY FILE STATISTICS

\*   2= PRINT PDF AMPLITUDE TABLES

\*   3= DISPLAY XPDF ON THE OSCILLISCOPE

\*   4= DISPLAY YPDF ON THE OSCILLISCOPE

\*   5= DISPLAY ZPDF ON THE OSCILLISCOPE

\*   6= RECONSTRUCT & DISPLAY CH X ON

\*        THE OSCILLISCOPE

\*   7= RECONSTRUCT & DISPLAY CH Y ON

\*        THE OSCILLISCOPE

\*   8= RECONSTRUCT & DISPLAY CH Z ON

\*        THE OSCILLISCOPE

\*   9= LOAD MEMORY FILE FRG1 DISK

\*

\*\*\*\*\*

\*\*\*\*\*

\*

\*                        START OF DISPL

\*

\*\*\*\*\*

\*

ORG    \$0100    PROGRAM ORIGIN

OPT    O        ASSB OPT. LIST ASSEMBLY

OPT    NOG      ASSB OPT. SUPPRESS FCC LIST

\*

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00570	*		
00580	*	LABEL DESCRIPTIONS	
	*		
00600	*	* SUPPORT FOR DISPLA-5	
00610	*		
00620	CA8F	OUTNCR EQU	SCA8F EPRO/DOS. OUTPUT STRING WITH
00630	CA36	KEYRDO LQU	SCA36 EPRO/DOS. INPUT ALPH STRING
00640	1D00	START EQU	\$1D00 EKG-EXEC. START ADDRESS
00650	CCA0	MAKNAA EQU	SCCA0 EPRO/DOS. GET FILE NAME FROM
00660	C803	RLIB EQU	\$C803 EPRO/DOS. SET UP READ I/O
00670	2800	DOS EQU	\$2800 DOS. RUSTART LOC
00680	C43C	READ0 EQU	\$C43C EPRO/DOS. READ DISK ROUTINE
00690	F055	BYTE EQU	\$E055 MIKBUG. INPUT 2 HEX DIG FMT T
00700	C7C8	DRIVE EQU	\$C7C8 EPRO/DOSIO. READ DISK ROUTINE
00710	0500	PRSTAT EQU	\$0500 PRINT COLLEC STAT ROUTINE
00720	*		
00730	*	* DATA BUFFERS	
00740	*		
00741	0020	CHNLBF EQU	\$0020 CHANNEL DESCRIPTOR FOR DECPRS
00750	3400	HDRSTR EQU	\$3400 MEM FILE HDR ADDR
00760	3500	XPDF EQU	\$3500 LOC OF CH X PDF BUFFER
00770	3700	YPDF EQU	\$3700 LOC OF CH Y PDF BUFFER
00780	3900	ZPDF EQU	\$3900 LOC OF CH Z PDF BUFFER
00790	FFP8	IROVEC FQU	\$FFF8 LOC OF INT VECTOR ADDRESS
00800	1C9D	VTCsav EQU	\$1C9D TEMP IROVEC SAVE LOC
00810	3002	ENDBUF EQU	\$3002 BUF WITH ADDR OF LAST CHAR IN
00820	1C96	STKSAV EQU	\$1C96 TEMP STACK STORAGE BUFFER
00830	0019	SAVEX EQU	\$0019 TEMP ENDPX REG STORAGE BUFFER
00840	000A	EMETH EQU	\$000A END MEM ADDR FOR DISK I/O
00850	001B	PROGX FQU	\$001B ERR VFC ADDR
00860	1C9F	OLAYCO EQU	\$1C9F OVRLAY MODE FLAG
00870	3057	TYPE EQU	\$3057 FILE TYPE BUFFER
00880	1CA9	PDFPLC EQU	\$1CA9 PDFPRT ADDR PASS BUFFER
00890	1CAB	OVRLLC EQU	\$1CAB OVRLAY ADDR PASS BUFFER
00900	1CAD	OVRLBUF EQU	\$1CAD OVRLAY PARAMETER BUFFER
00910	*		
00920	*	* HARDWARE ADDRESSES	
00930	*		
00940	E500	DACZRO EQU	\$E500 D/A CONVERTER ZERO
00950	E400	ADCZRO EQU	\$E400 A/D CONVERTER ZERO
00960	*		
00970	*		
00980	*FUNCTION	:	DISPL
00990	*INPUTS (REG)	:	NONE
01000	*OUTPUTS (REG)	:	None
01010	*CALLS	:	OUTNCR,KEYRDO
01020	*DESTROYS	:	A,B,CC,X
01030	*PURPOSE	:	DISPL IS THE COMMAND EXEC PRG. IT REPUTS COMMAND INPUT FROM THE TERMINAL AND T DIRECTING EXECUTION OF THE SELECTED OI
01040	*		ONS.
01050	*		
01060	*		
01070	*		
01080	*		
01090	0100 CE 0100	DISPL LDY	#DISPL

## DISPLA-5

01100 0103 FF 2801	STX	DOS+1		
01110 0106 CE 025B	LDX	#DISPMS		
01120 010D FF C1E8	JSR	DISPMS	"DISPMS CALLING ...."	
01130 010C BD C136	JSR	DISPMS	GET LAST CHAR INPUT	
01140 010F FE 3002	LDX	ENDBUF	SET UP IN DATA BUF POINTER	
01150 0112 09	DEX			
01160 0113 E6 00	LDA B	0,X	GET LAST CHAR INPUT FROM TERM	
01170 0115 C1 30	CMP B	#"0	WAS IT 0?	
01180 0117 26 03	BNE	DISPL1	NO. CHEK IF 1	
01190 0119 7E 1D00	JMP	START	YES. RET TO EKG-EXEC	
01200 011C C1 31	DISPL1	CMP B	#"1	WAS COMMAND 1?
01210 011E 26 0C	BNE	DISPL2	NO. CHEK IF 2	
01220 0120 CE 01E4	LDX	#PRSNAM	GET "PRSTAT*" TO SEND TO OVRL	
01230 0123 FF 1CAD	STX	OVRBUF	STORE "PRSTAT" TO OVRLAY BUFF	
01240 0126 7F 1C9F	CLR	OLAYGO	RESTE OLAYGO FLAG	
01250 0129 7E 04D6	JMP	OVRLAY		
01260 012C C1 32	DISPL2	CIP B	#"2	WAS COMMAND 2?
01270 012E 26 0F	BNE	DISPL3	NO. CHEK IF 3	
01280 0130 FD 04DB	JSR	PDFPRT	PRINT PDF AMPLITUDE TABLES	
01290 0133 CE 0463	LDX	#DISMS1		
01300 0136 FD CA8F	JSR	OUTNCR	"PRESS RETURN"	
01310 0139 BD CA36	JSR	KEYBD0		
01320 013C 7E 0100	JMP	DISPL		
01330 013F C1 33	DISPL3	CMP B	#"3	WAS COMMAND 3?
01340 0141 26 0F	BNE	DISPL4	NO. CHECK 4	
01350 0143 CE 3500	LDX	#XPDF	GET ADDRESS OF X PDF DATA	
01360 0146 BD 0470	JSR	MAXSCN	SCAN PDF & FIND SHIFT NECESSA	
01370 *			TO ALLOW OUTPUT OF PDF VIA 12	
01380 0149 BD 04A0	JSR	PDFTRN	TRANSFER X PDF TO WORK BUFFER	
01390 014C BD 04BA	JSR	SCALE	SCALE WORK BUFFER BY SHIFT FR	
01400 014F 7E 01F4	JMP	PDFOUT	NOW DISPLAY X PDF TO OSCOPE.	
01410 0152 C1 34	DISPL4	CIP B	#"4	WAS COMMAND 4?
01420 0154 26 0F	BNE	DISPL5	NO. CHECK 5	
01430 0156 CE 3700	LDX	#YPDF	GET ADDRESS OF Y PDF DATA	
01440 0159 BD 0470	JSR	MAXSCN		
01450 015C BD 04A0	JSR	PDFTRN		
01460 015F BD 04BA	JSR	SCALE		
01470 0162 7E 01F4	JMP	PDFOUT		
01480 0165 C1 35	DISPL5	CMP B	#"5	WAS COMMAND 5?
01490 0167 26 0F	BNE	DISPL6	NO. RET & PRINT COMMAND PRGMP	
01500 0169 CE 3900	LDX	#ZPDF		
01510 016C BD 0470	JSR	MAXSCN		
01520 016F BD 04A0	JSR	PDFTRN		
01530 0172 PD 04BA	JSR	SCALE		
01540 0175 7E 01F4	JMP	PDFOUT		
01550 0178 C1 36	DISPL6	CIP B	#"6	
01560 017A 26 0F	BNE	DISPL7		
01561 017C 7F 0020	CLR	CHNLBF	SET CHNL DESIGNATOR FOR CH X	
01570 017F 7F 1C9F	CLR	OLAYGO		
01580 0182 CE 01EC	LDX	#DECPRS	LOAD & RUN DECPRS ROUTINE	
01590 0185 FF 1CAD	STX	OVRBUF	STORE "DECPRS" TO OVRLAY BUFF	
01600 0188 7E 04D6	JMP	OVRPLAY		
01610 018B C1 37	DISPL7	CMP B	#"7	
01620 018D 26 10	BNE	DISPL8		

## DISPLA-5

01621 018F 86 01	LDA A #1	SET CHNL DESIGNATOR FOR CH Y
01622 0191 97 20	STA A CHNLBF	
01623 0193 7F 1C9F	CLR CLAYCO	
01624 0194 7E 011C	LDI	LOAD & RUN DECPRS ROUTINE
01650 0199 FF 1CAD	STX OVREBUF	STORE "DECPRS" TO OVRLAY BUFF
01660 019C 7E 04D6	JMP OVRLAY	
01670 019F C1 38	DISPL8 CMP B #'8	
01680 01A1 26 10	RNE DISPL9	
01681 01A3 86 02	LDA A #2	SET CHNL DESIGNATOR FOR CH Z
01682 01A5 97 20	STA A CHNLBF	
01690 01A7 7F 1C9F	CLR OLAYCO	
01700 01AA CE 01EC	LDX #DECPRS	LOAD & RUN DECPRS ROUTINE
01710 01AD FF 1CAD	STX OVREBUF	STORE "DECPRS" TO OVRLAY BUFF
01720 01E0 7E 04D6	JMP OVRLAY	
01730 01B3 C1 39	DISPL9 CMP B #'9	
01740 01B5 26 2A	BNE CMDERR	
01750 01E7 86 22	LDA A #\$22	SET UP FILE TYPE FOR LOAD
01760 01B9 B7 3057	STA A TYPE	
01770 01BC CE 024C	LDX #NAMEPG	"ENTER FILENAME?"
01780 01RF BD CA&F	JSR OUTNCR	
01790 01C2 BD CC&0	JSR MAKIAA	
01800 01C5 BD C803	JSR RLIB	
01810 01C8 CE FFFF	LDX #\$FFFF	
01820 01CB DF 0A	STX EMEMI	
01830 01CD CE 2800	LDX #DOS	
01840 01D0 DF 1B	STX PROGX	
01850 01D2 CE 3400	LDX #HDRSTR	PUT ADDR TO LOAD IN INDEX
01860 01D5 ED C43C	JSR READ0	
01870 01D8 CE 0456	LDX #FILE1S	
01880 01DB BD CA&F	JSR OUTNCR	
01890 01DE BD CA36	JSR KEYBD0	
01900 01E1 7E 0100	CMDERR JMP	DISPL
01910 *		
01920 01E4 50	PRSNAM FCC	/PRSTAT* /
01930 01EC 44	DECFRS FCC	/DECPRS* /
01960 *		
01970 01F4 86 AA	PDFOUT LDA A #SAA	SET INTERRUPT DONE TEST FLAG
01980 01F6 B7 0248	STA A DONTST	
01990 01F9 0E	PDFOU1 CLI	CLEAR INTERRUPT MASK FOR STOP
02000 01FA 7F 024B	CLR KOUNT	CLR BUFFER PTR COUNTER
02010 01FD FE FFF8	LDX IROVLC	GET CUR IRQ VECTOR
02020 0200 FF 1C9D	STX VEC\$AV	SAVF IN TEMP BUFFER
02030 0203 CE 0240	LDX #STOP	LOAD ADDRESS OF STOPPING INTR
02040 0206 FF FFF8	STX IROVLC	
02050 0209 CE 07FF	LDX #\$07FF	SEND MAX POS PULSE FOR SCOPE
02060 020C FF E500	STX DACZRO	
02070 020F CE 3200	LDX #\$3200	GET START ADDR OF BUFFER PTR
02080 0212 7C 024B	PDFCU2 INC KOUNT	BUMP UP PTR COUNT
02090 0215 27 0E	BFO PDFOU3	HAS WHOLE BUFFER FLAG SET
02100 0217 A6 00	LDA A 0,X	NO. GET USB BYTE OF BUFFER PTR
02110 0219 E6 01	LDA B 1,X	GET LSB
02120 021B B7 1500	STA A DACZRO	SIND TO D/A CONV
02130 021E F7 E501	STA B DACZRO+1	SIND TO D/A CONV
02140 0221 08	INX	INC INDEXED ADDRESS

## DISPLA-5

02150 0222 08		IMX		
02160 0223 20 ED	BRA	PDFOU2	LOOP AGAIN	
02170 0224 00 FF E5C0	*	PDFOU1	PUT BACK IN INT BUFFER	
02180 *		PDFOU4	YES. EXIT DISPLAY LOOP	
02190 0228 27 02	REQ	PDFOU4	NO. KEEP SENDING DISPLAY	
02200 022A 20 CD	BRA	PDFOU1		
02210 022C 0F	PDFOU4	SEI		
02220 022D CE 4000		LDX	#\$4000	
02230 0230 FF E5C0		STX	DAC7RO	
02240 0233 B7 E400		STA A	ADCZRO	
02250 0236 01		NOP	RESET ST6800 INT FLIP FLOP	
02260 0237 FE 1C9D		LDX	VECSAV	
02270 023A FF FFF8		STX	GET ORG INT VTC ADDRESS	
02280 023D 7E 0100		JMP	IRQVFC	
02290 *			PUT BACK IN INT VTC BUFFER	
02300 0240 7F 0248 STOP	CLR	DISPL	RET & PRINT COMMAND PROMPT	
02310 0243 B7 E400	STA A	DONTST	CLEAR DONE TEST FLAG	
02320 0246 01		ADCZRO	CLREA ST6800 INT FLIP FLOP	
02330 0247 3B		FTI		
02340 0248 0001	DONTST RIB	1	RETURN FROM INTERRUPT AND SET	
02350 0249 0001	SHPCNT RIB	1	INTR DONE TEST FLAG	
02360 024A 0001	MAXCNT RIB	1	BUF WITH NUM OF SHF PCH 12 BI	
02370 024B 0001	COUNT RIB	1	TEMP BUF TO COUNT 256 THRU NO	
02380 *			TEMP COUNTER	
02390 024C 45	NAMMSG FCC	/ENTER /		
02400 0252 04	FCB	4		
02410 0253 44	DRVMSG FCC	/DRIVE? /		
02420 025A 04	FCB	4		
02430 *				
02440 025B 1A07	DISPMIS FDB	\$1A07,\$0C0D,\$0A0A,\$0A0A,\$0A0A,\$0A0A		
02450 0267 0D0A	FDB	\$0D0A		
02460 0269 44	FCC	/DATA DISPLAY/		
02470 0275 0D0A	FDB	\$0D0A,\$0D0A		
02480 0279 43	FCC	/COMMAND OPTIONS:/		
02490 0289 0D0A	FDB	\$0D0A,\$0D0A		
02500 028D 20	FCC	/ 0=RETURN TO ENG-EXEC/		
02510 02A5 0D0A	FDB	\$0D0A,\$0D0A		
02520 02A6 20	FCC	/ 1=PRINT CURRENT MEMORY FILE STA		
02530 02D5 0D0A	FDB	\$0D0A,\$0D0A		
02540 02D9 20	FCC	/ 2=PRINT PDF AMPLITUDE TABLES/		
02550 02FA 0D0A	FDB	\$0D0A,\$0D0A		
02560 02FE 20	FCC	/ 3=DISPLAY CH X PDF ON OSCILLISC		
02570 0325 0D0A	FDB	\$0D0A,\$0D0A		
02580 0329 20	FCC	/ 4=DISPLAY CH Y PDF ON OSCILLISC		
02590 0350 0D0A	FDB	\$0D0A,\$0D0A		
02600 0354 20	FCC	/ 5=DISPLAY CH Z PDF ON OSCILLISC		
02610 037B 0D0A	FDB	\$0D0A,\$0D0A		
02620 037F 20	FCC	/ 6=RECONSTRUCT & DISPLAY CH X ON		
02630 03A4 4F	FCC	/OSCILLISCOPE/		
02640 0370 0D0A	FDB	\$0D0A,\$0D0A		
02650 03F4 20	FCC	/ 7=RECONSTRUCT & DISPLAY CH Y ON		
02660 03D9 4F	FCC	/OSCILLISCOPE/		
02670 03E5 0D0A	FDB	\$0D0A,\$0D0A		
02680 03E9 20	FCC	/ 8=RECONSTRUCT & DISPLAY CH Z ON		

02690 040E 47	FCC	/ASCII LISCPY/
02700 041A 0D0A	FDB	\$0D0A,\$0D0A
02710 041B 20	FCC	/LOAD MEMORY FILE FROM DISK/
02720 0431 0A	LIN	LINE , 10
02730 0443 45	FCC	/ENTER COMMAND NOW=/
02740 0455 04	FCB	4
02750 0456 46	FILEMS	FCC /FILE LOADED. /
02760 0463 50	DISMS1	FCC /PRESS RETURN/
02770 046F 04	FCB	4
02780	*	
02790	*	
02800	*FUNCTION	:MAXSCN
02810	*INPUTS (REG)	:X
02820	*OUTPUTS	:SHIFT REQUIRED IN SHFCNT
02830	*CALLS	:NOTHING
02840	*DESTROYS	:A,B,CC
02850	*PURPOSE	:THIS ROUTINE SCANS THE PDF BUFFER PO TO BY X AND RETURNS THE MINIMUM SHIFT RIGHTS NECESSARY TO BRING THE AMPLITUDE COUNTS DOWN TO 11 BITS MAG FOR OUTPUT FROM THE ST6200 D/A CONVE
02860	*	
02870	*	
02880	*	
02890	*	
02900	*	
02910	*	
02920 0470 DF 19	MAXSCN	STX SAVEX SAVE PDF ADDR IN X
02930 0472 09		DEX ADJUST X FOR PROPER INDEXED L
02940 0473 09		DEX
02950 0474 7F 0249		CLR SHFCNT
02960 0477 7F 024A		CLR MAXCNT
02970 047A 7C 024A	MAXSC1	INC MAXCNT 256 BYTES SCANNED YET?
02980 047D 27 1E		BEQ MAXSC4 YES. EXIT WITH CORRECT SHIFT
02990 047F 08		INX NO. INC X & LOOK AT NEXT 2 BY
03000 0480 08		INX
03010 0481 A6 00		LDA A 0,X GET MSB BYTE OF VALUE
03020 0483 81 08		CMP A #08 IS IT LESS THAN 8?
03030 0485 2D F3		BIT MAXSC1 YES. NO SHIFT NECESS, GO CHEC
03040 0487 5F		CLR B NO. SET UP COUNTER
03050 0488 0C		CLC INSURE CARRY CLEAR
03060 0489 5C	MAXSC2	INC B NOW SHIFT UNTIL MSB 1 BIT HIT
03070 048A 48		ASL A
03080 048B 25 02		BCS MAXSC3 CARRY SET?
03090 048D 20 FA		BRA MAXSC2 NO. KEEP SHIFTING LEFT
03100 048F 17	MAXSC3	TBA YES. PUT SHIFT COUNT IN A
03110 0490 86 06		LDA A #6 GET MAX SHIFT
03120 0492 10		SRA AND GET DIFFERENCE
03130 0493 B1 0249		CMP A SHFCNT IS SHIFT LE CURRENT SHIFT COU
03140 0496 2F E2		BLE MAXSC1 YES. IGNORE THIS VALUE & CONT
03150 0498 B7 0249		STA A SHFCNT NO. SAVE THIS SHIFT COUNT
03160 049B 20 DD		BRA MAXSC1 NOW GO CHECK NEXT IF: WHILE
03170 049D DE 19	MAXSC4	LDX SAVEX RETRIEVE INPUT X
03180 049F 39		RTS
03190	*	
03200	*	
03210	*FUNCTION	:PDFTRN
03220	*INPUTS	:X

## DISPLA-5

03230	*OUTPUTS	:NONE
03240	*CALLS	:NOTHING
03250	*DESTROYS	:A,CC
03260	*PURPOSE	:THIS ROUTINE TRANSFERS THE PDF DATA POINTED TO BY X INTO A WORK BUFFER AREA AT 3200-33FF.
03270	*	
03280	*	
03290	*	
03300	*	
03310 04A0 DF 19	PDFTRN STX	SAVEX SAVE INDEX IN TEMP BUFFER
03320 04A2 BF 1C96	STS	STKSAV SAVE STACK POINTER
03330 04A5 35	TXS	SET UP STACK FOR BUFF TRANSF
03340 04A6 CE 3200	LDX #S3200	POINT TO WORK BUFF START
03350 04A9 8C 3400	PDFTR1 CPX #\$3400	DONE TRANSFERRING YET?
03360 04AC 27 06	BEQ PDFTR2	YES. RET
03370 04AE 32	PUL A	NO. GET NEXT BYTE
03380 04AF A7 00	STA A 0,X	STORE IN WORK BUFFER
03390 04B1 08	INX	
03400 04B2 20 F5	BRA PDFTR1	GO GET NEXT BYTE TO TRANSFER
03410 04B4 DE 19	PDFTR2 LDX SAVEX	RETRIEVE ENTRY X
03420 04B6 BE 1C96	LDS STKSAV	RETRIEVE STACK PTR
03430 04B9 39	RTS	
03440	*	
03450	*	
03460	*FUNCTION	:SCALE
03470	*INPUTS	:DATA IN WORK BUFFER
03480	*OUTPUTS	:DATA IN WORK BUFFER SCALED BY SHFCNT
03490	*CALLS	:NOTHING
03500	*DESTROYS	:A,B,CC
03510	*PURPOSE	:THIS ROUTINE SCALPS THE PDF DATA IN THE WORK BUFFER (3200-33FF) BY THE AMOUNT CALCULATED BY MAXCN TO ENABL THE ST6800 D/A CONVERTER TO OUTPUT A POSITIVE VOLTAGE PROPORTIONAL TO THE FREQUENCY OF OCCURENCE OF EACH OF THE POSSIBLE 256 AMPLITUDE VALUES
03520	*	
03530	*	
03540	*	
03550	*	
03560	*	
03570	*	
03580	*	
03590	*	
03600 04RA DF 19	SCALE STX	SAVEX SAVE ENTRY X
03610 04BC CE 3200	LDX #\$3200	LOAD WORK BUFFER START ADDRES
03620 04BF 5F	CLR B	
03630 04C0 BC 0249	SCALE1 LDA A	SHFCNT PICK UP SCALE SHIFT COUNT
03640 04C3 27 0E	BEQ SCALE4	IF SHFCNT = 0, NO SCALING NEC
03650 04C5 27 07	SCALE2 BEQ SCALE3	IS SHFCNT = 0?
03660 04C7 67 00	ASR 0,X	NO. SCALE DATA POINTED TO BY
03670 04C9 66 01	ROR 1,X	
03680 04CB 4A	DEC A	
03690 04CC 20 F7	BRA SCALE2	GO SHIFT DOWN AGAIN
03700 04CE 08	SCALE3 INX	INCREMENT WORK BUFFER DATA TO
03710 04CF 08	INX	
03720 04D0 5C	INC B	
03730 04D1 26 ED	RNE SCALE1	256 2 BYTE WORDS SHIFTED YET?
03740 04D3 DE 19	SCALE4 IDX SAVEX	NO. KEEP LOADING & SHIFTING
03750 04D5 39	RTS	RETRIEVE ENTRY Y
03760	*	

03770 \*FUNCTION :OVERLAY  
03780 \*INPUTS :X  
03790 \*OUTPUTS :NONE  
03800 \*CALLS :EXECUTIVE ROUTINES TO SET X  
03810 \*DESTROYS :ALL REG  
03820 \*PURPOSE :THIS IS A VECTOR ROUTINE TO CALL OVERLAY  
03830 \* EKG-EXEC TO LOAD ROUTINES FOR EXECUTION.  
03840 \*  
03850 04D6 FE 1CAB OVERLAY LDX OVRLLC  
03860 04D9 6E 00 JMP 0,X  
03870 \*  
03880 04DB FE 1CA9 PDFPRT LDX PDFPLC  
03890 04DE 6E 00 JMP 0,X  
03900 \*  
03910 \*\*\*\*\*  
03920 \*  
03930 \* END OF OVERLAY DISPLAY ROUTINES  
03940 \*  
03950 \*\*\*\*\*  
03960 \*  
03970 END

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00040      *
00040      *
00040      * OVERLAY NAME : PRSTAT
00040      * MEMORY FILE TO CONSOLE
00070      * VERSION : 1:9
00080      * VERSION DATE : 2 OCT 80
00090      *
00100      *
00110      *
00120      * PROGRAM DESCRIPTION
00130      *
00140      * THIS ROUTINE, WHEN LOADED AND EXECUTED, TAKES
00150      * THE STATUS DATA CURRENTLY PRESENT IN THE
00160      * MEMORY FILE HEADER SECTOR AND PRINTS THIS DATA
00170      * TO THE CONSOLE. THIS ROUTINE CAN BE CALLED
00180      * IMMEDIATELY FOLLOWING A DATA COLLECTION OR
00190      * AFTER A PREVIOUSLY COLLECTED DATA FILE HAS BEEN
00200      * LOADED INTO MEMORY FROM DISK.
00210      *
00220      *
00230      *
00240      *
00250      * START OF PRSTAT
00260      *
00270      *
00280      *
00290 0500      ORG    $0500    OVERLAY START ADDRESS
00300      *
00310      OPT    O        ASSB OPT-GIN ORJ FILE
00320      OPT    NOG     ASSB OPT-SUPPRESS FCC LIST
00330      *
00340      *
00350      *
00360      * LABLE DESCRIPTIONS
00370      *
00380      * SUPPORT SUBROUTINE ADDRESSES
00390      *
00400  CA8F    OUTICR EQU    $CA8F    EPROMDOS. ALPH STRING TO CONS
00410  CA36    KEYFD0 EQU    $CA36    EPROTDOS. ALPH STRING IN FRM
00420  1D00    START   EQU    $1D00    EKG-EXEC. SOFT EKG-EXEC START
00430  2800    DOS     EQU    $2800    DOS ERROR VECTOR, USED FOR JM
00440      *
00450      * DATA BUFFERS
00460      *
00470  1CA0    LGOFLG EQU    $1CA0    LOAD VERSUS EXECUTE FLAG
00480  1CA5    HXASC1 EQU    $1CA5    HXASC ADDR PASS BUFFER
00490  1CA7    HXRUF  EQU    $1CA7    HXASC PARAMETER BUFFER
00500  1C96    STKSAV EQU    $1C96    TEMP STACK SAVE BUFFER
00510  3400    HDRSTR EQU    $3400    MEM FILE HEADER START
00520  3490    LOOPCT  EQU    $3490    TOTAL CAL LOOPS EXEC IN WAIT
00530  3494    SAMPLNO EQU    $3494    NUM OF SAMPLES TAKEN
00540  3496    LPCAL   EQU    $3496    NUM OF CAL LOOPS PER 1 INTR
00550  3497    MAXZ    EQU    $3497    MAX VIU IN CH Z
00560  3498    MAXZLO EQU    $3498    LOCATION OF MAX VIU IN CH Z

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00570	349A	MINZ EQU	\$349A	MIN VLU IN CH Z
00580	349B	MINZLO EQU	\$349B	LOCATION OF MIN VLU IN CH Z
00590	34A0	MINY EQU	\$34A0	MAX VLU IN CH Y
00600	34A1	MINYLO EQU	\$34A1	LOCATION OF MIN VLU IN CH Y
00610	34A2	MAXX EQU	\$34A2	MIN VLU IN CH X
00620	34A3	MAXY EQU	\$34A3	MAX VLU IN CH X
00630	34A4	MAXXLO EQU	\$34A4	LOCATION OF MAX VLU IN CH X
00640	34A6	MINX EQU	\$34A6	MIN VLU IN CH X
00650	34A7	MINXLO EQU	\$34A7	LOCATION OF MIN VLU IN CH X
00660	34A9	MBIT EQU	\$34A9	NUM OF BITS AVAILABLE FOR STO
00680	34AC	DTABIT EQU	\$34AC	NUM OF BITS USED TO STORE DATA
00690	34B0	XBITS EQU	\$34B0	NUM OF BITS USED TO STORE X DT
00700	34B3	YBITS EQU	\$34B3	NUM OF BITS USED TO STORE Y DT
00710	34B6	ZBITS EQU	\$34B6	NUM OF BITS USED TO STORE Z DT
00720	34B9	TBITS EQU	\$34B9	NUM OF BITS USED TO STORE TIME
00730	34BC	ACELCT EQU	\$34BC	# BITS FED TO VAR LIN CODUP
00740	34C2	BASEAV EQU	\$34C2	SAVE MEM LOC S3620 PG01 BASIC
00750	34C3	ENTRPT EQU	\$34C3	START OF ENTROPY BUFFER
00760	3460	MAXMIN EQU	\$3460	START OF MAX,MIN ASCII DATA
00770	*			
00780	*			
00790		*FUNCTION : PRSTAT		
00800		*INPUTS (REG) : NONE		
00810		*OUTPUTS (REG) : NONE		
00820		*CALLS : OUTPUT, IXASC		
00830		*DESTROYS (REG) : A,B,CC,X		
00840		*PURPOSE : TO LIST STATUS INFORMATION PG01 DATA		
00850		* COLLECTION FILE TO TERMINAL.		
00860	*			
00870	*			
00880	0500 86 08	PRSTAT LDA A #8		CLR MSG FIELDS IN STAT OUTPUT
00890	0502 CE 0A66	LDX #NAME		
00900	0505 BD 0930	JSR MSGCLR		
00910	0508 86 35	LDA A #53		
00920	050A CE 0A8C	LDX #SURJ		
00930	050D BD 0930	JSR MSGCLR		
00940	0510 86 0A	LDA A #10		
00950	0512 CE 0ADF	LDX #RATE		
00960	0515 PD 0930	JSR MSGCLR		
00970	0518 86 14	LDA A #20		
00980	051A CE 0B07	LDX #DATE		
00990	051D BD 0930	JSR MSGCLR		
01000	0520 86 08	LDA A #8		
01010	0522 CE 0B39	LDX #TIME		
01020	0525 BD 0930	JSR MSGCLR		
01030	0528 86 19	LDA A #25		
01040	052A CE 0B5F	LDX #TYPE		
01050	052D BD 0930	JSR MSGCLR		
01060	0530 86 35	LDA A #53		
01070	0532 CE 0F29	LDX #COMINT		
01080	0535 BD 0930	JSR MSGCLR		
01090	0538 86 08	LDA A #8		
01100	053A CE 0B96	LDX #XINT		

01110 053D RD 0930	JSR	MSGCLR
01111 053E RD 08	LDA A	#8
01112 053F RD 08	LDX	#FFFF
01140 0545 PD 0930	JSR	MSGCLR
01150 0548 86 08	LDA A	#8
01160 054A CE 0FFA	LDX	#ZENT
01170 054D RD 0930	JSR	MSGCLR
01180 0550 86 08	LDA A	#8
01190 0552 CE 0C2C	LDX	#TPNT
01200 0555 RD 0930	JSR	MSGCLR
01210 0558 86 08	LDA A	#8
01220 055A CE 0CCE	LDX	#MANCPR
01230 055D RD 0930	JSR	MSGCLR
01240 0560 86 08	LDA A	#8
01250 0562 CE 0D11	LDX	#CPTRACI
01260 0565 RD 0930	JSR	MSGCLR
01270 0568 86 05	LDA A	#5
01280 056A CE 0D17	LDX	#CPTRDP
01290 056D RD 0930	JSR	MSGCLR
01300 0570 86 05	LDA A	#5
01310 0572 CE 0D19	LDX	#TITTTF
01320 0575 PD 0930	JSR	MSGCLR
01330 0578 86 05	LDA A	#5
01340 057A CE 0F2B	LDX	#COLNUR
01350 057D RD 0930	JSR	MSGCLR
01360 0580 86 08	LDA A	#8
01370 0582 CE 0F2F	LDX	#AMINX
01380 0585 RD 0930	JSR	MSGCLR
01390 0588 86 08	LDA A	#8
01400 058A CE 0FFD	LDX	#AMAXX
01410 058D RD 0930	JSR	MSGCLR
01420 0590 86 08	LDA A	#8
01430 0592 CE 0E93	LDX	#AMINY
01440 0595 RD 0930	JSR	MSGCLR
01450 0598 86 08	LDA A	#8
01460 059A CE 0E61	LDX	#AMINY
01470 059D RD 0930	JSR	MSGCLR
01480 05A0 86 08	LDA A	#8
01490 05A2 CE 0E93	LDX	#AMINY
01500 05A5 RD 0930	JSR	MSGCLR
01510 05A8 86 08	LDA A	#8
01520 05AA CE 0F05	LDX	#AMANZ
01530 05AD RD 0930	JSR	MSGCLR
01540 05B0 86 08	LDA A	#8
01550 05B2 CE 0FF7	LDX	#AMINZ
01560 05B5 RD 0930	JSR	MSGCLR
01570 05B8 CE 3402	LDX	#AMINZ+2 NOW PUT LDA IN FULL HDP IN 11
01580 05BB FF 093D	STX	PICTMC
01590 05BE CE 0166	LDX	#EINT
01600 05C1 RD 093C	JSR	MSGDMT
01610 05C4 CE 0A8C	LDX	#EINT
01620 05C7 RD 093C	JSR	MSGDMT
01630 05CA CE 0A8C	LDX	#EINT
01640 05CD RD 093C	JSR	MSGDMT

01660 05D0 CE 0F07	LDX	#DATE
01660 05D3 BD 093C	JSR	MSGPUT
01670 05E5 CE 0139	LDX	#TIME
01680 05E5 CE 0139	JSR	MSGPUT
01690 05DC CE 0F29	LDX	#COMINT
01700 05DF BD 093C	JSR	MSGPUT
01710 05E2 FE 34C3	LDX	ENTRPY
01720 05E5 8C 2020	CPX	#\$2020
01730 05E8 27 03	BDQ	PRST00
01740 05EA 7E 06A4	JMP	PRSTA0
01750 05FD CE 09FD PRST00	LDX	#ENTCLR
01760 05F0 FF 095D	STX	FRMLOC
01770 05F3 CE 0B96	LDX	#XINT
01780 05F6 BD 093C	JSR	MSGPUT
01790 05F9 CE 09FD	LDX	#ENTCLR
01800 05FC FF 095D	STX	FRMLOC
01810 05FF CE 0FC8	LDX	#YINT
01820 0602 FD 093C	JSR	MSGPUT
01830 0605 CE 09FD	LDX	#ENTCLR
01840 0608 FF 095D	STX	FRMLOC
01850 060B CE 0BFA	LDX	#ZINT
01860 060E BD 093C	JSR	MSGPUT
01870 0611 CE C9FD	LDX	#ENTCLR
01880 0614 FF 095D	STX	FRMLOC
01890 0617 CE 0C2C	LDX	#TINT
01900 061A BD 093C	JSR	MSGPUT
01910 061D CE 09FD	LDX	#ENTCLR
01920 0620 FF 095D	STX	FRMLOC
01930 0623 CE 0CCE	LDX	#MAXCLR
01940 0626 BD 093C	JSR	MSGPUT
01950 0629 CE 09FD	LDX	#ENTCLR
01960 062C FF 095D	STX	FRMLOC
01970 062F CF 0D11	LDX	#CPRACH
01980 0632 BD 093C	JSR	MSGPUT
01990 0635 CE 09FD	LDX	#ENTCLR
02000 0638 FF 095D	STX	FRMLOC
02010 063B CE 0D57	LDX	#CPREFF
02020 063E BD 093C	JSR	MSGPUT
02030 0641 CE 09FD	LDX	#ENTCLR
02040 0644 FF 095D	STX	FRMLOC
02050 0647 CE 0D99	LDX	#TIMEFF
02060 064A BD 093C	JSR	MSGPUT
02070 064D CE 09FD	LDX	#ENTCLR
02080 0650 FF 095D	STX	FRMLOC
02090 0653 CE 0DCB	LDX	#COLNUR
02100 0656 FD 093C	JSR	MSGPUT
02110 0659 CE 09FD	LDX	#ENTCLR
02120 065C FF 095D	STX	FRMLOC
02130 065F CE 0DFF	LDX	#MAXX
02140 0662 FD 093C	JSR	MSGPUT
02150 0665 CE 09FD	LDX	#ENTCLR
02160 0668 FF 095D	STX	FRMLOC
02170 066B CE 0FCF	LDX	#AMINX
02180 066E BD 093C	JSR	MSGPUT

02190 0671 CE 09FD	LDX	#ENTCLR
02200 0674 FF 095D	STX	FRMLOC
02210 067A BD 093C	LDX	#ENTCLR
02220 067A BD 093C	JSR	#AMINX
02230 067D CE 09FD	LDX	#ENTCLR
02240 0680 FF 095D	STX	FRMLOC
02250 0683 CE 0993	LDX	#AMINY
02260 0686 BD 093C	JSR	MSGPUT
02270 0689 CE 091D	LDX	#ENTCLR
02280 068C FF 095D	STX	FRMLOC
02290 068F CW 0EC5	LDX	#AMAXZ
02300 0692 BD 093C	JSR	MSGPUT
02310 0695 CE 09FD	LDX	#ENTCLR
02320 0698 FF 095D	STX	FRMLOC
02330 069B CE 0EF7	LDX	#AMINZ
02340 069E BD 093C	JSR	MSGPUT
02350 06A1 7E 070A	JMP	PRSTAL
02360 06A4 CE 34C3 PRSTA0	LDX	#ENTRPY GET ADDRESS OF ENTRPY BUFFER
02370 06A7 FF 095D	STX	FRMLOC
02380 06AA CE 0EB6	LDX	#XENT NOW PRINT APPROPRIATE DAT TO
02390 06AD BD 093C	JSR	MSGPUT
02400 06B0 CE 0FC8	LDX	#YENT
02410 06B3 BD 093C	JSR	MSGPUT
02420 06B6 CE 0FFA	LDX	#ZENT
02430 06B9 BD 093C	JSR	MSGPUT
02440 06BC CE 0C2C	LDX	#TENT
02450 06BF BD 093C	JSR	MSGPUT
02460 06C2 CE 0CCE	LDX	#MAXCPR
02470 06C5 BD 093C	JSR	MSGPUT
02480 06C8 CE 0D11	LDX	#CPRACH
02490 06CB BD 093C	JSR	MSGPUT
02500 06CE CE 0D57	LDX	#CPRFFF
02510 06D1 BD 093C	JSR	MSGPUT
02520 06D4 CE 0D99	LDX	#TIMEFF
02530 06D7 PD 093C	JSR	MSGPUT
02540 06DA CE 0DCB	LDX	#COLDUR
02550 06DD BD 093C	JSR	MSGPUT
02560 06E0 CE 3460	LDX	#MAXHIN NOW GET ADDR WHERE ASCII MAX,
02570 06E3 FF 095D	STX	FRMLOC
02580 06E6 CE 0DFD	LDX	#AMAXX
02590 06E9 BD 093C	JSR	MSGPUT
02600 06EC CE 0E2F	LDX	#AMINX
02610 06FF BD 093C	JSR	MSGPUT
02620 06F2 CE 0E61	LDX	#AMAXY
02630 06F5 BD 093C	JSR	MSGPUT
02640 06F8 CE 0E93	LDX	#AMINY
02650 06FB BD 093C	JSR	MSGPUT
02660 06FE CE 0FC5	LDX	#AMAXZ
02670 0701 BD 093C	JSR	MSGPUT
02680 0704 CE 0EF7	LDX	#AMINZ
02690 0707 BD 093C	JSR	MSGPUT
02700 070A FE 3400 PRSTAL	LDX	HDRSIR NOW DECODE CONTRS TYPE & PUT
02710 070D 8C 4E43	CPX	#\$4E43 (NC) NOCPRS?
02720 0710 26 08	PNE	PRSTA2 NO

02700 0712 CE 0979	LDX	#NOCOMP	YES. POINT TO NOCOMP MSG	
02740 0715 FF 095D	STX	PRMLOC		
02750 0718 DD 42	PRA	PRSTA7		
02760 071A CE 0951	LDX	#PRSTA7	(TP) TURNING?	
02770 071D 26 08	LXI	PRSTA3	NO.	
02780 071F CE 09C2	LDX	#TOLA	YES. POINT TO TOLAN-A MSG	
02790 0722 FF 095D	STX	FRMLOC		
02800 0725 20 35	PRA	PRSTA7		
02810 0727 8C 5442	PRSTA3	CPX	#\$5442 (TB) TOLAN-B?	
02820 072A 26 08	ENE	PRSTA4	NO	
02830 072C CE 09CA	LDX	#TOLB	YES. POINT TO TOLAN-B MSG	
02840 072F FF 095D	STX	FRMLOC		
02850 0732 20 28	BRA	PRSTA7		
02860 0734 8C 444F	PRSTA4	CPX	#\$444F (DO) DOWER?	
02870 0737 26 08	BNE	PRSTA5	NO	
02880 0739 CE 09D2	LDX	#DOWR	YES. POINT TO DOWER MSG	
02890 073C FF 095D	STX	FRMLOC		
02900 073F 20 1B	PRA	PRSTA7		
02910 0741 8C 5450	PRSTA5	CPX	#\$5450 (TP) TURNING POINT?	
02920 0744 26 08	ENE	PRSTA6	NO.	
02930 0746 CE 09D8	LDX	#TURNPT	YES. POINT TO TURNING POINT M	
02940 0749 FF 095D	STX	FRMLOC		
02950 074C 20 0E	BRA	PRSTA7		
02960 074E 8C 494E	PRSTA6	CPX	#\$494E (IN) 2ND ORDER INTERPOLATOR?	
02970 0751 27 03	EEQ	PRSTA1	YES. TYPE IS RECOGNIZED. CONT	
02980 0753 7E 095F	JMP	ERROR	NO. TYPE NOT RECOGNIZED. EPR	
02990 0756 CE 09E6	PRSTA6	LDX	#INTER	
03000 0759 FF 095D	STX	FRMLOC		
03010 075C CE 0B5F	PRSTA7	LDX	#TYPE PUT CVERS TYPE IN MSG	
03020 075F BD 093C	JSR	MSGPUT		
03030 0762 CE 0A12	LDX	#HDRMSG	NOW PT TO STATMSG PAGE 1	
03040 0765 BD CA8F	JSR	OUTNCR	AND PRINT PAGE 1	
03050 0768 BD CA36	JSR	KEYDD0		
03060 076B CE 0C6B	LDX	#HDPMIS2	NOW PT TO STATMSG PAGE 2	
03070 076E BD CA8F	JSR	OUTNCR	AND PRINT PAGE 2	
03080 0771 BD CA36	JSR	KEYDD0		
03090 0774 CE 0F64	LDX	#NUMSMP+2	PICK UP NUM OF SMPLS,	
03100 0777 D6 3495	LDA A	SAMPNO+1	CONVRT TO ASCII, & STORE	
03110 077A BD 0955	JSR	HXASC	TO OUTPUT STRING.	
03120 077D CN 0F62	LDX	#NUMSMP		
03130 0780 EG 3494	LDA A	SAMPNO		
03140 0783 FD 0955	JSR	HXASC		
03150 0786 CE 1019	LDX	#NUMLPS	PICK UP VLU OF CAL TST	
03160 0789 EG 3496	LDA A	LPCAL	& STORE TO OUTPUT STRING	
03170 078C FD 0955	JSR	HXASC		
03180 078F CE 1060	LDX	#TOTLUP+6	PICK UP TOTL LOOP COUNT	
03190 0792 EG 3493	LDA A	LOOPCT+3	, CONVRT TO ASCII & STORE	
03200 0795 FD 0955	JSR	HXASC	TO OUTPUT STRING	
03210 0798 CE 105E	LDX	#TOTLUP+4		
03220 079B FD 3492	LDA A	LOOPCT+2		
03230 079E BD 0955	JSR	HXASC		
03240 07A1 CF 105C	LDX	#TOTLUP+2		
03250 07A4 FD 3491	LDA A	LOOPCT+1		
03260 07A7 BD 0955	JSR	HXASC		

03270 07AA CE 105A	LDX #TOTLUP
03270 07B0 BD 3490	LDA A LOOPCT
03270 07B1 BD 3490	JSR HXASC
03300 07B3 CE 10C5	LDX #XMAXNM PICK UP MAX VLU ON CH X
03310 07B6 B6 34A3	LDA A MAXX CON TO ASCII & STR IN
03320 07B9 BD 0955	JSR HXASC TO OUTPUT STRING
03330 07BC CE 10E1	LDX #XMAXLO+2
03340 07BF B6 34A5	LDA A MAXXLO+1
03350 07C2 BD 0955	JSR HXASC
03360 07C5 CF 10DF	LDX #XMINLO
03370 07C8 B6 34A4	LDA A MAXXLO
03380 07CB BD 0955	JSR HXASC
03390 07CE CE 10F2	LDX #XMINNM PICK UP MIN VLU ON CH X
03400 07D1 B6 34A6	LDA A MINX CONV TO ASCII & STORE IN
03410 07D4 BD 0955	JSR HXASC OUTPUT STRING
03420 07D7 CE 110E	LDX #XMINLO+2
03430 07DA B6 34A8	LDA A MINXLO+1
03440 07DD BD 0955	JSR HXASC
03450 07E0 CE 110C	LDX #YMAXLO
03460 07E3 B6 34A7	LDA A MINXLO
03470 07E6 BD 0955	JSR HXASC
03480 07E9 CE 111F	LDX #YMAXNM PICK UP MAX VLU ON CH Y
03490 07EC B6 349D	LDA A MAXY CONV TO ASCII & STORE IN
03500 07FF BD 0955	JSR HXASC OUTPUT STRING
03510 07F2 CE 113B	LDX #YMAXLO+2
03520 07F5 B6 349F	LDA A MAXYLO+1
03530 07F8 BD 0955	JSR HXASC
03540 07FB CE 1139	LDX #YMINLO
03550 07FE B6 349D	LDA A MAXYLO
03560 0801 BD 0955	JSR HXASC
03570 0804 CE 114C	LDX #YMINNM PICK UP MIN VLU ON CH Y
03580 0807 B6 34A0	LDA A MINY CONV TO ASCII & STORE IN
03590 080A BD 0955	JSR HXASC OUTPUT STRING
03600 080D CE 1168	LDX #YMINLO+2
03610 0810 B6 34A2	LDA A MINYLO+1
03620 0813 BD 0955	JSR HXASC
03630 0816 CE 1166	LDX #ZMAXLO
03640 0819 B6 34A1	LDA A MINYLO
03650 081C BD 0955	JSR HXASC
03660 081F CE 1179	LDX #ZMAXNM PICK UP MAX VLU ON CH Z
03670 0822 B6 3497	LDA A MAXZ CONV TO ASCII & STORE IN
03680 0825 BD 0955	JSR HXASC OUTPUT STRING
03690 0828 CE 1195	LDX #ZMAXLO+2
03700 082B B6 3499	LDA A MAXZLO+1
03710 082E BD 0955	JSR HXASC
03720 0831 CE 1193	LDX #ZMINLO
03730 0834 B6 3498	LDA A MAXZLO
03740 0837 BD 0955	JSR HXASC
03750 083A CE 11A6	LDX #ZMINNM PICK UP MIN VLU ON CH Z
03760 083D B6 349A	LDA A MINZ CONV TO ASCII & STORE IN
03770 0840 BD 0955	JSR HXASC OUTPUT STRING
03780 0843 CE 11C2	LDX #ZMINLO+2
03790 0846 B6 349C	LDA A MINZLO+1
03800 0849 BD 0955	JSR HXASC

03810 084C CE 11C0	LDX #ZMINIO
03820 084F B6 349B	LDA A MINZIO
03830 0850 BD 0955	JSR HXASC
03840 0850 CE 110B	LDX #AVABIT+4 PICK UP OF AVAILBITS
03850 0858 B6 34AB	LDA A MEMBIT+2 AVAILABLE FOR DATA STOR
03860 085B BD 0955	JSR HXASC , CONV TO ASCII & STOR
03870 085E CE 120C	LDX #AVABIT+2 IN OUTPUT STRING
03880 0861 B6 34AA	LDA A MEMBIT+1
03890 0864 BD 0955	JSR HXASC
03900 0867 CE 120A	LDX #AVABIT
03910 086A B6 34A9	LDA A MEMBIT
03920 086D BD 0955	JSR HXASC
03930 0870 CE 124B	LDX #TCPBIT+6
03940 0873 B6 34BF	LDA A ACELCT+3
03950 0876 BD 0955	JSR HXASC
03960 0879 CE 1249	LDX #TCPBIT+4
03970 087C B6 34BE	LDA A ACELCT+2
03980 087F BD 0955	JSR HXASC
03990 0882 CE 1247	LDX #TCPBIT+2
04000 0885 B6 34ED	LDA A ACELCT+1
04010 0888 BD 0955	JSR HXASC
04020 088B CE 1245	LDX #TCPBIT
04030 088E B6 34BC	LDA A ACELCT
04040 0891 BD 0955	JSR HXASC
04050 0894 CE 127F	LDX #TOTBIT+6 PICK UP NUM OF BITS
04060 0897 F6 34AF	LDA A DTABIT+3 ACTUALLY STORED
04070 089A BD 0955	JSR HXASC BECAUSE OF DATA
04080 089D CE 127D	LDX #TOTBIT+4 COMPRESSION, CONV TO
04090 08A0 F6 34AE	LDA A DTABIT+2 ASCII & STORE IN
04100 08A3 ED 0955	JSR HXASC OUTPUT STRING
04110 08A6 CE 127B	LDX #TOTBIT+2
04120 08A9 B6 34AD	LDA A DTABIT+1
04130 08AC BD 0955	JSR HXASC
04140 08AF CE 1279	LDX #TOTBIT
04150 08B2 F6 34AC	LDA A DTABIT
04160 08B5 BD 0955	JSR HXASC
04170 08B8 CE 12B7	LDX #XBITNM+4 PICK UP NUM OF BITS
04180 08BB B6 34B2	LDA A XBITS+2 ACTUALLY STORED FOR
04190 08BE BD 0955	JSR HXASC DATA IN CHANNEL X,
04200 08C1 CE 12B5	LDX #XBITNM+2 CONV TO ASCII &
04210 08C4 F6 34B1	LDA A XBITS+1 STORE IN OUTPUT STRING
04220 08C7 BD 0955	JSR HXASC
04230 08CA CE 12B3	LDX #XBITNM
04240 08CD F6 34B0	LDA A XBITS
04250 08D0 ED 0955	JSR HXASC
04260 08D3 CE 12FF	LDX #YBITNM+4 PICK UP NUM OF BITS
04270 08D6 B6 34B5	LDA A YBITS+2 ACTUALLY STORED FOR
04280 08D9 BD 0955	JSR HXASC DATA IN CHANNEL Y,
04290 08DC CE 12FD	LDX #YBITNM+2 CONV TO ASCII &
04300 08DF B6 34B4	LDA A YBITS+1 STORE IN OUTPUT STRING
04310 08E2 PD 0955	JSR HXASC
04320 08E5 CE 12FB	LDX #YBITNM
04330 08E8 F6 34B3	LDA A YBITS
04340 08EB BD 0955	JSR HXASC

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04350 08EE CE 1327 LDX #ZBITNM+4 PICK UP NUM OF BITS
04360 08F1 B6 34B8 LDA A ZBITS+2 ACTUALLY STORED FOR
04370 08F4 BD 0955 JSR HXASC DATA IN CHANNEL N,
04380 08F7 C1 1325 LDX #ZBITNM+2 CONV TO ASCII &
04390 08FA B6 34B7 LDA A ZBITS+1 STORE IN OUTPUT STRING
04400 08FD BD 0955 JSR HXASC
04410 0900 CE 1323 LDX #ZBITNM
04420 0903 B6 34B6 LDA A ZBITS
04430 0906 BD 0955 JSR HXASC
04440 0909 CE 136D LDX #TBITNM+4 PICK UP NUM OF BITS
04450 090C B6 34EB LDA A TBITS+2 ACTUALLY STORED FOR
04460 090F BD 0955 JSR HXASC TIME (OR OTHER PARAMETER)
04470 0912 CE 136B LDX #TBITNM+2 CONV TO ASCII &
04480 0915 B6 34BA LDA A TBITS+1 STORE IN OUTPUT STRING
04490 0918 BD 0955 JSR HXASC
04500 091B CE 1369 LDX #TBITNM
04510 091E B6 34B9 LDA A TBITS
04520 0921 BD 0955 JSR HXASC
04530 0924 CE 0F89 LDX #COSTAT GET CUTPUT STRING ADDR
04540 0927 BD CA8F JSR OUTNCR & SEND TO CONSOLE
04550 092A BD CA36 JSR KEYBD0 WAIT FOR TERM INPUT
04560 092D 7E 2800 JMP DOS JUMP BACK TO CALLING ROUTINE
04570 *
*FUNC: MSGCLR
*INPUTS: A (# OF SPACES TO CLR), X (LOC TO PUT SPACES
*OUTPUTS: ASCII $20 TO MEM AT X
*CALS: NOTHING
*DESTROYS: A,B,CC
*PURPOSE: THIS ROUTINE CLRS THE MSG BUFFR EACH TIME
* PRSTAT IS CALLED FOR NEW INFO OFF OF THE MEM DATA
*
04660 0930 4C MSGCLR INC A INC COUNTER
04670 0931 C6 20 LDA B #$20 ASCII SPACE
04680 0933 4A MSGCLL DEC A
04690 0934 27 05 BEQ MSGCLL2 IS BUFFER CLRRED?
04700 0936 E7 00 STA B 0,X NO. KEEP LOOPING
04710 0938 08 INX
04720 0939 20 F8 BRA MSGCLL
04730 093B 39 MSGCLL2 RTS
04740 *
04750 * END OF MSGCLR
04760 *
04770 ****
04780 *
04790 *FUNC: MSGPUT
04800 *INPUTS: X (ADDR WHERE DATA IN FRMLOC GOING)
04810 *OUTPUTS: ASCII DATA TO ADDR IN X
04820 *CALLS: NOTHING
04830 *DESTROYS: A,X,CC
04840 *PURPOSE: THIS ROUTINES TRANSFER ASCII TEXT
04850 * FROM MEM FILE HDR TO STAT MSG BUFFR
04860 *
04870 093C BF 1C96 MSGPUT STS STKSAV SAVE CURRENT STACK PTR
04880 093F BE 095D LDS FRMLOC GET LOC OF DATA FOR TRANSMITTER

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0480 0942 34      DES
04900 0943 32     MSGPUL PUL A
04910 0944 64     CMP A #4      END OF MSG?
04911 0945 27 05   INQ MSGPU2  YES, EXIT
04930 0948 A7 00   STA A 0,X    NO. PUT CHAR IN MSG
04940 094A 08      INX
04950 094B 20 F6   BRA MSGPUL
04960 094D 31      MSGPU2 INS   UPDATE STK WITH NEXT TRANSFR L
04970 094E BF 095D   STS FRMLOC STR BACK IN FRMLOC
04980 0951 BE 1C96   LDS STKSAV PICK UP ENTRY STACK PTR
04990 0954 39      RTS
05000 *
05010 * END MSGPUT
05020 *
05030 ****
05040 *
05050 *FUNC: HXASC
05060 *INPUTS:A,X (DATA TO BE CONV'D,ADDR TO STORE ASCII O
05070 *CALLS: HSASC ROUTINE IN EKG-EXEC
05080 *DESTROYS: A,X,CC
05090 *PURPOSE: THIS IS A RELOC PASS ROUTINE TO PICK UP
05100 * ADDRO OF HXASC IN EKG-EXEC AND JUMP TO IT
05110 *
05120 0955 FF 1CA7 HXASC STX HXBUF SAVE X IN PARAMETER BUFFER
05130 0958 FE 1CA5   LDX HXASLC GWT ADDR OF CONV ROUTINE IN E
05140 095B 6E 00   JMP 0,X   JUMP TO IT
05150 *
05160 * END OF HXASC
05170 *
05180 ****
05190 *
05200 095D 0002     FRMLOC RMB 2      DATA FROM LOC FOR MSGPUT
05210 *
05220 095F CE 096B ERROR LDX #ERRMSG
05230 0962 BD CA8F   JSR OUTNCR
05240 0965 BD CA36   JSR KEYDDO
05250 0968 7E 1D00   JMP START
05260 *
05270 096B 0707 ERRMSG FDB $0707,$0D0A
05280 096F 4D FCC /MEMORY FILE COMPRESSION TYPE /
05290 098C 4E FCC /NOT RECOGNIZED. PRESS RETURN/
05300 09A8 04 FCB 4
05310 *
05320 ****
05330 * END PRSTAT
05340 ****
05350 *
05360 ****
05370 *
05380 * OUTPUT STRING TO LIST TO CONSOLE
05390 *
05400 ****
05410 *
05420 09A9 4E NOCOMP FCC /NO COMPRESSION PERFORMED/

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05430 09C1 04		FCB	4
05440 09C2 54	TOLA	FCC	/TOLAN-A/
05450 09C9 04		FCB	4
05460 09CA 54	TOLB	FCC	/TOLAN-B/
05470 09D1 04		FCB	4
05480 09D2 44	DOWR	FCC	/DOWER/
05490 09D7 04		FCB	4
05500 09D8 54	TURNPT	FCC	/TURNING POINT/
05510 09E5 04		FCB	4
05520 09E6 32	INTER	FCC	/2ND ORDER INTERPOLATOR/
05530 09FC 04		FCB	4
05540 09FD 43	ENTCLR	FCC	/CALCULATION NOT MADE/
05550 0A11 04		FCB	4
05560 0A12 1A07	HDRMSG	FDB	\$1A07,\$0C0D,\$0A0A,\$0ACA,\$0ACA,\$0ACA
05570 0A1E 3A		FCC	/: EKG SAMPLE COLLECTION STATISTICS/
05580 0A3F 20		FCC	/ : PAGE 1/
05590 0A48 0D0A		FDB	\$0D0A,\$0D0A
05600 0A4C 46		FCC	/FILENAME. . . . . /
05610 0A66 0008	NAME	RMB	8
05620 0A6E 0D0A		FDB	\$0D0A,\$0D0A
05630 0A72 53		FCC	/SUBJECT . . . . . /
05640 0A8C 0035	SUBJ	RME	53
05650 0AC1 0D0A		FDB	\$0D0A,\$0D0A
05660 0AC5 53		FCC	/SAMPLING RATE . . . . . /
05670 0ADF 000A	RATE	RMB	10
05680 0AE9 0DCA		FDB	\$0D0A,\$0D0A
05690 0AED 44		FCC	/DATE OF COLLECTION. . . . /
05700 0B07 0014	DATE	RMB	20
05710 0B1B 0D0A		FDB	\$0D0A,\$0D0A
05720 0B1F 54		FCC	/TIME OF COLLECTION. . . . /
05730 0B39 0008	TIME	RMB	8
05740 0B41 CDOA		FDB	\$0D0A,\$0D0A
05750 0B45 43		FCC	/COMPRESSION USED. . . . /
05760 0B5F 0019	TYPE	RMB	25
05770 0B78 0D0A		FDB	\$0D0A,\$0D0A
05780 0B7C 43		FCC	/CHANNEL X ENTROPY . . . . /
05790 0B96 0008	XENT	RMB	8
05800 0BE9 20		FCC	/ BITS /
05810 0BAA 0D0A		FDB	\$0D0A,\$0D0A
05820 0BAE 43		FCC	/CHANNEL Y ENTROPY . . . . /
05830 0BC8 0008	YENT	RMB	8
05840 0BD0 20		FCC	/ BITS /
05850 0BDC 0DCA		FDB	\$0D0A,\$0D0A
05860 0BE0 43		FCC	/CHANNEL Z ENTROPY . . . . /
05870 0BFA 0008	ZENT	RMB	8
05880 0C02 20		FCC	/ BITS /
05890 0C0E CDOA		FDB	\$0D0A,\$0D0A
05900 0C12 54		FCC	/TOTAL SOURCE ENTROPY. . . . /
05910 0C2C 0008	TENT	RMB	8
05920 0C34 20		FCC	/ BITS /
05930 0C40 0D0A		FDB	\$0D0A,\$0D0A
05940 0C44 50		FCC	/PRESS RETURN FOR PAGE 1 OF STATISTIC
05950 0C6A 04		FCB	4
05960 0C6B 1A07	HDRMS2	FDB	\$1A07,\$0DCA

05970 0C6F 45	FCC	/EKG SAMPLE COLLECTION STATISTICS/
05980 0C8F 20	FCC	/ : PAGE 2/
06000 0C9A 10A	FCC	MAXIMUM CHANNEL COUNT
06000 0C9C 41	FCC	,INTERNAL MAX CONNECTIONS
06010 0CB2 0D0A	FDB	\$0D0A
06020 0CB4 52	FCC	/RATIO POSSIBLE. . . . . /
06030 0CCE 0008	MAXCPR RMB	8
06040 0CD6 20	FCC	/ : 1 /
06050 0CE2 0D0A	FDB	\$0D0A
06060 0CE4 43	FCC	/COMPRESSION RATIO/
06070 0CF5 0D0A	FDB	\$0D0A
06080 0CF7 41	FCC	/ACHIEVED. . . . . . /
06090 0D11 0008	CPRACH RMB	8
06100 0D19 20	FCC	/ : 1 /
06110 0D25 0D0A	FDB	\$0D0A
06120 0D27 41	FCC	/ACHIEVED COMPRESSION/
06130 0D3B 0D0A	FDB	\$0D0A
06140 0D3D 45	FCC	/EFFICIENCY. . . . . /
06150 0D57 0005	CPREFF RMB	5
06160 0D5C 20	FCC	/ % OF MAXIMUM /
06170 0D6B 0D0A	FDB	\$0D0A
06180 0D6D 43	FCC	/COMPRESSION TIME/
06190 0D7D 0D0A	FDB	\$0D0A
06200 0D7F 45	FCC	/EFFICIENCY OBTAINED . . . /
06210 0D99 0005	TIMEFF RMB	5
06220 0D9E 20	FCC	/ % SMP INTERVAL/
06230 0DAD 0D0A	FDB	\$0D0A,\$0D0A
06240 0DF1 43	FCC	/COLLECTION DURATION . . . /
06250 0DCB 0005	COLDUR RMB	5
06260 0DD0 20	FCC	/ SECONDS /
06270 0DDF 0D0A	FDB	\$0D0A,\$0D0A
06280 0DE3 43	FCC	/CHANNEL X MAXIMUM . . . . /
06290 0DFD 0008	AMAXX RMB	8
06300 0E05 20	FCC	/ VOLTS /
06310 0E11 0D0A	FDB	\$0D0A,\$0D0A
06320 0E15 .3	FCC	/CHANNEL X MINIMUM . . . . /
06330 0E2F 0008	AMINX RMB	8
06340 0E37 20	FCC	/ VOLTS /
06350 0E43 0D0A	FDB	\$0D0A,\$0D0A
06360 0E47 43	FCC	/CHANNEL Y MAXIMUM . . . . /
06370 0E61 0008	AMAYY RMB	8
06380 0E69 20	FCC	/ VOLTS /
06390 0E75 0D0A	FDB	\$0D0A,\$0D0A
06400 0E79 43	FCC	/CHANNEL Y MINIMUM . . . . /
06410 0E93 0008	AMINY RMB	8
06420 0E9B 20	FCC	/ VOLTS /
06430 0EA7 0D0A	FDB	\$0D0A,\$0D0A
06440 0EB3 43	FCC	/CHANNEL Z MAXIMUM . . . . /
06450 0EC5 0008	AMAXZ RMB	8
06460 0ECB 20	FCC	/ VOLTS /
06470 0ED9 0D0A	FDB	\$0D0A,\$0D0A
06480 0EDD 43	FCC	/CHANNEL Z MINIMUM . . . . /
06490 0EF7 0008	AMINZ RMB	8
06500 0EFF 20	FCC	/ VOLTS /

06510 0F0B CD0A	FDB	\$0D0A, \$CD0A
06520 0F0F 43	FCC	/COMMENTS. . . . . . /
06530 0F11 000A	FCC	
06550 0F62 50	FCC	/PRESS RETURN FOR PAGE 3 OF STATISTIC
06560 0F88 04	FCB	4
06570 0F89 1A07	COSTAT FDB	\$1A07, \$OC0D, \$CACA, \$0A0A, \$CA0A, \$0A0A
06580 0F95 45	FCC	/EKG SAMPLE COLLECTION STATISTICS: PA
06590 0FD0 0D0A	FDB	\$0D0A CR/LF
06600 0FBF 20	FCC	/ NUMBER OF SAMPLES TAKEN /
06610 0FD8 28	FCC	/(SAMPLE) = /
06620 0FE2 0004	NUMSMP RMB	4
06630 0FE6 20	FCC	/ (HEX) /
06640 0FFC 0D0A	FDB	\$0D0A CR/LF
06650 0FFE 20	FCC	/ MAXIMUM LOOP COUNT PER /
06660 1005 20	FCC	/ INTERRUPT (IPCNT) = /
06670 1019 0002	NUMPS RMB	2
06680 101B 20	FCC	/ (HFX) /
06690 1021 000A	FDB	\$0D0A CR/LF
06700 1023 20	FCC	/ TOTAL WAITING LOOP COUNTS DURING /
06710 1044 20	FCC	/ COLLECTION (LOOPCT) = /
06720 105A 0008	TOTLUP RMB	8
06730 1062 20	FCC	/ (HEX) /
06740 1068 CD0A	FDB	\$0D0A
06750 106A 54	FCC	/TIME EFFICIENCY = (1-(ICOPCT)*(SAMPLE
06760 1097 2A	FCC	/*100/
06770 109B CD0A	FDB	\$0D0A, \$0D0A
06780 109F 43	FCC	/CHANNEL MAXIMUMS AND MINIMUMS /
06790 10BC CD0A	FDB	\$0D0A
06800 10FE 20	FCC	/ XMAX= /
06810 10C5 0002	XMAXNM RMB	2
06820 10C7 20	FCC	/ (HEX) AT SAMPLE NUMBER /
06830 10DF 0004	XMAXLO RMB	4
06840 10E3 20	FCC	/ (HEX) /
06850 10E9 0D0A	FDB	\$0D0A
06860 10EB 20	FCC	/ XMN= /
06870 10F2 0002	XMINNM RMB	2
06880 10F4 20	FCC	/ (HEX) AT SAMPLE NUMBER /
06890 110C 0004	XMINLO RMB	4
06900 1110 20	FCC	/ (HEX) /
06910 1116 0D0A	FDB	\$0D0A
06920 1118 20	FCC	/ YMAX= /
06930 111F 0002	YMAXNM RMB	2
06940 1121 20	FCC	/ (HEX) AT SAMPLE NUMBER /
06950 1139 0004	YMAXLO RMB	4
06960 113D 20	FCC	/ (HEX) /
06970 1143 CD0A	FDB	\$0D0A
06980 1145 20	FCC	/ YMN= /
06990 114C 0002	YMINNM RMB	2
07000 114E 20	FCC	/ (HEX) AT SAMPLE NUMBER /
07010 1166 0004	YMINLO RMB	4
07020 116A 20	FCC	/ (HEX) /
07030 1170 0D0A	FDB	\$0D0A
07040 1172 20	FCC	/ ZMAX= /

## PRSTAT-9

07050 1179 0002	ZMAXNM RMB	2	/ (HEX) AT SAMPLE NUMBER /
07060 117D 20	FCC		
07070 1181 40	ZMINNM RMB	1	
07080 1197 20	FCC	/ (HEX)/	
07090 119D 0D0A	FDB	\$0D0A	
07100 119F 20	FCC	/ ZMIN= /	
07110 11A6 0002	ZMINNM RMB	2	/ (HEX) AT SAMPLE NUMBER /
07120 11A8 20	FCC		
07130 11C0 0004	ZMINLO RMB	4	
07140 11C4 20	FCC	/ (HEX)/	
07150 11CA 0D0A	FDB	\$0D0A,\$CD0A	
07160 11CE 43	FCC	/COMPRESSION STATISTICS:/	
07170 11E5 0D0A	FDB	\$0D0A	
07180 11E7 20	FCC	/ NUMBER OF MEMORY BITS /	
07190 11FE 41	FCC	/AVAILABLE = /	
07200 120A 0006	AVABIT RMB	6	
07210 1210 20	FCC	/ (HEX)/	
07220 1216 0D0A	FDB	\$0D0A	
07230 1218 20	FCC	/ NUMBER OF BITS AVAILABLE/	
07240 1231 20	FCC	/ TO VAR LEN CODER = /	
07250 1245 0008	TCPBIT RMB	8	
07260 124D 20	FCC	/ (HEX)/	
07270 1253 0D0A	FDB	\$0D0A	
07280 1255 20	FCC	/ TOTAL NUMBER OF DATA BITS/	
07290 126F 20	FCC	/ STORED = /	
07300 1279 0008	TOTBIT RMB	8	
07310 1281 20	FCC	/ (HEX)/	
07320 1287 0D0A	FDB	\$0D0A	
07330 1289 20	FCC	/ NUMBER OF BITS USED TO/	
07340 12A0 20	FCC	/ STORE CHANNEL X = /	
07350 12B3 0006	XBITNM RMB	6	
07360 12B9 20	FCC	/ (HEX)/	
07370 12BF 0D0A	FDB	\$0D0A	
07380 12C1 20	FCC	/ NUMBER OF BITS USED TO/	
07390 12D8 20	FCC	/ STORE CHANNEL Y = /	
07400 12EB 0006	YBITNM RMB	6	
07410 12F1 20	FCC	/ (HEX)/	
07420 12F7 0D0A	FDB	\$0D0A	
07430 12F9 20	FCC	/ NUMBER OF BITS USED TO/	
07440 1310 20	FCC	/ STORE CHANNEL Z = /	
07450 1323 0006	ZBITNM RMB	6	
07460 1329 20	FCC	/ (HEX)/	
07470 132F 0D0A	FDB	\$0D0A	
07480 1331 20	FCC	/ NUMBER OF BITS USED TO/	
07490 1348 20	FCC	/ STORE TIME OR OTHER PARAM/	
07500 1362 45	FCC	/ETER = /	
07510 1369 0006	TBITNM RMB	6	
07520 136F 20	FCC	/ (HEX)/	
07530 1375 0D0A	FDB	\$0D0A	
07540 1377 43	FCC	/COMPRESSION RATIO = TOTAL /	
07550 1391 44	FCC	/DATA BITS SHARED PER MSG /	
07560 13A8 42	FCC	/BITS AVAILABLE/	
07570 13B8 0D0A	FDB	\$0D0A,\$0D23	
07580 13PC 50	FCC	/PRNG RETURN =/	

PRSTAT-9

07590 13CA 04                  FCB      4  
07600                  \*  
07610                  \* END OF PRSTAT  
07620                  \*  
07630                  END

```

00030 *****  

00040 *  

00050 * OWNER NAME : NOCPRS  

00060 * AUTHOR : CAPT. MEL TOWNSEND  

00070 * VERSION : 1.7  

00080 * VERSION DATE : 3 OCT 80  

00090 *  

00100 *****  

00110 *  

00120 * OVERLAY DESCRIPTION  

00130 *  

00140 * THIS OVERLAY SAMPLES THE EKG DATA VIA THE  

00150 * A/D CONVERTERS AND STORES THE 8 BIT ROUNDED  

00160 * VALUES INTO MEMORY LOCATIONS 3C00-7FFF. THE  

00170 * THE X,Y,Z CHANNELS ARE SAMPLED AND STORED  

00180 * SEQUENTIALLY STARTING AT 3C00 AND WORKING UP.  

00190 *  

00200 *****  

00210 *****  

00220 *  

00230 * START OF NOCPRS  

00240 *  

00250 *****  

00260 *  

00270 0100 ORG $0100 OVERLAY START ADDRESS  

00280 *  

00290 OPT O ASSEMBLY-CODE FILE  

00300 OPT NOG ASSB OPT-SUPPRESS FULL FCC LI  

00310 *  

00320 *****  

00330 *  

00340 * LABEL DECLARATIONS  

00350 *  

00360 * SUPPORT SUBROUTINE ADDRESSES  

00370 *  

00380 CA87 OUTPUT EQU SCA87 FREQDOS. ALPH STRING TO CONS  

00390 CA8F OUTCR EQU SCA8F FREQDOS. ALPH STRNG, NO CR/LF  

00400 CA2C KEYFD INU SCAC2C FREQDOS. CONSOLE INPUT ROTTI  

00410 CA36 KEYDD INU SCAC6 FREQDOS. CONSOLE INPUT. NO ?  

00420 1D00 START EQU $1D00 ERG-EMC. START ERG-EMC  

00430 *  

00440 * DATA BUFFERS  

00450 *  

00460 3400 HDRSTR EQU $3400 DATA MEMORY BUFFER PTR/ADR  

00470 FFF8 IREQVC EQU $FFF8 INITVEC ADDRESS PTR  

00480 1C96 STKSAV EQU $1C96 STACK SAVE BUFFER  

00490 1C98 CPRTYP EQU $1C98 CENTER SCREEN BUFFER PTR  

00500 3002 ENDBUF EQU $3002 ADP C/LIST CHR PTR  

00510 1C9D VECSAV EQU $1C9D IPO VECTOR SAVE PTR  

00520 1CA1 FILHLC EQU $1CA1 FILTER AND HOLD PTR  

00530 1CA3 SAVHLC EQU $1CA3 SAVHLP AND HOLD PTR  

00540 3490 LOOPCT EQU $3490 TOTAL NUMBER OF LOOPS  

00550 3494 SAMMO EQU $3494 NUMBER OF SAMPLES  

00560 3496 LICAL EQU $3496 AVG'D MAX DENSITY

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00570	3497	MAXZ EQU	\$3497	MAX VLU IN CH Z
00571	3498	MAXXLO EQU	\$3498	MAX VLU LOC IN CH 7
		MAXXHI EQU	\$3499	MAX VLU LOC IN CH 8
00600	349A	MINXLO EQU	\$349A	MIN VLU LOC IN CH Z
00610	349D	MAXY EQU	\$349D	MAX VLU IN CH Y
00620	349E	MAXYLO EQU	\$349E	MAX VLU LOC IN CH Y
00630	34A0	MINY EQU	\$34A0	MIN VLU IN CH Y
00640	34A1	MINYLO EQU	\$34A1	MIN VLU LOC IN CH Y
00650	34A3	MAXX EQU	\$34A3	MAX VLU IN CH X
00660	34A4	MAXXLO EQU	\$34A4	MAX VLU LOC IN CH X
00670	34A6	MINX EQU	\$34A6	MIN VLU IN CH X
00680	34A7	MINXLO EQU	\$34A7	MIN VLU LOC IN CH X
00690	34AC	DTABIT EQU	\$34AC	NUM OF BITS USED TO STR DTA.
00700	34B0	XBITS EQU	\$34B0	NUM OF BITS USED TO STR X
00710	34B3	YBITS EQU	\$34B3	NUM OF BITS USED TO STR Y
00720	34B6	ZBITS EQU	\$34B6	NUM OF BITS USED TO STR Z
00730	34B9	TBITS EQU	\$34B9	NUM OF BITS USED TO STR T
00740	34C1	ACILCT EQU	\$34C1	# BITS SET TO VLG IN CIR HRC
00750	3600	XPDYI EQU	\$3600	0 VLG LOC OF X FILE
00760	3800	YPDYI EQU	\$3800	0 VLG LOC OF Y FILE
00770	3A00	ZPDFM EQU	\$3A00	0 VLG LOC OF Z FILE
00780	3B00	TPDF EQU	\$3B00	0 VLG LOC TO TIME FILE
00790	3C00	SECZRO EQU	\$3C00	SWAIT OF DATA ITEM FILE
00800	*			
00810		* HARDWARE ADDRESSES		
00820		*		
00830	E400	ADCZRO EQU	\$E400	DAC CHANNEL ZERO
00840	E404	ADCTWO EQU	\$E404	NDC CHANNEL TWO
00850	E500	DACZRO DQU	\$F500	DAC CHANNEL ZERO
00860		*		
00870		*		
00880		*FUNCTION :NOCPRS		
00890		*INPUTS :STATUS REGISTERS FROM EKG-EXEC		
00900		*OUTPUTS :DATA TO DISK		
00910		*CALLS :OUTPUT,PUTADR,PUTRD,OUTNCR		
00920		* KEYBD0,SVTTL,SWINT		
00930		*DESTROYS :ALL REGISTERS		
00940		*PURPOSE :TO COLLECT 3 CHANNELS OF EKG		
00950		* DATA AND STORE INTO MEMORY.		
00960		*		
00970		*		
00980	0100 OF	NOCPRS SEI		
00990	0101 CE 0233	LDX #STRING		
01000	0104 DD C1.8F	JSR OUTWR	"THIS WHILE SAVING THE..."	
01010	0107 DD C1.2C	JSR KEYRD	GMR ELEMENTS FROM CONSOLE	
01020	010A FE 3002	LDX ENDPR		
01030	010D 09	DEX		
01040	010E F6 00	LDA B 0,X		
01050	0110 C1 59	CMP B #'Y	IS INPUT YES?	
01060	0112 27 03	BHQ NOCPRI	YES. KEEP EXECUTING THIS ROUT	
01070	0114 7E 1D00	JMP START	NO. RET TO EKG-EXEC	
01080	0117 CM 4E43 NOCPRI	LDX "S4E43"	FLAG COMMISSION TYPE (NC)	
01090	011A FF 1C98	STX CRNTYP?		
01100	011D BD 051A	JSR FILHDR	SET UP DATA FILE HEADER	

01110 0120 CE 02B5	LDX	#IROMSG		
01120 0121 FF 0000	STX	CONTINU	"INITIAL COUNT" AND ETC..."	
01140 0129 CE 3C00 NOCIGO	LDX	#SECZRO	INITIALIZE VALUE OF NEXT LO.	
01150 012C FF 0230	STX	BUFPTR	STORE IN BUFFER POINTER	
01160 012F 86 03	LDA A	#3	PICK UP COUNTER FOR THE CAL	
01170 0131 B7 0227	STA A	CALCNT	STORE IN BUFFER	
01180 0134 FE FFF8	LDX	IRQVEC	PICK UP CURRENT IRQ VECTOR	
01190 0137 FF 1C9D	STX	VFCSAV	SAVE IN BUFFER	
01200 013A CE 051F TIMCAL	LDX	#CALINT	GET INTR VECTOR ADDR FOR CAL	
01210 013D FF FFFF	STX	IRQVEC	PUT IN VECTOR ADDRESS	
01220 0140 86 00	LDA A	#0	INIT COUNTER FOR 256 TEST LOG	
01230 0142 B7 0232	STA A	DONTST	STORE IN DONTST BUFFER	
01240 0145 CE 0000	LDX	#0	CLEAR LOOP COUNTER	
01250 0148 FF 022C	STX	IRQCJT		
01260 014E FF 022E	STX	IRQCJT+2		
01270 014E FF E500	STX	DACZRO	PULSE INT' ENABLE CIPCJIT	
01280 0151 20 09	BKA	SPO01	PPA TO COUNTING LOOP	
01290	*			
01300	*****			
01310	*			
01320	* BASIC TIMING LOOP FOR EFFICIENCY TEST			
01330	*			
01340	*****			
01350	*			
01360 0153 01	SPLOOP	NOP	DELAY TO MATCH TIME	
01370 0154 01		NOP	IT TAKES TO EXECUTE ONE	
01380 0155 01		NOP	INCREMENT OF BYTES 3&4	
01390 0156 01		NOP	OF THE LOOP COUNTER	
01400 0157 01		NOP	WHEN COUNT CARRIES TO	
01410 0158 01		NOP	HIGH 2 BYTES OF 4-BYTE	
01420 0159 7E 015F	JMP	SPO03	JUMP TO COUNTING LOOP	
01430 015C 05	SPO01	CLI	PREPARE FOR TIMER INT	
01440 015D 3E		WAI	STOP PROCESSOR & WAIT	
01450 015E 0E	SPO04	CLI		
01460 015F FD 022E	LDX	IRQCJT+2	RET'ED FROM INT, INC COUNT	
01470 0162 08	INX			
01480 0163 FF 022E	STX	IRQCJT+2	SAVE IT	
01490 0166 26 FB	BNE	SPLOOP	COUNT GONE FFFF TO 0000?	
01500 0168 FE 022C	LDX	IRQCJT	YES. INC BYTES 3&4	
01510 016B 08	INX			
01520 016C FF 022C	STX	IRQCJT	SAVE COUNT	
01530 016F B6 0232	SPO03	LDA A	DONEST	IS DONE TEST SATISIFIED?
01540 0172 81 00		CMP A	#0	
01550 0174 26 E8	BNE	SPO04	NO. KEEP LOOPING	
01560	*			
01570	*****			
01580	*			
01590	* END BASIC TIMING LOOP FOR EFF TEST			
01600	*			
01610	*****			
01620	*			
01630 0176 0F	SEI		PREVENT SER TO NVM. INTR	
01640 0177 CE 4000	IDX	#\$4000	PICK UP INTR OUT WORD	

01650 017A FF F500	STX	DACTRO	DISPLAY INTR FLAG TIMER
01660 017D B7 E400	STA A	ADCTRO	CLR VME6800 INT FLIP FLOP
01670 0180 0000	LDA A	CM C1	LOAD TIME COUNT
01680 0184 27 38	REQ	SPXNE	IS COUNT 32 HRS, GO SAVE DATA
01700 0185 F6 0227	LDA A	CALCNT	NO, GLT COUNT
01710 0189 81 02	CMP A	#2	IS COUNT 2?
01720 018B 26 22	ENR	PARSAV	NO, SAVE PARAMETER FROM INPUT
01730 018D FE 022E	LDX	IRQGT+2	YES, SAVE TIME LOOP CENTER
01740 0190 FF 022A	STX	CALONE	
01750 0193 CE 0316	LDX	#SAV1E	NOW PUT SAMPLE ADDR IN IRQVEC
01760 0195 FF FFFF	STX	IRQVEC	
01770 0199 86 AA	LDA A	#SAA	SET UP DCNTST
01780 019D B7 0232	STA A	DONIST	
01790 019E CE 0000	LDX	#0	CLR LOOP COUNT AND SAMPLE COU
01800 01A1 FF 022C	STX	IRQGT	
01810 01A4 FF 022E	STX	IRQGT+2	
01820 01A7 FF 3194	STX	SAMING	
01830 01A8 FF F500	STX	DACTRO	ENABLE TIMER INTERRUPTS
01840 01AD 20 AD	ENR	SPOLL	GO WAIT FOR INTERRUPT
01850 01AF FE 022C PARSAV	LDX	IRQGT	SAVE LOOP COUNT
01860 01B2 FF 3490	STX	LOOPCP	
01870 01B5 FE 022E	LIX	IRQGT+2	
01880 01E8 FF 3492	STX	LOOPCP+2	
01890 01EB 7F 013A	JMP	TINCL	GO EXECUTE ANOTHER TIME CAL
01900 01BE B6 022E SPDCNE	LEA A	CALC0+1	AVERAGE TWO TIME CAL RESULTS
01910 01C1 B8 022F	ADD A	IRQGT+3	ADD LSB BYTES
01920 01C4 B7 0229	STA A	CALZ0+1	STORE IN BUFFER
01930 01C7 F6 022A	LDA A	CALCNE	PICK UP USE OF FIRST
01940 01CA B9 022E	ADC A	IRQGT+2	ADD WITH CARRY
01950 01CD B7 0228	STA A	CALZ0	STORE IN BUFFER
01960 01D0 77 0229	ASR	CALZ0	DIVIDE BY 2 TO AVERAGE
01970 01D3 76 0229	ROR	CALZ0+1	
01980 01D6 85 68	LDA A	#8	SET UP COUNTER FOR /256
01990 01D8 77 0228 CMISHF	ASR	CALZ0	NOT DIVIDE BY 256 FOR
02000 01D9 76 0229	ROR	CALZ0+1	LOCK/TIMER/INT
02010 01DE 4A	DEC A		
02020 01D9 26 F7	ENR	CALSHF	8 SHIFTS YET?
02030 01E1 B6 0229	LDA A	CALZ0+1	PICK UP SHIFTED INPUT
02040 01E4 B7 3496	STA A	LICNL	STORE IN FILE PAR BUFFER
02050 01E7 B6 34A7	LDA A	DTNPIT+3	PUT DATA BIT COUNT INTO ACFLC
02060 01EA B7 34FF	STA A	ACFLCP+3	SINCE NO COMPRESSION PERFORM
02070 01FB B6 34AE	LDA A	DTNPIT+2	
02080 01F0 B7 34DE	STA A	ACFLCP+2	
02090 01F3 B6 34AD	LDA A	DTNPIT+1	
02100 01F6 B7 341D	STA A	ACFLCP+1	
02110 01F9 B6 34AC	LDA A	DTADIT	
02120 01FC B7 34DC	STA A	ACFLCT	
02130 01FF FE 1C9D	LDX	VRCSSW	YES, RESTORE IRQ VECTOR
02140 0202 FF FFFF	STX	IRQVEC	
02150 0205 CE 0312	LDX	#SAVMSG	
02160 0208 FD GME	JSR	OUTNCR	"SAMPLING COMPLETE..."
02170 020B FD CMC	JSR	KEYTD	GLT SAVE DECISION
02180 020E FE 3002	LDX	ENDBUF	

02190 0211 09	DIX		
02200 0212 F6 00	LDA B	0,X	
02210 0214 C1 F9	CMP B	#'Y	IS P'cision YES?
02220 0216 28 0C	PUL	EXJMP	NO. JUMP TO EXJMP
02230 0218 FD 0515	JSR	SAVEFL	YES. SAVE THE FILE ON DISK
02240 021B CF 0381	IDX	#SDONE	
02250 021E BD CA8F	JSR	OUTCR	
02260 0221 FD CA2C	JSR	KEYED	
02270 0224 7E 1D00	EXJMP	JMP	JUMP BACK TO EKG-FNDC
02280 *			
02290 0227 0001	CALCNT RMB	1	TEST CAL. LOOP COUNTER
02300 0228 0002	CALZHO RMB	2	AVG'D CAL LOGIC DURING LAST CO
02310 022A 0002	CALONE RMB	2	TEMP RUF FOR PRE CHLLG CAL
02320 022C 0004	IRQCIT RMB	4	TEMP INTR. LOOP COUNTER
02330 0230 0002	BUFFPTR RMB	2	BUFFER POINTER : OT NEXT AVAIL
02340 0232 0001	DONTSTT RMB	1	DONE TEST FLG FOR MEM FULL
02350 *			
02360 0233 1A07	STRMSG FDR	\$1A07	
02370 0235 54	FCC	/THIS MODULE SAMPLES THE EKG /	
02380 0251 49	FCC	/INPUT AND STORES THE DATA WITH/	
02390 026F 0D0A	FDB	\$0D0A	
02400 0271 4E	FCC	/NO DATA COMPRESSION./	
02410 0285 0D0A	FDB	\$0D0A,\$0D0A	
02420 0289 44	FCC	/DO YOU WISH TO EXECUTE THIS /	
02430 02A5 4D	FCC	/MODULE (Y OR N)/	
02440 02B4 04	FCB	4	
02450 02B5 1A	IRQMSG FCB	\$1A	
02460 02B6 49	FCC	/INSURE SUBJECT AND EKG DEVICE /	
02470 02D4 52	FCC	/READY! /	
02480 02DA 0D0A	FDB	\$0D0A,\$0D0A,\$0D0A	
02490 02E0 50	FCC	/PRESS RETURN, THEN CLOSE INTERRUPT/	
02500 0302 20	FCC	/ENABLE SWITCH./	
02510 0311 04	FCB	4	
02520 0312 1A07	SAVMSG FDB	\$1A07	
02530 0314 53	FCC	/SAMPLING COMPLETE. PLEASE OPEN /	
02540 0334 49	FCC	/INTERRUPT ENBLIN SWITCH./	
02550 034C 0D0A	FDB	\$0D0A,\$0D0A,\$0D0A	
02560 0352 44	FCC	/DO YOU WISH TO SAVE THIS DATA ON /	
02570 0373 44	FCC	/DISK (Y OR N)/	
02580 0380 04	FCB	4	
02590 0381 20	SDONE FCC	/ PRESS RETURN/	
02600 038C 04	FCB	4	
02610 *			*****
02620			*****
02630			JD NOCRS
02640			*****
02650			*
02660			*
02670			*FUNCTION :SAMPLE
02680			*INPUTS :STATUS BUFFERS
02690			*OUTPUTS :COMPRESSED,ROUNDED DATA IN MM BUFF
02700			*CALLS :NOTHING
02710			*PURPOSE :THIS ROUTINE SAMPLES THE EKG LEADS,
02720			* ROUNDS THE VALUES TO 3 BITS (PREC 12)

02730	* CALCULATES MAX,MIN, # OF BITS, SAMPLES		
02740	* ETC., AND SAVES THESE PARAM AND DATA		
	* TO INPUT IN ROM.		
02760	*		
02770	*		
02780 038F 0001	SHFRUF RMB	1	TEMP SHIFT BUFFER
02790 0390 0001	TEMPA1 RMB	1	TEMP REG, MSB CH 0
02800 0391 0001	TEMIB1 RMB	1	TEMP REG, LSB CH 0
02810 0392 0001	TEMPA2 RMB	1	TEMP REG, MSB CH 1
02820 0393 0001	TEMIB2 RMB	1	TEMP REG, LSB CH 1
02830 0394 0001	TEMPA3 RMB	1	TEMP REG, MSB CH 2
02840 0395 0001	TEMIB3 RMB	1	TEMP REG, LSB CH 2
02850	*		
02860 0396 FE 3494	SAMPLE LDX	SAMPNO	GET CUR SAMPL COUNT
02870 0399 08	INX		
02880 039A FF 3494	STX	SAMPNO	
02890 039D BF 1C96	STS	STKSAV	SAVE STACK PRT
02900 03A0 8E 3280	LDS	#\$3280	INITIALIZE STACK IN UNSHD MEM
02910 03A3 CE E400	LDX	#ADCZRO	NOW PULSE A/D TO START CONV
02920 03A6 A7 00	STA A	0,X	FOR CHANNEL 0
02930 03A8 01	NOP		
02940 03A9 A6 00	LDA A	0,X	
02950 03AB F6 01	LDA B	1,X	
02960 03AD B7 0390	STA A	TEMPA1	
02970 03E0 F7 0391	STA B	TEMIB1	
02980 03B3 08	INX		
02990 03B4 08	INX		
03000 03B5 A7 00	STA A	0,X	NOW PULSE A/D FOR CONV
03010 03B7 01	NOP		ON CHANNEL 1
03020 03B8 A6 00	LDA A	0,X	
03030 03BA E6 01	LDA B	1,X	
03040 03BC B7 0392	STA A	TEMPA2	
03050 03BF F7 0393	STA B	TEMIB2	
03060 03C2 08	INX		
03070 03C3 08	INX		
03080 03C4 A7 00	STA A	0,X	NOW PULSE A/D FOR CONV
03090 03C6 01	NOP		ON CHANNEL 2
03100 03C7 A6 00	LDA A	0,X	
03110 03C9 F6 01	LDA B	1,X	
03120 03CB B7 0394	STA A	TEMPA3	
03130 03CE F7 0395	STA B	TEMIB3	
03140 03D1 CE 0390	IDX	#TEMPA1	
03150 03D4 A6 00	SAMPLE LDX	LDA A	0,X
03160 03D6 B6 01	LDA B	1,X	
03170 03D8 47	ASR A		NOW SHIFT 2 BYTE VLN IN REG'
03180 03D9 56	ROR B		4 POSITIONS TO RIGHT IN REG
03190 03DA 47	ASR A		12 BIT TO 8 BIT ROUTED CONV
03200 03DB 56	ROR B		
03210 03DC 47	ASR A		
03220 03DD 56	ROR B		
03230 03DE 47	ASR A		
03240 03DF 56	ROR B		
03250 03E0 F7 039F	STA B	SHFRUF	SAVE 8 BIT RESULT OF SHIFT
03260 03E3 86 00	LDA A	#0	AND ADD CARRY OUT OF LAST

03270 03F5 B9 039F	ADC A	SUBBUF	SHIFT ROUNDING UP OF DATA
03280 03F8 36	PSH A		SAVE VLU IN STACK TEMPORARILY
03290 03F9 00	PUL A		
03300 03A1 05	LDX		
03310 03EB 8C 0396	CPX	#TEMPA3+2	CHANNEL Z ROUNDED TO 8 BITS
03320 03EE 26 E4	BNE	SAMPL1	NO. GO SAMPLE NEXT CHANNEL
03330 03F0 32	PUL A		GET Z DATA
03340 03F1 FE 0230	LDX	BUFPTR	PICK UP CUR MEM FILE PTR
03350 03F4 A7 02	STA A	2,X	SAVE DATA TO MEM FILE
03360 03F6 FE 3494	LDX	SAMPNO	GET CUR SAMPLE COUNT
03370 03F9 B1 3497	CMP A	MAXZ	IS CUR Z MAX OVER SAMPLE SET?
03380 03FC 2F 06	BLE	SPZMIN	NO. GO CHECK FOR MIN
03390 03FE B7 3497	STA A	MAXZ	YES. KEEP CUR VLU
03400 0401 FF 3493	STX	MAXZLO	KEEP CUR SAMPLE NUM
03410 0404 B1 349A SPZMIN	CMP A	MINZ	IS CUR Z MIN OVER SAMPLE SET?
03420 0407 2C 06	BGE	SPYMAX	NO. GO CHECK FOR Y MAX
03430 0409 B7 349A	STA A	MINZ	YES. KEEP CUR VLU
03440 040C FF 349B	STX	MINZLO	KEEP CUR SAMPLE NUM
03450 040F 32	SPYMAX	PUL A	GET CUR Y DATA
03460 0410 FE 0230	LDX	BUFPTR	PICK UP CURRENT MEM FILE PTR
03470 0413 A7 01	STA A	1,X	SAVE DATA TO MEM FILE
03480 0415 FE 3494	LDX	SAMPNO	GET CUR SAMPLE COUNT
03490 0418 B1 349D	CMP A	MAXY	IS CUR Y MAX OVER SAMPLE SET?
03500 041B 2F 06	BLE	SPYMIN	NO. GO CHEK FOR MIN
03510 041D B7 349D	STA A	MAXY	YES. KEEP CUR Y VLU
03520 0420 FF 349E	STX	MAXYLO	KEEP CUR SAMPLE NUM
03530 0423 B1 34A0 SPYMIN	CMP A	MINY	IS CUR Y MIN OVER SAMPLE SET?
03540 0426 2C 06	BGE	SPYMAX	NO. GO CHEK FOR X MAX
03550 0428 B7 34A0	STA A	MINY	YES. KEEP CUR Y VLU
03560 042B FF 34A1	STX	MINYLO	KEEP CUR SAMPLE NUM
03570 042E 32	SPYMAX	PUL A	GET CUR X DATA
03580 042F FE 0230	LDX	BUFPTR	PICK UP CURRENT MEM FILE PTR
03590 0432 A7 00	STA A	0,X	SAVE DATA TO MEM FILE
03600 0434 FE 3494	LDX	SAMPNO	GET CUR SAMPLE COUNT
03610 0437 B1 34A3	CMP A	MAXX	IS CUR X MAX OVER SAMPLE SET?
03620 043A 2F 06	BLE	SPXMIN	NO. GO CHEK X MIN
03630 043C B7 34A3	STA A	MAXX	YES. KEEP CUR X VLU
03640 043F FF 34A4	STX	MAXXLO	KEEP CUR SAMPLE VLU
03650 0442 B1 34A6 SPXMIN	CMP A	MINX	IS CUR X MIN OVER SAMPLE SET?
03660 0445 2C 06	BGE	MNDONE	NO. EXIT MAX/MIN UPDATE
03670 0447 B7 34A6	STA A	MINX	YES. KEEP CUR X VLU
03680 044A FF 34A7	STX	MINXLO	KEEP CUR SAMPLE COUNT
03690 044D 0C	MNDONE	CLC	DONE WITH MAX & MINS. NOW
03700 044E 86 18	LDA A	#24	UPDATE ACTUAL DATA BIT SWRD
03710 0450 C6 00	LDA B	#0	
03720 0452 EB 34AF	ADD A	DTABIT+3	
03730 0455 B7 34AF	STA A	DTABIT+3	
03740 0458 86 00	LDA A	#0	
03750 045A F9 34AE	ADC B	DTABIT+2	
03760 045D F7 34AE	STA B	DTABIT+2	
03770 0460 C6 00	LDA B	#0	
03780 0462 B9 34AD	ADC A	DTABIT+1	
03790 0465 B7 34AD	STA A	DTABIT+1	
03800 0468 F9 34AC	ADC B	DTABIT	

03810 046B F7 34AC	STA B	DYBIT	
03820 046E 86 08	LDA A	#8	NOW UPDATE XBIT, YBIT, ZBIT COU
03830 0470 C6 0C	ADD A	TEMP+1	UPD X, Y, Z COUNTERS & INC
03840 0472 BB 34B2	ADD A	TEMP+1	X, Y, Z COUNTERS & INC.
03850 0475 B7 34B2	STA A	XBITS+2	
03860 0478 B7 34B5	STA A	YBITS+2	
03870 047B B7 34B8	STA A	ZBITS+2	
03880 047E 86 00	LDA A	#0	
03890 0480 F9 34B1	ADC B	XBITS+1	
03900 0483 F7 34B1	STA B	XBITS+1	
03910 0486 F7 34B4	STA B	YBITS+1	
03920 0489 F7 34B7	STA B	ZBITS+1	
03930 048C B9 34B0	ADC A	XBITS	
03940 048F B7 34B0	STA A	XBITS	
03950 0492 B7 34B3	STA A	YBITS	
03960 0495 B7 34B6	STA A	ZBITS	
03970 0498 7F 34B8	CLR	TBITS+2	
03980 049B 7F 34BA	CLR	TBITS+1	
03990 049E 7F 34B9	CLR	TBITS	
04000 04A1 34	DES		NOW POINT STACK BACK TO DATA
04010 04A2 34	DES		
04020 04A3 34	DES		
04030 04A4 32	PUL A		GET Z DATA
04040 04A5 CE 3A00	LDX	#ZPDF1	
04050 04A8 BD 04D2	JSR	PDFSTR	UPDATE Z PDF BIN COUNTER
04060 04AB 32	PUL A		GET Y DATA
04070 04AC CE 3800	LDX	#YPDF1	
04080 04AF BD 04D2	JSR	PDFSTR	UPDATE Y PDF BIN CONTER
04090 04B2 32	PUL A		
04100 04B3 CE 3600	LDX	#XPDF1	
04110 04B6 BD 04D2	JSR	PDFSTR	
04120 04B9 FE 0230	LDX	BUFPTR	NOW GET & UPDATE MEM BUFF PTR
04130 04BC 08	INX		
04140 04BD 08	INX		
04150 04BE 08	INX		
04160 04BF FF 0230	STX	BUFPTR	
04170 04C2 BE 1C96	LDS	STKSAV	RETRIEV ENTRY STACK POINTER
04180 04C5 8C 7FFE	CPX	#\$7FFE	IS MEM BUF FULL?
04190 04C8 26 03	ENE	SAMRTI	NO. RET FROM INTR & SAMPLE AG
04200 04CA 7F 0232	CLR	DONTST	YES. RESET MEM FULL FLAG
04210 04CD 3B	SAMRTI	RTI	
04220 *			
04230 04CE 0002	TEMPWF RMB	2	TEMP WORKING BUFFER
04240 04D0 0002	TEMPST RMB	2	TEMP STACK SAVE PTR
04250 *			
04260 04D2 BF 04D0	PDFSTR SIS	TEMPST	SAVE STACK PTR
04270 04D5 FF 04CE	STX	TEMPDF	SAVE PDF PTR
04280 04D8 16	TAB		SAVE INPUT DATA VLI
04290 04D9 BB 04CF	ADD A	TEMPDF+1	NOW CALCULATE ADDRESS OF INC
04300 04DC B7 04CF	STA A	TEMPDF+1	IN THE PDF MEM BUFFER.
04310 04DF 86 00	LDA A	#0	
04320 04E1 C5 80	BIT B	#\$80	IS THE DATA VLI NEG?
04330 04E3 2B 05	EPI	PDFSTI	YES. SRC MSB VS ADC
04340 04E5 B9 04CE	ADC A	TEMPDF	NO. ADD WITH CARRY MSB

04350 04E8 20 03	RRA	PDFST2	AND STORE
04360 04FA B2 04CE	SPC A	TEMPDF	VLU IS NEG. SBC
04370 04F0 B7 04CF	STA A	.....	TO TEMPDF.
04380 04F0 17	TBA	.....	LOAD THE VLU AND ADD IT TO TEMPDF.
04390 04F1 BB 04CF	ADD A	TEMPDF+1	FOR PROPER ADDR CALCULATION
04400 04F4 B7 04CF	STA A	TEMPDF+1	
04410 04F7 86 00	LDA A	#0	
04420 04F9 C5 80	BIT B	#\$80	IS THE DATA NEG?
04430 04FB 2B 05	BMI	PDFST3	YES. SBC MSB VS ADC
04440 04FD B9 04CE	ADC A	TEMPDF	NO. ADD WITH CARRY MSB BYTE
04450 0500 20 03	ERA	PDFST4	AND STORE
04460 0502 B2 04CE	SPC A	TEMPDF	VLU IS NEG. SUB WITH CARRY
04470 0505 B7 04CF	STA A	TEMPDF	AND STORE
04480 0508 FE 04CE	LDX	TEMPDF	NOW LOAD CALC ADDR FOR INDEX
04490 050B AE 00	LDS	0,X	GRAB VLU IN CALC ADDR
04500 050D 31	INS	.....	INCREMENT IT
04510 050E AF 00	STS	0,X	AND STORE IT BACK IN BUFFER
04520 0510 BE 04D0	LDS	TEMPST	RECOVER STACK POINTER
04530 0513 17	TBA	.....	RECOVER ORIGINAL DATA
04540 0514 39	RTS	.....	
04550	*	.....	
04560	*FUNC: EXEC JUMPS	.....	
04570	*INPUTS: ACCUMULATORS	.....	
04580	*OUTPUTS: NCNE	.....	
04590	*CALLS: SAVEFL,FILHDR (FEG-EXEC) VIA ADDR BUFFERS	.....	
04600	*DESTROYS: X,A,B,CC	.....	
04610	*PURPOSE: TO JUMP TO DESIRED ROUTINES VIA RELOC ADDR	.....	
04620	*	.....	
04630 0515 FE 1CA3	SAVEFL LDX	SAVELC	GET ADDR OF SAVEFL
04640 0518 6E 00	JMP	0,X	JUMP TO SAVEFL
04650	*	.....	
04660 051A FE 1CA1	FILHDR LDX	FILHLC	GET ADDR OF FILHLC
04670 051D 6E 00	JMP	0,X	JUMP TO SAVEFL
04680	*	.....	
04690	*****	*****	*****
04700	* END SAMPLE	*****	*****
04710	*****	*****	*****
04720	*	*****	*****
04730	*	*****	*****
04740	*FUNCTION :CALINT	*****	*****
04750	*INPUTS (REG) :NCNE	*****	*****
04760	*OUTPUTS (REG) :NONE	*****	*****
04770	*CALLS :NOTHING	*****	*****
04780	*DESTROYS (FEG):NCNE (INTERRUPT HANDLER)	*****	*****
04790	*PURPOSE :THIS ROUTINE IS USED FOR CALIBRATING	*****	*****
04800	* THE MAX NUMBER OF LOOPS POSSIBLE	*****	*****
04810	* DURING AN INTERRUPT INTRO. THIS	*****	*****
04820	* ROUTINE UPDATES THE DONST FLAG AND	*****	*****
04830	* RESETS THE ST6800 INT FLIP FLOP &	*****	*****
04840	* THE RETURNS.	*****	*****
04850	*	*****	*****
04860 051F 7C 0232	CNINT INC	DONST	INCREMENT DONF TEST FLAG
04870 0522 B7 E400	STA A	ADCZRO	CLR ST6800 INTR FLIP FLOP
04880 0525 01	NOP	.....	

NOCPRS-7

04890 0526 3B                   RTI  
04900  
04910  
04920                           \* END CALINT  
04930                           \*\*\*\*\*  
04940                           \*  
04950                           \*  
04960                           \*\*\*\*\*  
04970                           \*  
04980                           \* END OF NOCPRS OVERLAY ROUTINES  
04990                           \*  
05000                           \*\*\*\*\*  
05010                           \*  
05020                           END

## TOLAN-A

```

00030 ****
00040 *
00050 *
00060 * OWNER NAME: TOLAN-A
00070 * AUTHOR : CAPT. MEL TOLAND
00080 * VERSION : 1.8
00090 * VERSION DATE : 22 OCT 80
00100 *
00110 ****
00120 * OVERLAY DESCRIPTION
00130 *
00140 * THIS OVERLAY SAMPLES THE EKG DATA VIA THE
00150 * A/D CONVERTERS, ROUNDS THE DATA TO 8 BITS,
00160 * AND THEN COMPRESSES THE DATA VIA THE TOLAN-A
00170 * ALGORITHM. THE COMPRESSED DATA IS THEN STORED INTO
00180 * MEMORY DATA FILE FROM 3C00-7FFF.
00190 *
00200 ****
00210 *
00220 * START OF TOLAN-A
00230 *
00240 ****
00250 *
00260 0100 ORG $0100 OVERLAY START ADDRESS
00270 *
00280 OPT O ASSB OPT-GLN OBJ FILE
00290 OPT NOG ASSB OPT-SUPPRESS FULL FCC LI
00300 *
00310 ****
00320 *
00330 * LABLE DECLARATIONS
00340 *
00350 * SUPPORT SUBROUTINE ADDRESSES
00360 *
00370 CA87 OUTPUT EQU SCA87 EPRGDOS. ALPH STRING TO CONS
00380 CA8F OUTNCR EQU SCA8F EPRGDOS. ALPH STRING, NO CR/LF
00390 CA2C KEYPD EQU SCA2C EPRGDOS. CONSOLE INPUT ROUTI
00400 CA36 KEYBD0 EQU SCA36 EPRGDOS. CONSOLE INPUT. NO ?
00410 1D00 START EQU $1D00 EKG-EXEC. FILE CREATE ROUTINE
00420 *
00430 * DATA BUFFERS
00440 *
00450 3400 HERSTR EQU $3400 DATA MEMORY BUFFER HEADER
00460 FFF8 IFQVNC EQU $FFP8 INTERRUPT VECTOR ADDR
00470 1C96 STKSAV EQU $1C96 STACK SAVE BUFFER
00480 1C98 CPRTYP EQU $1C98 COMPRESSION ALGOR FLAG
00490 3002 ENDBUF EQU $3002 ADDR OF LAST CHAR PRINT
00500 1C9D VLCSAV EQU $1C9D IRQ VECTOR SAVE BUFFER
00510 1CA1 FILHLC EQU $1CA1 FILTER ADDE LAST FILTER
00520 1CA3 SAVHLC EQU $1CA3 SAVITL ADDR OF TS FILTER
00530 3490 LOOPCT EQU $3490 TOTAL NUM TIME LOOP EXEC'D
00540 3494 SAMPLS EQU $3494 NUMBER OF SAMPLES TAKEN
00550 3496 LPCLL EQU $3496 AVG'D MAX LOOP COUNT/INTR
00560 3497 MAXZ EQU $3497 MAX VLU IN CH Z

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00570	3498	MAXZLO EQU	\$3498	MAX VLU LOC IN CH Z
00580	349A	MINZ EQU	\$349A	MIN VLU IN CH Z
00590	349C	MINZLO EQU	\$349C	MIN VLU LOC IN CH Z
00600	349D	MAXY EQU	\$349D	LAD Y IN CH Y
00610	349E	MAXYL0 EQU	\$349E	MAX VLU LOC IN CH Y
00620	34A0	MINY EQU	\$34A0	MIN VLU IN CH Y
00630	34A1	MINYLO EQU	\$34A1	MIN VLU LOC IN CH Y
00640	34A3	MAXX EQU	\$34A3	MAX VLU IN CH X
00650	34A4	MAXXL0 EQU	\$34A4	MAX VLU LOC IN CH X
00660	34A6	MINX EQU	\$34A6	MIN VLU IN CH X
00670	34A7	MINXL0 EQU	\$34A7	MIN VLU LOC IN CH X
00680	34AC	DITABIT EQU	\$34AC	NUM OF BITS USED TO STR DTA
00690	34B0	XBITS EQU	\$34B0	NUM OF BITS USED TO STR X
00700	34B3	YBITS EQU	\$34B3	NUM OF BITS USED TO STR Y
00710	34B6	ZBITS EQU	\$34B6	NUM OF BITS USED TO STR Z
00720	34B9	TBITS EQU	\$34B9	NUM OF BITS USED TO STR T
00730	34BC	ACELCT EQU	\$34BC	# BITS FED TO VAR LEN CODE
00740	3600	XPDFM EQU	\$3600	0 VLU LOC OF X PDF
00750	3800	YPDFM EQU	\$3800	0 VLU LOC OF Y PDF
00760	3A00	ZPDFM EQU	\$3A00	0 VLU LOC OF Z PDF
00770	3B00	TPDF EQU	\$3B00	0 VAL LOC OF TIME VAR HIST
00780	3C00	SECZRO EQU	\$3C00	DATA STORE ADDR START
00790	*			
00800		* HARDWARE ADDRESSES		
00810	*			
00820	E400	ADCZRO EQU	\$E400	ADC CHANNEL ZERO
00830	E404	ADCTWO EQU	\$E404	ADC CHANNEL TWO
00840	E500	DACZRO EQU	\$E500	DAC CHANNEL ZERO
00850	*			
00860	*			
00870		*FUNCTION :TOLAN-A		
00880		*INPUTS :STATUS BUFFERS FROM EKG-EXEC		
00890		*OUTPUTS :DATA TO DISK		
00900		*CALLS :OUTPUT,FILHDR,KEYPD,OUTNCR		
00910		* KEYBD0,SAVEFL,START		
00920		*DESTROYS :ALL REGISTERS		
00930		*PURPOSE :TO COLLECT 3 CHANNELS OF EKG		
00940		* DATA AND STORE INTO MEMORY.		
00950	*			
00960	*			
00970	0100 0F	TOLAN1 SEI		
00980	0101 CE 0252	LDX #STR'ISG		
00990	0104 ED CA8F	JSR OUTNCR	"THIS MODULE SAMPLES TH...'."	
01000	0107 ED CA2C	JSR KEYPD	GET RESPONSE FROM CONSOLE	
01010	010A FE 3002	LDX FNDIRUF		
01020	010D 09	DEX		
01030	010E F6 03	LDA B 0,X		
01040	0110 C1 59	CMP B #'Y	IS INPUT YES?	
01050	0112 27 03	BEQ TOLAN1	YES. KEEP EXECUTING THIS ROUT	
01060	0114 7E 1D00	JMP START	NO. RIN TO EKG-EXEC	
01070	0117 CE 5441	TOLAN1 LDX #\$5441	FLAG CORRECTION TAPE (TA)	
01080	011A FF 1C98	STX CINTYP		
01090	011D BD 0233	JSR FILHDR	SET UP DATA FILE HEADER	
01100	0120 CE 02D9	IDX #IRQMSG		

01110	0123 BD CA8F	JSR	OUTNCR	"INSURE SUBJECT AND MSG..."
01120	0126 BD CA36	JSR	KEYPDO	
01130	012C FF 024E	STX	BITPTR	INITIALIZE BIT PTR START LOC
01140	013C FF 024E	LDA A	#4	STORE IN BUFFER
01150	012F 86 04	STA A	CALCNT	PICK UP COUNTER FOR TIME CAL
01160	0131 B7 0245	LDA A	#\$80	STORE IN BUFFER
01170	0134 86 80	STA A	BITPTR	INITIALIZE BIT POINTER TO LEF
01180	0136 B7 03C7	CLR A		INIT COMPRESS VAR
01190	0139 4F	STA A	FIRST	CLEAR FIRST SAMPLE FLAG
01200	013A B7 0251	STA A	XSLOPE	
01210	013D B7 03BC	STA A	YSLOPE	
01220	0140 B7 03BD	STA A	ZSLOPE	
01230	0143 B7 03BE	LDX	IRQVEC	PICK UP CURRENT IRQ VECTOR
01240	0146 FE FFFF	STX	VECSAV	SAVE IN BUFFER
01250	0149 FF 1C9D	LDX	#\$CALINT	GET INTR VECTOR ADDR FOR CAL
01260	014C CE 06AE TIMCAL	STX	IRQVEC	PUT IN VECTOR ADDRESS
01270	014F FF FFFF	LDA A	#0	INIT COUNTER FOR 256 TEST LOO
01280	0152 86 00	STA A	DONIST	STORE IN DONIST BUFFER
01290	0154 B7 0250	LDX	#0	CLEAR LOOP COUNTER
01300	0157 CE 0000	STX	IRQCNT	
01310	015A FF 024A	STX	IRQCNT+2	
01320	015D FF 024C	STX	DACZRO	PULSE INT ENABLE CIRCUIT
01330	0160 FF E500	BRA	SPOOL	BRA TO COUNTING LOOP
01340	0163 20 09	*		
01350		*****		
01360		*		
01370		*		
01380		*	BASIC TIMING LOOP FOR EFFICIENCY TEST	
01390		*		
01400		*****		
01410		*		
01420	0165 01	SPLCOOP	NOP	DELAY TO MATCH TIME
01430	0166 01		NOP	IT TAKES TO EXECUTE THE
01440	0167 01		NOP	INCREMENT OF BYTES 3&4
01450	0168 01		NOP	OF THE LOOP COUNTER
01460	0169 01		NOP	WHEN COUNT CARRIES TO
01470	016A 01		NOP	HIGH 2 BYTES OF 4 FFFF
01480	016B 7E 0181	JMP	SPO03	JUMP TO CONTINUE LOOP
01490	016E 0E	SPOOL	CLI	PREPARE FOR TIMER INT
01500	016F 3E		WAI	STOP PROCESSOR & WAIT
01510	0170 0E	SPO04	CLI	
01520	0171 FE 024C	LDX	IRQCNT+2	RET'ED FROM INT, INC COUNT
01530	0174 08	INX		
01540	0175 FF 024C	STX	IRQCNT+2	SAVE IT
01550	0178 26 E3	PNE	SPLCOOP	COUNT COME FFFF TO 0600?
01560	017A FE 024A	LDX	IRQCPT	YES, INC BYTES 3&4
01570	017D 08	INX		
01580	017E FF 024A	STX	IRQCPT	SAVE COUNT
01590	0181 B6 0250	LDA A	DONIST	IS DONE TEST SATISIFIED?
01600	0184 81 00	CMP A	#0	
01610	0186 26 E8	PNE	SPO04	NO, KEEP LOOPING
01620		*		
01630		*****		
01640		*		

01690	* END BASIC TIMING LOOP FOR EKG TEST		
01690	*	*****	
01690	*	*****	
01690 0188 0F	SEI		PREVENT SER TO MORE INTR
01700 0189 CE 4000	LDX #S4000		PICK UP INTR OFF WORD
01710 018C FF E500	STX DACZRO		DISABLE INTR FROM TIMER
01720 018F B7 E400	STA A ADCZRO		CLR ST6200 INT FLIP FLOP
01730 0192 01	NOP		
01740 0193 7A 0245	DEC CALCNT		DEC TIME CAL LOOP COUNTER
01750 0196 27 4F	BEQ SPDONE		IS COUNT 3? YES, GO SAVE DATA
01760 0198 B6 0245	LDA A CALCNT		NO. GET COUNT
01770 019B 81 03	CMP A #3		IS COUNT 3?
01780 019D 26 25	BNE MAINSP		NO. GO TO MAIN SPL LOOP
01790 019F FE 024C	LDX IRQCNT+2		YES. SAVE TIME LOOP COUNTER
01800 01A2 FF 0248	STX CALONE		
01810 01A5 CE 03C8	LDX #SAMPLE		NOW PUT SAMPLE ADDR IN IRQVEC
01820 01A8 FF FFF8	STX IRQVEC		
01830 01AB 86 00	LDA A #\$0		SET UP DONTST
01840 01AD B7 0250	STA A DONTST		
01850 01B0 7F 0251	CLR FIRST		RESET FIRST SMP FLAG
01860 01B3 CE 0000	LDX #0		CLR LOOP COUNT AND SAMPLE COU
01870 01B6 FF 024A	STX IRQCNT		
01880 01B9 FF 024C	STX IRQCNT+2		
01890 01BC FF 3494	STX SAMPLE		
01900 01BF FF E500	STX DACZRO		ENABLE TIMER INTERRUPTS
01910 01C2 20 AA	BRA SPOOL1		GO WAIT FOR INTERRUPT
01920 01C4 81 02	MAINSP CMP A #2		IS COUNT 2?
01930 01C6 26 10	BNE PARSAV		NO. SAVE PARAMETER FROM EKG
01940 01C8 86 55	LDA A #\$55		SET FIRST FLAG
01950 01CA B7 0251	STA A FIRST		
01960 01CD B7 0250	STA A DONTST		SET DONTST FOR SMPL COLLEC
01970 01D0 CF 0000	LDX #0		ENABLE INTR CLOCK
01980 01D3 FF E500	STX DACZRO		
01990 01D6 20 95	BRA SPOOL1		GO WAIT FOR INTERRUPT
02000 01D8 FE 024A	PARSAV LDX IRQCNT		SAVE LOOP COUNT
02010 01DB FF 3490	STX LOOPCPT		
02020 01DE FE 024C	LDX IRQCNT+2		
02030 01E1 FF 3492	STX LOOPCPT+2		
02040 01E1 7E 014C	JMP TIMCAL		GO EXECUTE ANOTHER TIME CM.
02050 01E7 B6 0249	SPDGNE LDA A CALONE+1		AVERAGE TWO TIME CM. HNS
02060 01EA BB 024D	ADD A IRQCNT+3		ADD LSB BYTES
02070 01ED B7 0247	STA A CALZRO+1		STORE IN BUFFER
02080 01FC B6 0248	LDA A CALONE		PICK UP MSB OF FIRST
02090 01F3 B9 024C	ADC A IRQCNT+2		ADD WITH CARRY
02100 01F6 B7 0246	STA A CALZRO		STORE IN BUFFER
02110 01F9 77 0246	ASR CALZRO		DIVIDE BY 2 TO ALIGN
02120 01FC 76 0247	ROR CALZRO+1		
02130 01FF 86 09	LDA A #8		SET UP COUNTER FOR SHIFT
02140 0201 77 0246	CNLSHF ASR CALZRO		NOW DIVIDE BY 2^6 = 64
02150 0204 76 0247	ROR CALZRO+1		LOOPS/INTERVAL
02160 0207 4A	DEC A		
02170 0208 26 F7	PNE CALSHF		8 SHIFTS YET?
02180 020A B6 0247	LDA A CALZRO+1		PICK UP SHIFTED RESULT

02190 020D B7 3496	STA A	LPCAL	STORE IN FILE BUFFER
02200 0210 BD 059E	JSR	DTACIT	NOW COUNT = OF DATA BITS FOR
02210 0212 11 1010	LXI	11,1010	FILE. PULL IN DATA
02220 0216 FF FFFS	SIX	11,FFFS	
02230 0219 CE 0336	LDX	#SAVMSG	
02240 021C BD CA8F	JSR	OUTNCR	"SAMPLING COMPLETE..."
02250 021F BD CA2C	JSR	KEYED	GET SAVE DECISION
02260 0222 FE 3002	LDX	ENDBUF	
02270 0225 09	DEX		
02280 0226 E6 00	LDA B	0,X	
02290 0228 C1 59	CMP B	#"Y	IS DECISION YES?
02300 022A 26 0C	BNE	EXJMP	NO. RIN TO EKG-EXEC
02310 022C BD 0240	JSR	SAVEFL	YES. SAVE THE FILE ON DISK
02320 022F CE 03A4	LDX	#SDONE	
02330 0232 BD CA8F	JSR	OUTNCR	
02340 0235 BD CA2C	JSR	KEYRD	
02350 0238 7E 1D00	EXJMP	JMP	START JUMP BACK TO EKG-EXEC
02360 *			
02370 *FUNC: RELOC JUMP			
02380 *INPUTS: A			
02390 *OUTPUTS: NONE			
02400 *CALLS: FILHDL,SAVEFL (EKG-EXEC ROUTINES)			
02410 *DESTROYS: A,B,CC,X			
02420 *PURPOSE : THIS ROUTINE ENABLES TOLAN-A TO PER RFLOCA			
02430 *WITHOUT WORRY OF CHANGING CALLS TO EKG-EXEC.			
02440 *			
02450 023E FE 1CA1	FILHDL	LDX	FILHLC GET FILHDL ADDRESS
02460 023E 6E 00		JMP	0,X JUMP TO IT
02470 *			
02480 0240 FE 1CA3	SAVEFL	LDX	SAVEFL GET SAVEFL ADDRESS
02490 0243 6E 00		JMP	0,X JUMP TO IT
02500 *			
02510 0245 0001	CALCNT	RMB	1 TEST CAL LOOP COUNTER
02520 0246 0002	CALZFO	RMB	2 AVG'D CAL LOOPS DURING DAT CO
02530 0248 0002	CALCFL	RMB	2 TEMP BUF FOR FRE CALLIC CAL
02540 024A 0004	IRQCTT	RMB	4 TEMP INTR LOOP COUNTER
02550 024E 0002	DUTYTR	RMB	2 BUFFER POINTER OF INTNT AVNL
02560 0250 0001	DONTST	RMB	1 DONE TEST FLG FOR MEM FULL
02570 0251 0001	FIRST	RMB	1 FIRST A/D SAMPL FLG
02580 *			
02590 0252 1A07	STRMSG	FDR	\$1A07
02600 0254 54	FCC		/THIS MODULE SAMPLES THE EKG /
02610 0270 49	FCC		/INPUT AND STORES THE DATA /
02620 028A 0D0A	FDB		\$0D0A
02630 028C 43	FCC		/COMPRESSED VIA MAJOR TOLAN-A/
02640 02A9 0D0A	FDB		\$0D0A,\$0D0A
02650 02AD 44	FCC		/DO YOU WISH TO EXECUTE THIS
02660 02C9 4D	FCC		/MODULE (Y OR N)/
02670 02D8 04	FCB		4
02680 02D9 1A	IRQMSG	FDR	\$1A
02690 02DA 49	FCC		/INSURE SUBJECT AND EKG LIVING /
02700 02F8 52	FCC		/READY!/
02710 02FD 0D0A	FDB		\$0D0A,\$0D0A,\$0D0A
02720 0304 50	FCC		/PRESS RETURN, THIS CLASS IS TYPICAL/

## TOLAN-8

02730 0326 20		FCC	/ ENABLE SWITCH./
02740 0335 04		FCB	4
02750 0336 1107	SAVING INT		\$1707
02760 0338 53		FCC	/SWITCHING CHANNEL, INPUT, OUTPUT
02770 0357 19		FCC	/INTERSOFT ENABLE SWITCH./
02780 036F 0D0A		FDB	\$0D0A,\$0D0A,\$0D0A
02790 0375 44		FCC	/DO YOU WISH TO SAVE THIS DATA ON /
02800 0396 44		FCC	/DISK (Y OR N)/
02810 03A3 04		FCB	4
02820 03A4 20	SDONE	FCC	/ PRESS RETURN/
02830 03B1 04		FCB	4
02840	*		
02850	*		
02860	*		END TOLAN-A
02870	*		
02880	*		
02890	*		FUNCTION :SAMPLE
02900	*		INPUTS :STATUS BUFFERS
02910	*		OUTPUTS :CCW/DIREC, ROUNDED DATA IN MEM BUFF
02920	*		CALLS :NOTHING
02930	*		PURPOSE :THIS ROUTINE SAMPLES THE EKG LEADS,
02940	*		ROUNDS THE VALUES TO 8 BITS (PROG 12)
02950	*		CALCULATES MAX,MIN,# OF BITS, SAMPLES
02960	*		ETC., AND SAVES THESE PARAM AND DATA
02970	*		TO DATA MEM FILE.
02980	*		
02990	*		
03000 03B2 0001	SIFBUF RMB	1	TEMP SHIFT BUFFER
03010 03B3 0001	TMPPA1 RMB	1	TEMP REG, MSB CH 0
03020 03B4 0001	TMPPA1 RMB	1	TEMP REG, LSB CH 0
03030 03B5 0001	TMPPA2 RMB	1	TMPP REG, MSB CH 1
03040 03B6 0001	TMPPA2 RMB	1	TEMP REG, LSB CH 1
03050 03B7 0001	TMPPA3 RMB	1	TEMP REG, MSB CH 2
03060 03B8 0001	TMPPA3 RMB	1	TEMP REG, LSB CH 2
03070 03B9 0001	XDATA RMB	1	TEMP X CH DTA BUF
03080 03BA 0001	YDATA RMB	1	TEMP Y CH DTA BUF
03090 03BB 0001	ZDATA RMB	1	TEMP Z CH DTA BUF
03100 03BC 0001	XSLOPE RMB	1	1ST DIFF X(N)-X(N+1)
03110 03BD 0001	YSLOPE RMB	1	1ST DIFF Y(N)-Y(N+1)
03120 03BE 0001	ZSLOPE RMB	1	1ST DIFF Z(N)-Z(N+1)
03130 03BF 0001	XEXP RMB	1	EXPECTED VAL OF X(N+1)
03140 03C0 0001	YEXP RMB	1	EXPECTED VAL OF Y(N+1)
03150 03C1 0001	ZEXP RMB	1	EXPECTED VAL OF Z(N+1)
03160 03C2 0001	XCCFL RMB	1	DIFF X(N+1)-E(X(N+1))
03170 03C3 0001	YACCEL RMB	1	DIFF Y(N+1)-E(Y(N+1))
03180 03C4 0001	ZACCEL RMB	1	DIFF Z(N+1)-E(Z(N+1))
03190 03C5 0001	STPLIG INT	1	I HAVE JUST GAINED 100%
03200 03C6 0001	ICIT INT	1	TIME CUT VAR (<127)
03210 03C7 0001	BITPR RMB	1	BIT POINTER FOR SET & RESET
03220	*		
03230 03C8 FF 3494	SAMPLE LDX	SAMENO	GET CUR SAMPLE COUNT
03240 03CB 08	INX		
03250 03CC FF 3494	STX	SAMPIO	SAVE STACK INT
03260 03CF FF 1C96	SIS	STKSAV	SAVE STACK INT

03270 03D2 8E 3280	LDS	#\$3280	INITIALIZE STACK IN UNSED MEM
03280 03D4 F0 0001	LDA A	#1	CK FIRST FLAG
(a)			IF LT 100
03300 03DD F0 05C8	JSR	SAMPL0	YLS. SML & RLS
03310 03DD F0 024E	LDX	BUFFPTR	NOW STR FIRST DATA VALUES
03320 03E0 F6 03F9	LDA A	XDATA	
03330 03E3 F7 00	STA A	0,X	
03340 03E5 F8	INX		
03350 03E6 F6 03FA	LDA A	YDATA	
03360 03E9 F7 00	STA A	0,X	
03370 03FB F8	INX		
03380 03EC F6 03FB	LDA A	ZDATA	
03390 03EF F7 00	STA A	0,X	
03400 03F1 F8	INX		
03410 03F2 FF 024E	STX	BUFFPTR	UPDATE NEW BUFFPTR
03420 03F5 F6 55	LDA A	#\$55	SET STRFLG & FIRST FLAG
03430 03F7 F7 03C5	STA A	STRFLG	
03440 03FA F7 0251	STA A	FIRST	
03450 03FD F6 08	LDA A	#8	UPDATE X,Y,Z BIT CTRS FOR IN
03460 03FF F7 34E2	STA A	XBITS+2	
03470 0402 F7 34E5	STA A	YBITS+2	
03480 0405 F7 34E8	STA A	ZBITS+2	
03490 0408 7F 0250	CLR	DONTST	CLEAR DONTST FLAG FOR 1 SMPL
03500 040B F6 1C96	LDS	STKSAV	
03510 040E F6 3B	RTI		
03520 040F 7D 03C5 COMPRES	TST	STRFLG	HAS DATA JUST BEEN STORED ?
03530 0412 27 05	BEQ	COMPRI	NO. KEEP COUNTING TIME
03540 0414 F6 01	LDA A	#1	YES. SET TIME COUNT TO 1
03550 0416 F7 03C6	STA A	ICNT	
03560 0419 F6 03F9 COMPRI	LDA A	XDATA	GET X. VLU
03570 041C F6 03IC	ADD A	XSL0P0	ADD {X(N-1)-X(N-2)}
03580 041F F7 03FF	STA A	XEXP	CREATE EXPECTED VAL OF X(N+1)
03590 0422 F6 03FA	LDA A	YDATA	GET LAST Y SML VALUE
03600 0425 BB 03FD	ADD A	YSLOPE	ADD {(Y(N-1)-Y(N-2))}
03610 0428 F7 03C0	STA A	YEXP	CREATE EXPECTED VAL OF Y(N+1)
03620 042B F6 03FB	LDA A	ZDATA	GET LAST Z SML VALUE
03630 042E BB 03FB	ADD A	ZSLOPE	ADD {(Z(N-1)-Z(N-2))}
03640 0431 F7 03C1	STA A	ZEXP	CREATE EXPECTED VAL OF Z(N+1)
03650 0434 F6 05C8	JSR	SAMPL0	NOW GO GET NEXT SML.. N=N+1
03660 0437 F6 03F9	LDA A	XDATA	CALC DIF X(N)-L(X(N))
03670 043A F0 03FF	SUB A	YEXP	
03680 043D F7 03C2	STA A	XACCEL	SAVE DIF IN VR
03690 0440 F6 03FA	LDA A	YDATA	CALC DIF Y(N)-L(Y(N))
03700 0443 F0 03C0	SUB A	YEXP	
03710 0446 F7 03C3	STA A	YACCEL	SAVE DIF IN VR
03720 0449 F6 03FB	LDA A	ZDATA	CALC DIF Z(N)-L(Z(N))
03730 044C F0 03C1	SUB A	ZEXP	
03740 044F F7 03C4	STA A	ZACCEL	SAVE DIF IN VR
03750 0452 F6 03C2	LDA A	XACCEL	ARE ALL EXTRAPOLATIONS GOOD?
03760 0455 26 21	FNE	COMPRI2	IF NO, NEED TO STR ALL DIFT'S
03770 0457 F6 03C3	LDA A	YACCEL	
03780 045A 26 1C	FNE	COMPRI2	
03790 045C F6 03C2	LDA A	YACCEL	
03800 045F 26 17	FNE	COMPRI2	

## TOLAN-8

03810 0461 B6 03C6	LDA A	ICNT	UPDATE X N TIME COUNTER
03820 0464 AC	INC A		
03830 0467 2F 71	CMP A		TEST N & 1 (N>1)?
03840 0469 B7 03C6	BLD	XCNTS	YEP. COUNT THE PULL & ADD IC
03850 046C 7F 03C5	STA A	ICNT	NO. INC ICNT & RTI
03860 0470 86 55	CLR	SIRFLG	RESET I HAVE JUST STRED FLG
03870 047F 86 55	LDA A	#\$55	
03880 0471 B7 0250	STA A	DONIST	KEEP DONE TEST SET
03890 0474 BE 1C96	LDS	STKSAV	RETRIEV STACK
03900 0477 3B	RTI		
03910 *			
03920 0478 CE 34B0	COMPR2	IDX	UPDATE X BIT CNTR
03930 047E B6 03C2	LDA A	XACCEL	
03940 047E BD 0568	JSR	BITCNT	
03950 0481 CE 3600	LDX	#XPDFM	UPDATE CH X FREQ OF OCCUR HIS
03960 0484 B6 03C2	LDA A	XACCEL	
03970 0487 BD 066B	JSR	PDFSTR	
03980 048A B6 03C2	LDA A	XACCEL	
03990 048D B6 03C2	LDA A	XACCEL	
04000 0490 BB 03FC	ADD A	XSLOPE	CALC & SAVE X(N-1)-X(N-2)
04010 0493 B7 03FC	STA A	XSLOPE	
04020 0496 CE 34B3	LDX	#YBITS	UPDATE Y BIT CNTR
04030 0499 B6 03C3	LDA A	YACCEL	
04040 049C BD 0568	JSR	BITCNT	
04050 049F CE 38C0	LDX	#YPDFM	UPDATE CH Y FREQ OF OCCUR HIS
04060 04A2 B6 03C3	LDA A	YACCEL	
04070 04A5 BD 066B	JSR	PDFSTR	
04080 04A8 B6 03C3	LDA A	YACCEL	
04090 04AB B6 03C3	LDA A	YACCEL	
04100 04AE BB 03ED	ADD A	YSLOPE	CALC & SAVE Y(N-1)-Y(N-2)
04110 04B1 B7 03FD	STA A	YSLOPE	
04120 04B4 CE 34B6	LDX	#ZBITS	UPDATE Z BIT CNTR
04130 04B7 B6 03C4	LDA A	ZACCEL	
04140 04FA BD 0568	JSR	BITCNT	
04150 04FD CE 100	LDX	#ZPDFM	UPDATE CH Z FREQ OF OCCUR HIS
04160 04C0 B6 03C4	LDA A	ZACCEL	
04170 04C3 FD 066B	JSR	PDFSTR	
04180 04C6 B6 03C4	LDA A	ZACCEL	
04190 04C9 B6 03C4	LDA A	ZACCFI	
04200 04CC BB 03FE	ADD A	ZSLOPE	CALC & SAVE Z(N-1)-Z(N-2)
04210 04CF B7 03FE	STA A	ZSLOPE	
04220 04D2 CE 34B9	LDX	#TBITS	NOW UPDATE TIME BIT CTR
04230 04D5 86 06	LDA A	#6	
04240 04D7 FD 0568	JSR	BITCNT	
04250 04DA B6 03C2	COMPR3	LDA A	PICK UP YACCEL FOR
04260 04DD BD 051B	JSR	COMSTR	GO CODE & STORE IT
04270 04E0 B6 03C3	LDA A	YACCEL	PICK UP YACCFI FOR
04280 04E3 FD 051B	JSR	COMSTR	GO CODE & STORE IT
04290 04E6 B6 03C4	LDA A	ZACCEL	PICK UP ZACCEL FOR
04300 04E9 FD 051B	JSR	COMSTR	GO CODE & STORE IT
04310 04LC B6 03C6	LDA A	ICNT	NEW GET ICNT FOR TIME CTR
04320 04LF CB 3F00	IDX	TPPDF	UPDAT TIME BIT CTR
04330 04F2 FD 066B	JSR	PDFSTR	
04340 04F5 FD 056A	JSR	RESET	SET TO 0 IN TIME BIT CTR

AD-A100 799

AIR FORCE INST OF TECH WRIGHT-PATTERSON AFB OH SCHOOL—ETC F/S 6/5  
ANALYSIS AND PERFORMANCE EVALUATION OF ELECTROCARDIOGRAM DATA C—ETC(U)  
DEC 80 M D TOWNSEND

UNCLASSIFIED AFIT/GE/EE/80D-46

NL

3 OF 3  
AFIT/EE/80D-46  
100/99

END  
DATE  
FILED  
7-31  
DTIC

04350 04F8 79 03C6		ROL	ICNT	ROL 7 BIT ICNT TO MS POSITION
04360 04FB 86 08		LDA A	#8	SET CNT FOR TIME STORE LOOP
04370 04FD 4A	CON'PR4	DEC A		DECREMENT LOOP CNT
04380 04FE 27 0F		BEQ	CON'PR6	TIME STORED YET?
04390 0500 79 03C6		ROL	ICNT	NO. ROTATE TIME WORD ONCE LEF
04400 0503 25 05		BCS	COMPR5	1 IN CARRY?
04410 0505 BD 055A		JSR	RESET	NO. PUT 0 IN MEM FILE
04420 0508 20 F3		BRA	COMPR4	
04430 050A BD 0534	COMPR5	JSR	SET	YES. PUT 1 IN MEM FILE
04440 050D 20 EE		BRA	COMPR4	
04450 050F 86 55	COMPR6	LDA A	#\$55	NOW SET "I HAVE JUST STORED"
04460 0511 B7 03C5		STA A	STRFLG	SET DONTST FLAG
04470 0514 B7 0250		STA A	DONTST	
04480 0517 BE 1C96		LDS	STKSAV	
04490 051A 3B		RTI		
04500 *				
04510 051B BD 055A	COMSTR	JSR	RESET	PUT 0 IN MEM FILE FOR DELIM
04520 051E 4D		TST A		NOW TEST STATUS OF A
04530 051F 2A 06		BPL	COMST1	IS IT +?
04540 0521 40		NEG A		NO. GET 2'S COMP FOR MAG
04550 0522 BD 0534		JSR	SET	NO. STORE 1 FOR SIGN BIT
04560 0525 20 03		BRA	COMST2	
04570 0527 BD 055A	COMST1	JSR	RESET	YES. STORE 0 FOR SIGN BIT
04580 052A 4D	COMST2	TST A		NOW CNT DOWN & STORE DATA
04590 052B 27 06		BEQ	COMST3	CNT 0?
04600 052D BD 0534		JSR	SET	NO. STORE 1 TO MEM FILE
04610 0530 4A		DEC A		DEC CNTR
04620 0531 20 F7		BRA	COMST2	
04630 0533 39	COMST3	RTS		
04640 *				
04650 *				
04660 *FUNC: SET				
04670 *INPUTS: BUFPTR,BITPTR				
04680 *OUTPUTS: BIT SET PT'D TO BY ABOVE				
04690 *CALLS: NOTHING				
04700 *DESTROYS: B,X,CC				
04710 *PURPOSE: THIS ROUTINE SETS THE BIT POINTED TO				
04720 * BY BITPTR & BUFPTR AND THEN UPDATES THESE COUNTERS				
04730 *				
04740 0534 FE 024E	SET	LDX	BUFPTR	GET CUR MEM WORD FOR DATA STO
04750 0537 F6 03C7		LDA B	BITPTR	GET BIT IN THAT WORD TO BE SE
04760 053A EA 00		ORA B	0,X	SET THE BIT
04770 053C E7 00		STA B	0,X	SAVE IN MEM FILE
04780 053E 0C	SET0	CLC		CLR CARRY FOR PROPER LEFT ROT
04790 053F 76 03C7		ROR	BITPTR	ROLL BITPTR ONCE RIGHT
04800 0542 24 15		BCC	SETRTS	WAS BITPTR DOWN TO 1ST BIT?
04810 0544 76 03C7		ROR	BITPTR	YES. IN BUFPTR \$ SET BITPTR T
04820 0547 08		INX		
04830 0548 FF 024E		STX	BUFPTR	
04840 054B F6 024E		LDA B	BUFPTR	
04850 054E C1 80		CMP B	#\$80	
04860 0550 26 07		ENE	SETRTS	
04870 0552 7F 0250		CLR	DONTST	YES. STOP DATA COLLEC & RTI
04880 0555 BE 1C96		LDS	STKSAV	

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04890 C558 3B          RTI
04900 0559 39          SETRTS RTS
04910
04920
04930 *FUNC: RIGET
04940 *INPUTS: BUFPTR,BITPTR
04950 *OUTPUTS: BIT RESET IN MEM FILE PT'ED TO BY ABOVE
04960 *CALLS: NOTHING
04970 *DESTROYS: B,X,CC
04980 *PURPOSE: THIS ROUTINE RESETS THE BIT IN THE MEMFILE
04990 * POINTED TO BY BUFPTR & BITPTR & UPDATES THESE COUN
05000 055A FE 024E RESET LDX   BUFPTR
05010 055D F6 03C7      LDA B   BITPTR
05020 0560 53           COM B
05030 0561 E4 00         AND B  0,X     RESET BIT IN MEM WORD
05040 0563 E7 00         STA B  0,X     SAVE BIT IN MEM FILE
05050 0565 7E 053E      JMP    SETO    UPDATE CNTR'S AND CK IF MEM F
05060 *
05070 *FUNC: BITCNT
05080 *INPUTS: ACCEL VLU IN A, BIT CNTR IN X
05090 *OUTPUTS: UPDATED BIT COUNT
05100 *CALLS:NOTHING
05110 *DESTROYS:A,B,CC
05120 *PURPOSE: THIS ROUTINE INC 3 BYTE BIT COUNTER WITH
05130 * NUMBER OF BIT USED TO STORE X,Y,Z,& TIME.
05140 *
05150 0568 2A 01         BITCNT BPL   BITCN1  IS ACCEL +
05160 056A 40           NEG A
05170 056B 4C           BITCN1 INC A
05180 056C 4C           INC A
05190 056D C6 00         LDA B  #0
05200 056F AB 02         ADD A  2,X
05210 0571 A7 02         STA A  2,X
05220 0573 86 00         LDA A  #0
05230 0575 E9 01         ADC B  1,X
05240 0577 E7 01         STA B  1,X
05250 0579 C6 08         LDA B  #8
05260 057B A9 00         ADC A  0,X
05270 057D A7 00         STA A  0,X
05280 057F 86 00         LDA A  #0
05290 0581 FB 34BF      ADD B  ACELCT+3 NO. INC ACELCT COUNT
05300 0584 F7 34BF      STA B  ACELCT+3
05310 0587 C6 00         LDA B  #0
05320 0589 B9 34BE      ADC A  ACELCT+2
05330 058C B7 34BE      STA A  ACELCT+2
05340 058F 86 00         LDA A  #0
05350 0591 F9 34BD      ADC B  ACELCT+1
05360 0594 F7 34BD      STA B  ACELCT+1
05370 0597 B9 34BC      ADC A  ACELCT
05380 059A B7 34BC      STA A  ACELCT
05390 059D 39           BITRTS RTS
05400 *
05410 *FUNC: DTACNT
05420 *INPUTS: NUM ON SMPLS COLLECTED

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05430 *OUTPUTS: BITCNT=TO SMPLS*24
05440 *CALLS: NOTHING
05450 *DESTROYS: A,B,X,CC
05460 *PURPOSE: TO CALC NUM OF BITS STORED DURING COLLECTI
05470 *
05480 059E FE 3494 DTACNT LDX SAMPNO
05490 05A1 86 13 DTACN1 LDA A #24
05500 05A3 C6 00 LDA B #0
05510 05A5 BB 34AF ADD A DTABIT+3
05520 05A8 B7 34AF STA A DTABIT+3
05530 05AB 86 00 LDA A #0
05540 05AD F9 34AE ADC B DTABIT+2
05550 05B0 F7 34AE STA B DTABIT+2
05560 05B3 C6 00 LDA B #0
05570 05B5 B9 34AD ADC A DTABIT+1
05580 05B8 B7 34AD STA A DTABIT+1
05590 05BB F9 34AC ADC B DTABIT
05600 05BE F7 34AC STA B DTABIT
05610 05C1 09 DEX
05620 05C2 8C 0000 CPX #0
05630 05C5 26 DA BNE DTACN1
05640 05C7 39 RTS
05650 *
05660 *FUNC: SAMPL0
05670 *INPUTS: DATA VIA A/D CONV
05680 *OUTPUTS: ROUNDED DATA TO XDTA,YDTA,ZDTA
05690 *CALLS:PDFSTR (UPDATE PDF HISTOGRAMS)
05700 *DESTROYS:A,B,X,CC
05710 *PURPOSE: THIS ROUTINES SAMPLES THE A/D AND THEN
05720 * ROUNDS THE DATA TO 8 BITS FROM 12. THE DATA
05730 * IS PLACED IN XDTA,YDTA,ZDTA AND THE COLLEC
05740 * STAT ARE UPDATED.
05750 *
05760 05C8 CE E400 SAMPL0 LDX #ADCZRO NOW PULSE A/D TO START CONV
05770 05CB A7 00 STA A 0,X ON CHANNEL 0
05780 05CD 01 NOP
05790 05CE A6 00 LDA A 0,X
05800 05D0 E6 01 LDA B 1,X
05810 05D2 B7 03B3 STA A TEMPAL
05820 05D5 F7 03B4 STA B TEMPB1
05830 05D8 08 INX
05840 05D9 08 INX
05850 05DA A7 00 STA A 0,X NOW PULSE A/D TO CONV
05860 05DC 01 NOP ON CHANNEL 1
05870 05DD A6 00 LDA A 0,X
05880 05DF E6 01 LDA B 1,X
05890 05E1 B7 03B5 STA A TEMPAL
05900 05E4 F7 03B6 STA B TEMPB2
05910 05E7 08 INX
05920 05E8 08 INX
05930 05E9 A7 00 STA A 0,X NOW PULSE A/D TO CONV
05940 05EB 01 NOP ON CH 2
05950 05EC A6 00 LDA A 0,X
05960 05EE E6 01 LDA B 1,X

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05970	05F0	B7	03B7	STA A	TEMPA3		
05980	05F3	F7	03B8	STA B	TEM1PB3		
05990	05F6	CE	03B3	LDX	#TEM1PAL		
06000	05F9	A6	00	SAMPL1	LDA A	0,X	
06010	05FB	E6	01		LDA B	1,X	
06020	05FD	47			ASR A		
06030	05FE	56			ROR B		
06040	05FF	47			ASR A		
06050	0600	56			ROR B		
06060	0601	47			ASR A		
06070	0602	56			ROR B		
06080	0603	47			ASR A		
06090	0604	56			ROR B		
06100	0605	F7	03B2	STA B	SHFBUF	SAVE 8 BIT RESULT OF SHIFT	
06110	0608	86	00	LDA A	#0	AND ADD CARRY OUT OF LAST	
06120	060A	B9	03B2	ADC A	SHFBUF	SHIFT ROUNDING UP OR DOWN	
06130	060D	36		PSH A		SAVE TO MEM FILE BUFFER VIA S	
06140	060E	08			INX		
06150	060F	08			INX		
06160	0610	8C	03B9	CPX	#TEMPA3+2	CHANNEL Z ROUNDED TO 8 BITS	
06170	0613	26	E4	BNE	SAMPL1	NO. GO SAMPLE NEXT CHANNEL	
06180	0615	32		PUL A		GET Z DATA	
06190	0616	B7	03BB	STA A	ZDATA		
06200	0619	FE	3494	LDX	SAMPNO	GET CUR SAMPLE COUNT	
06210	061C	B1	3497	CMP A	MAXZ	IS CUR Z MAX CVE SAMPLE SET?	
06220	061F	2F	06	BLE	SPZMIN	NO. GO CHECK FOR MIN	
06230	0621	B7	3497	STA A	MAXZ	YES. KEEP CUR VLU	
06240	0624	FF	3498	STX	MAXZLO	KEEP CUR SAMPLE NUM	
06250	0627	B1	349A	CMP A	MINZ	IS CUR Z MIN OVER SAMPLE SET?	
06260	062A	2C	06	BGE	SPYMAX	NO. GO CHECK FOR Y MAX	
06270	062C	B7	349A	STA A	MINZ	YES. KEEP CUR VLU	
06280	062F	FF	349B	STX	MINZLO	KEEP CUR SAMPLE NUM	
06290	0632	32		SPYMAX	PUL A	GET CUR Y DATA	
06300	0633	B7	03BA	STA A	YDATA		
06310	0636	B1	349D	CMP A	MAXY	IS CUR Y MAX OVER SAMPLE SET	
06320	0639	2F	06	BLE	SPYMIN	NO. GO CHEK FOR MIN	
06330	063B	B7	349D	STA A	MAXY	YES. KEEP CUR Y VLU	
06340	063E	FF	349E	STX	MAXYLO	KEEP CUR SAMPLE NUM	
06350	0641	B1	34A0	SPYMIN	CMP A	MINY	IS CUR Y MIN OVER SAMPLE SFT?
06360	0644	2C	06	BGE	SPXMAX	NO. GO CHEK FOR X MAX	
06370	0646	B7	34A0	STA A	MINY	YES. KEEP CUR Y VLU	
06380	0649	FF	34A1	STX	MINYLO	KEEP CUR SAMPLE NUM	
06390	064C	32		SPXMAX	PUL A	GET CUR X DATA	
06400	064D	B7	03B9	STA A	XDATA		
06410	0650	B1	34A3	CMP A	MAXX	IS CUR X MAX OVER SAMPLE SET?	
06420	0653	2F	06	BLE	SPXMIN	NO. GO CHECK X MIN	
06430	0655	B7	34A3	STA A	MAXX	YES. KEEP CUR X VLU	
06440	0658	FF	34A4	STX	MAXXLO	KEEP CUR SAMPLE VLU	
06450	065B	B1	34A6	SPXMIN	CMP A	MINX	IS CUR X MIN OVER SAMPLE SET?
06460	065E	2C	06	BGE	MMDONE	NO. EXIT MAX,MIN UPDATE	
06470	0660	B7	34A6	STA A	MINX	YES. KEEP CUR X VLU	
06480	0663	FF	34A7	STX	MINXLO	KEEP CUR SAMPLE COUNT	
06490	0666	39		MMDONE	RTS		
06500			*				

12 BIT TO 8 BIT ROUNDED CONV

## TOLAN-8

06510 0667 0002	TEMPDF RMB	2	TEMP WORKING BUFFER
06520 0669 0002	TEMPST RMB	2	TEMP STACK SAVE BUFFER
06530 *			
06540 066E BF 0669	PDFSTR STS	TEMPST	SAVE STACK PTR
06550 066E FF 0667	STX	TEMPDF	SAVE PDF PTR
06560 0671 16	TAB		
06570 0672 BB 0668	ADD A	TEMPDF+1	NOW CALCULATE ADDRESS IN INC
06580 0675 B7 0668	STA A	TEMPDF+1	IN THE PDF MEM BUFF
06590 0678 86 00	LDA A	#0	
06600 067A C5 80	BIT B	#\$80	IS THE DATA VLU NEG
06610 067C 2B 05	BMI	PDFST1	YES. SBC MSB VS ADC
06620 067E B9 0667	ADC A	TEMPDF	NO. ADC MSB
06630 0681 20 03	BRA	PDFST2	AND STORE
06640 0683 B2 0667	SBC A	TEMPDF	VLU IS NEG. SBC
06650 0686 B7 0667	STA A	TEMDFD	AND STORE
06660 0689 17	TBA		RECOVER VIU AND ADD (SUB) AGA
06670 068A BB 0668	ADD A	TEMPDF+1	FOR PROPER ADDR CALC
06680 068D B7 0668	STA A	TEMPDF+1	
06690 0690 86 00	LDA A	#0	
06700 0692 C5 80	BIT B	#\$80	IS VLU NEG?
06710 0694 2B 05	BMI	PDFST3	YES. SBC MSB VS ADC MSB
06720 0696 B9 0667	ADC A	TEMPDF	NO. ADC MSB
06730 0699 20 03	BRA	PDFST4	
06740 069B B2 0667	SBC A	TEMPDF	VLU IS NEG. SBC
06750 069E B7 0667	STA A	TEMPDF	AND STORE
06760 06A1 FE 0667	LDX	TEMPDF	NO LOAD CALC ADDR FOR INDEX
06770 06A4 AE 00	LDS	0,X	GRAB VLU IN CALC ADDR
06780 06A6 31	INS		INC IT
06790 06A7 AF 00	STS	0,X	AND STORE IT BACK IN BUFFER
06800 06A9 BE 0669	LDS	TEMPST	RECOVER STACK POINTER
06810 06AC 17	TBA		
06820 06AD 39	RTS		
06830 *			
06840	*FUNC: CALINT		
06850	*INPUTS: NONE		
06860	*OUTPUTS: NONE		
06870	*CALLS: NOTHING		
06880	*DESTROYS :NO REG (INTR HANDLER)		
06890	*PURPOSE: THIS ROUTINE IS USED FOR CALIBRATING		
06900	* THE MAX NUMBER OF LOOPS POSSIBLE DURING		
06910	* AN INTERRUPT PERIOD. THIS ROUTINE UPDATES THE		
06920	* DONTST FLAG AND RESETS THE ST6800 INT FLIP FLOP		
06930	* AND THEN RETURNS.		
06940	*		
06950 06AE 7C 0250	CALINT INC	DONTST	INCREMENT DONE TEST FLAG
06960 06B1 B7 E400	STA A	ADCZRO	CLR ST6800 INTR FLIP FLOP
06970 06B4 01	NOP		
06980 06B5 3B	RTI		
06990	*		
07000	* END OF TOLAN-A OVERLAY ROUTINES		
07010	*		
07020	END		

```

00030      *
00040      *      PROGRAM NAME: DECPRS
00050      *      AUTHOR : CAPT. MEL TOWNSEND
00060      *      VERSION : 1.3
00070      *      VERSION DATE : 29 NOV 80
00080      *
00090      ****
00100      *
00110      *      PROGRAM DESCRIPTION
00120      *
00130      *      THIS PROGRAM PERFORMS THE DECOMPRESSION OPERATI
00140      *      ON THE DATA COMPRESSED BY THE TOLAN-A,TOLAN-B,DOWE
00150      *      ETC. THE ROUTINE ASSUMES THAT THE DATA IS IN MEMO
00160      *      AND READS THE COMPRESSION TYPE FROM THE MEMORY FIL
00170      *      BUFFER HEADER. AFTER PROMPT TO THE TERMINAL, THIS
00180      *      PROGRAM BEGINS DECOMPRESSION AND OUTPUTS THE DECOM
00190      *      DATA TO AN ANALOG DISPLAY DEVICE (IE OSCILLISCOPE)
00200      *      D/A CONVERTER 0. THE ROUTINE REQUIRES AN INTERRUPT
00210      *      CLOCK VIA THE A/D BOARD (ST16800) AND THE ASSUMPTIO
00220      *      IS MADE THAT THE INTERRUPT CLOCK FREQUENCY IS ADJU
00230      *      TO THE SAMPLE RATE AT WHICH THE DATA WAS TAKEN.
00240      *
00250      ****
00260      *
00270      *      START OF DECPRS
00280      *
00290      ****
00300      *
00310 0500      ORG    $0500    PROGRAM ORIGIN
00320      *
00330      OPT     O      ASSB OPT. LIST ASSEMBLY
00340      OPT     NOG    ASSB OPT. SUPPRESS FCC LIST
00350      *
00360      ****
00370      *
00380      *      LABLE DESCRIPTIONS
00390      *
00400      *      SUPPORT SUBROUTINES
00410      *
00420  CA8F  OUTNCK EQU    $CA8F    EPROMDOS. OUTPUT STRING
00430  CA2C  KEYBD0 EQU    $CA2C    EPROMDOS. INPUT ALPH STRING
00440  0100  DISPLA EQU    $0100    EKG-EXEC. TERMINAL INTFC DRIV
00450      *
00460      *      DATA BUFFERS
00470      *
00480  0020  CHNLBF EQU    $0020    CHNL SELECT FLG FRM DISPLAY
00490  3C00  XINIT EQU    $3C00    INITIAL COND CH X
00500  3C01  YINIT EQU    $3C01    INITIAL COND CH Y
00510  3C02  ZINIT EQU    $3C02    INITIAL COND CH Z
00520  3002  ENDBUF EQU    $3002    END-OF-STRING FRM TERM INPUT
00530  FFF8  IRQVEC EQU    $FFF8    MASKABLE INTR JMP VECTOR ADDR
00540  3400  HDRSTR EQU    $3400    MEM FILE HEADER ADDR
00550  3C00  SECZRO EQU    $3C00    DATA START LOCATION
00560      *

```

C0570 \* HARDWARE ADDRESSES  
 00580 \*  
 00590 E500 DACZRO EQU \$E500 D/A CH 0  
 00600 E400 ADCZRO EQU \$E400 A/D CH 0  
 00610 \*  
 00620 \*  
 00630 \*FUNCTION: DECPRS  
 00640 \*INPUTS : CHNL SELECT VIA CHNLBF  
 00650 \*OUTPUTS : DECRPRSED DATA VIA D/A  
 00660 \*CALLS : OUTNCR,KEYBDO,DISPLAY  
 00670 \*DESTROYS (REG) : A,B,X,CC  
 00680 \*PURPOSE : THIS IS THE SET UP ROUTINE WHICH INITIALI  
 00690 \* THE PROGRAM CONSTANTS AND, DETERMINES WHAT TYPE OF  
 00700 \* COMPRESSION WAS USED, AND ENABLES RECONSTRUCTION  
 00710 \* INTERRUPTS.  
 00720 \*  
 00730 0500 0F DECPRS SEI  
 00740 0501 FE 3400 LDX HDRSTR GET CMPRS TYPE FROM MEM FILE  
 00750 0504 FF 0538 STX NOTDNL SIR CPR TYPE IN ERROR MSG  
 00760 0507 8C 5441 CPX #\$5441 IS IT TOLAN-A?  
 00770 050A 27 61 BEQ TADEC P YES. GO DECPRS IT.  
 00780 050C CE 0518 LDX #NOTDNE NO. CPRS ROUTINE NOT DONE  
 00790 050F BD CA8F JSR OUTNCR PRINT ERROR MSG & RETURN  
 00800 0512 BD CA2C JSR KEYBDO  
 00810 0515 7E 0100 JMP DISPLA RETURN TO DISPLAY ROUTINE  
 00820 \*  
 00830 0518 1A07 NOTDNE FDB \$1A07,\$0D0A  
 00840 051C 44 FCC /DECOMPRESSION FOR FILE TYPE /  
 00850 0538 0002 NOTDNL RMB 2  
 00860 053A 20 FCC / NOT YET IMPLEMENTED./  
 00870 054F 0D0A FDB \$0D0A,\$0D0A  
 00880 0553 50 FCC /PRESS ANY KEY TO CONTINUE/  
 00890 056C 04 FCB 4  
 00900 \*  
 00910 056D 96 20 TADEC P LDA A CHNLBF GET WHICH CHANNEL TO BE DECPR  
 00920 056F 81 00 CMP A #0 IS IT CHNL X?  
 00930 0571 26 07 TADEC0 BNE TADEC1 NO. CHECK IF Y.  
 00940 0573 C6 58 LDA B #'X PUT ASCII X IN ERROR MSG  
 00950 0575 F7 0846 STA B CHNASC  
 00960 0578 20 10 BRA TADEC3  
 00970 057A 81 01 TADEC1 CMP A #1 IS IT CH Y?  
 00980 057C 26 07 BNE TADEC2 NO. MUST BE Z.  
 00990 057E C6 59 LDA B #'Y PUT ASCII Y IN ERROR MSG  
 01000 0580 F7 0846 STA B CHNASC  
 01010 0583 20 05 BRA TADEC3  
 01020 0585 C6 5A TADEC2 LDA B #'Z PUT ASCII Z IN ERROR MSG  
 01030 0587 F7 0846 STA B CHNASC  
 01040 058A CE 0811 TADEC3 LDX #GOMSG GET PROMPT MSG LOC  
 01050 058D BD CA8F JSR OUTNCR NOW PRINT PRMPT  
 01060 0590 BD CA2C JSR KEYBDO  
 01070 0593 FE 3002 LDX ENDBUF GET ANSWER FROM TERMINAL  
 01080 0596 09 DEX  
 01090 0597 E6 00 LDA B 0,X  
 01100 0599 C1 4E CMP B #'N WAS IT NO?

01110 059B 26 03	BNE	TADEC4	YES. EXIT ROUTINE
01120 059D 7E 0100	JMP	DISPLA	
01130 05A0 CP 3C03 TADEC4	LDX	*SCHM+3	POSITION PTR TO DTA. STR
01140 05A3 FF 0611	STX	PUPTR	
01150 05A6 86 80	LDA A	#\$80	INITIALIZE BITPTR TO LEFT BIT
01160 05A8 B7 0613	STA A	BITPTR	
01170 05AB 7F 0617	CLR	XSLOPE	INIT SLOPES TO ZERO (2ND INIT)
01180 05AE 7F 0618	CLR	YSLOPE	
01190 05B1 7F 0619	CLR	ZSLOPE	
01200 05B4 B6 3C00	LDA A	XINIT	GET FIRST INIT COND
01210 05B7 B7 0614	STA A	XDTA	& PUT IN DATA BUF
01220 05BA B6 3C01	LDA A	YINIT	
01230 05BD B7 0615	STA A	YDTA	
01240 05C0 B6 3C02	LDA A	ZINIT	
01250 05C3 B7 0616	STA A	ZDTA	
01260 05C6 7F 0629	CLR	TARFLG	CLR DECPRSN FLG
01270 05C9 7F 061A	CLR	DELT	INIT TIME COUNTER
01280 05CC FE FFF8	LDX	IRQVEC	PICK UP CURRENT INT VEC
01290 05CF FF 061D	STX	VECSAV	SAVE IT IN TEMP BUFFER
01300 05D2 CE 062F	LDX	#TARECN	PUT NEW INT VEC IN
01310 05D5 FF FFF8	STX	IRQVEC	
01320 05D8 86 AA	LDA A	#\$AA	SET DONTST FLAG
01330 05DA B7 062A	STA A	DONTST	
01340 05DD CE 0000	LDX	#0	
01350 05E0 FF E500	STX	DACZRO	SEND ENABLE PLS TO INT CIRCUIT
01360 *			
01370 05E3 0E	TAINTR	CLI	ENABLE CPU FOR RECONS INTR'S
01380 05E4 B6 062A	LDA A	DONTST	CHECK IF DECPRS DONE
01390 05E7 81 00	CMP A	#0	DONE YET?
01400 05E9 26 F8	BNE	TAINTR	NO. KEEP LOOPING
01410 05EB 0F	SEI		INSURE INTR DISABLE
01420 05EC CE 0878	LDX	#GOAGAIN	PRINT GO AGAIN MSG
01430 05EF BD CA8F	JSR	OUTNCR	"ENTER CONTINUATION COMMAND"
01440 05F2 BD CA2C	JSR	KEYBD0	
01450 05F5 FE 3002	LDX	ENDBUF	NOW GET COMMAND
01460 05F8 09	DEX		
01470 05F9 E6 00	LDA B	0,X	
01480 05FB 17	TBA		PUT IT IN A
01490 05FC 90 28	SUB A	\$28	TRANSFORM X TO 0
01500 05FE 97 20	STA A	CHNLBF	PUT IN CHNL FLAG
01510 0600 F7 0846	STA B	CHNASC	PUT ASCII IN ERROR MSG
01520 0603 C1 58	CMP B	#'X	WAS IT .LT. 'X?
01530 0605 2D 07	BLT	TADEC5	YES. EXIT ROUTINE
01540 0607 C1 5A	CMP B	#'Z	WAS CMMD .GT. 'Z ?
01550 0609 2E 03	BGT	TADEC5	YES. EXIT ROUTINE
01560 060B 7E 05A0	JMP	TADEC4	NO. GO EXECUTE DECPRS AGAIN
01570 060E 7E 0100	TADEC5	JMP	RETURN TO DISPLAY
01580 *			
01590	*****		
01600 *	END DECPRS INIT ROUTINE		
01610	*****		
01620 *			
01630 *			
01640	* PARAMETER BUFFER		

01650	*		
01660 0611 0002	PUPPTR RMB	2	CURRENT MEM FILE WORD POINTER
01670 0612 0001	PIPPTR RMB	1	CH X DATA BUF
01680 0614 0001	XDTA RMB	1	CH Y DATA BUF
01690 0615 0001	YDTA RMB	1	CH Z DATA BUF
01700 0616 0001	ZDTA RMB	1	CH X 1ST DIFF BUFFER
01710 0617 0001	XSLAPE RMB	1	CH Y 1ST DIFF BUF
01720 0618 0001	YSLAPE RMB	1	CH Z 1ST DIFF
01730 0619 0001	ZSLAPE RMB	1	TIME COMPRESSION RUN LEN BUFF
01740 061A 0001	DELT RMB	1	DATA RET BUF FROM GETVLU
01750 061B 0001	VALUE RMB	1	COUNTER IN GETTIM
01760 061C 0001	CNT RMB	1	IRQ VECTOR SAVE BUF
01770 061D 0002	VECSAV RMB	2	BIT SET/RESET FLAG FROM GETBI
01780 061F 0001	BITVLU RMB	1	SIGN BIT ON DECODED DATA
01790 0620 0001	SIGN RMB	1	UNCODED CH X ACCEL DATA
01800 0621 0001	XACCEL RMB	1	TEMP BUF USED IN DECPRS
01810 0622 0001	XTEMP RMB	1	UNCODED CH Y ACCEL DATA
01820 0623 0001	YACCEL RMB	1	TEMP BUF USED IN DECPRS
01830 0624 0001	YTEMP RMB	1	UNCODED CH Z ACCEL DATA
01840 0625 0001	ZACCEL RMB	1	TEMP BUF USED IN DECPRS
01850 0626 0001	ZTEMP RMB	1	TEMP WORKING BUFFER
01860 0627 0002	WORKBF RMB	2	FLG USED IN DECPRS ALGORITHM
01870 0629 0001	TARFLG RMB	1	FLG USED TO EXIT INTR LOOP
01880 062A 0001	DCNTST RMB	1	D/A CONV OUTPUT BUFFER
01890 062B 0002	DADUF RMB	2	TEMP STACK SAVE BUFFER
01900 062D 0002	STKSAV RMB	2	*
01910	*		
01920	*FUNCTION : TARECN		
01930	*INPUTS : BUFFER VALUES		
01940	*OUTPUTS : D/A VALUES VIA ST-6800		
01950	*CALLS : GETVLU,GETTIM,OUTDA		
01960	*DESTROYS A,B,X,CC		
01970	*PURPOSE : THIS IS THE INTERRUPT SERVICE ROUTINE WHICH		
01980	* PERFORMS RECONSTRUCTION ON THE TOLAN-A ENCODED DATA		
01990	* FILE IN MEMORY. THE DATA IS DECODED AND OUTPUT ON		
02000	* D/A (ST6800) CH 0.		
02010	*		
02020	*****		
02030	*		
02040 062F B7 E400	TARECN STA A ADCZRO		ACK INTR. RESET INTR FF ON ST
02050 0632 01	NOP		
02060 0633 BF 062D	STS STKSAV		SAVE ENTRY STACK FOR RTI
02070 0636 F6 0846	LDA A CHNASC		CHECK CHNL SELECT, OUTPUT CUR
02080 0639 81 58	CMP A #'X		IS CHNL SELECT X ?
02090 063B 26 08	BNE TAREC1		NO. CHK IF Y
02100 063D B6 0614	LDA A XDTA		YES. GET CUR XDTA VALUE.
02110 0640 FD 0796	JSR OUTDA		OUTPUT VALUE VIA D/A CH 0
02120 0643 20 12	BRA TARE22		
02130 0645 81 59	TAREC1 CMP A #'Y		IS CHNL SELECT Y ?
02140 0647 26 08	BNE TAREC2		NO. MUST BE Z
02150 0649 F6 C615	LDA A YDTA		GET CUR YDTA VALUE.
02160 064C BD 0796	JSR OUTDA		OUTPUT Y VALUE VIA D/A CH 0
02170 064F 20 06	BRA TARE22		
02180 0651 F6 0616	TAREC2 LDA A ZDTA		GET CUR ZDTA VALUE.

02190 0654 BD 0756	JSR	OUTDA	OUTPUT Z VALUE VIA D/A CH 0
02200 0657 B6 0629	LDA A	TARFLG	WAS LAST DATA PT EXTROPOLOATED
02210 065A 81 00	CMP A	#0	
02210 065C 26 33	BRA	MAIN3	YES. NOW GO GET NEXT ACCEL. N
02230 065E 86 58	LDA A	#'X	PUT ASCII X IN GETVLU ERROR M
02240 0660 B7 0A3D	STA A	OVFLCH	
02250 0663 BD 0700	JSR	GETVI.U	NOW GO GET NEXT XACCEL VALUE
02260 0666 B6 061B	LDA A	VALUE	AND PUT XACCEL
02270 0669 B7 0621	STA A	XACCEL	
02280 066C B7 0622	STA A	XTEMP	SAVE XACCEL IN TEMP BUFF
02290 066F 86 59	LDA A	#'Y	PUT 'Y IN GETVLU ERROR MSG
02300 0671 B7 0A3D	STA A	OVFLCH	
02310 0674 BD 0700	JSR	GETVI.U	GET YACCEL DATA
02320 0677 B6 061B	LDA A	VALUE	
02330 067A B7 0623	STA A	YACCEL	PUT IN YACCEL VAR
02340 067D B7 0624	STA A	YTEMP	PUT IN YACCEL TEMP VAR
02350 0680 86 5A	LDA A	#'Z	PUT 'Z IN GETVLU ERROR MSG
02360 0682 B7 0A3D	STA A	OVFLCH	
02370 0685 BD 0700	JSR	GETVLU	GET ZACCEL DATA
02380 0688 B7 0625	STA A	ZACCEL	PUT IN ZACCEL VAR
02390 068B B7 0626	STA A	ZTEMP	PUT IN ZACCEL TEMP VAR
02400 068E BD 075E	JSR	GETTIM	NOW GET TIME COPRESSION RUN L
02410 0691 7F 0629	TAREC3	CLR	RESET FLAGS
02420 0694 7A 061A	DEC	DELT	NOW CK IF RUN LEN ONLY 1
02430 0697 27 09	BEQ	TAREC4	YES. ACCEL VAR NOT ZERO
02440 0699 7F 0621	CLR	XACCEL	ACCEL WAS ZRO SO CLR X-Y-ZACC
02450 069C 7F 0623	CLR	YACCEL	
02460 069F 7F 0625	CLR	ZACCEL	
02470 06A2 B6 0617	TAREC4	LDA A	NOW GET LAST SLOPE
02480 06A5 BB 0621	ADD A	XACCEL	& CALC NEW SLOPE WITH ACCEL
02490 06A8 B7 0617	STA A	XSLOPE	
02500 06AB B6 0618	LDA A	YSLOPE	
02510 06AE BB 0623	ADD A	YACCEL	
02520 06E1 B7 C618	STA A	YSLOPE	
02530 06B4 B6 0619	LDA A	ZSLOPE	
02540 06B7 BB 0625	ADD A	ZACCEL	
02550 06BA B7 0619	STA A	ZSLOPE	
02560 06BD B6 0614	LDA A	XDTA	NOW CALC DATA VALUE USING SLOPES
02570 06C0 BB 0617	ADD A	XSLOPE	
02580 06C3 B7 0614	STA A	XDTA	
02590 06C6 B6 0615	LDA A	YDTA	
02600 06C9 BB 0618	ADD A	YSLOPE	
02610 06CC B7 0615	STA A	YDTA	
02620 06CF B6 0616	LDA A	ZDTA	
02650 06D8 B6 061A	LDA A	DELT	
02660 06DB 27 22	BEQ	TARRTI	WAS RUN LEN .GT. 2 PIS ?
02670 06DD 81 01	CMP A	#1	IF YES, DONT GET NEW TIM FIL
02680 06DF 27 07	BEQ	TAREC5	JUST EXTROPOLOATE POINTN UNTIL
02690 06E1 86 37	LDA A	#55	DELTA IS REDUCED TO 1
02700 06E3 B7 0629	STA A	TARFLG	SET TARFL1
02710 06E6 20 17	BRA	TARRTI	
02720 06E8 B6 0622	TAREC5	LDA A	RETRIEVE ACCEL VLUS PREVIOUSL
02730 06FB B7 0621	STA A	XACCEL	
02740 06EE B6 0624	LDA A	YTEMP	

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02750 06F1 B7 0623 STA A YACCEL
02760 06F4 B6 0626 LDA A ZTE1P
02770 06F7 B7 0625 STA A ZACCL
02780 06FA 86 55 LDA A #55 NOW SET TARFL2 FLAG
02790 06FC B7 0629 STA A TARFLG
02800 06FF 3B TARRTI RTI
02810 *
02820 **** END OF INTERRUPT HANDLER ****
02830 *
02840 *
02850 *
02860 * SUBROUTINES
02870 *
02880 *
02890 *FUNCTION : GETVLU
02900 *INPUTS : BUFPTR,BITPTR,SIGN BUFFERS
02910 *OUTPUTS : VALUE BUFFER
02920 *CALLS : GETEIT
02930 *DESTROYS : A,X,CC
02940 *PURPOSE : THIS RUTINE SCANS MEMORY STARTING AT POI
02950 * DEFINED BY BUFPTR,BITPTR AND GETS ACCEL VALUE. RUF
02960 * AND BITPTR ARE UPDATED.
02970 *
02980 *
02990 0700 7F 061B GETVLU CLR VALUE INSURE VALUE IS CLEAR ON STAR
03000 0703 ED 07ED JSR GETBIT GET DELIMITER BIT
03010 0706 E6 061F LDA A BITVLU
03020 0709 81 00 CMP A #0 IS DELIM BIT =0 ?
03030 070B 27 06 BEQ GETVLL YES. GET NEXT BIT
03040 070D CE 097D LDX #SYNCMS NO. SYNC ERROR HAS OCCURED. E
03050 0710 7E 07F5 JMP ERROR
03060 0713 7F 0620 GETVLL CLR SIGN INSURE SIGN CLR
03070 0716 ED 07ED JSR GETBIT GET SIGN BIT
03080 0719 B6 061F LDA A BITVLU
03090 071C 81 00 CMP A #0 IS ACCEL POS?
03100 071E 27 03 BEQ GETVL2 YES. LEAVE SIGN CLR
03110 0720 73 0620 COM SIGN NO. SET SIGN FLAG
03120 0723 BD 07ED GETVL2 JSR GETRIT GET FIRST CODE BIT
03130 0726 B6 061F LDA A BITVLU
03140 0729 26 11 RNE GETVLL4
03150 072B FE 0611 GETVLL3 LDX BUFPTR END OF VALUE DECODE. PACK UP
03160 072E 78 0613 ASL BITPTR FOR CORRECT POSITIONING FOR N
03170 0731 24 2A BCC GETVLL6 VALUE FETCH
03180 0733 79 0613 ROL BITPTR
03190 0736 09 DEX
03200 0737 FF 0611 STX BUFPTR
03210 073A 20 21 BRA GETVLL6
03220 073C 7C 061B GETVLL4 INC VALUE NOW PAST DELIM & SIGN. NOW CO
03230 073F 2A 06 BPL GETVLL5 BITS. IS CODE WORD > 127 ?
03240 0741 CE 09FB LDX #OVFLMS YES. OVERFLOW HAS OCCURED. PR
03250 0744 7E 07F5 JMP ERROR
03260 0747 BD 07ED GETVLL5 JSR GETBIT NOW GET DATA VALUE COMM BITS
03270 074A B6 061F LDA A BITVLU
03280 074D 81 00 CMP A #0 DELIM DETECTED YET?

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03290 074F 26 EB      BNE    GETVL4  NO. KEEP FETCHING BITS
03300 0751 F6 0620    LDA A  SIGN   YES. GET SIGN AND ADJUST CNT
03310 0752 40 00       CMP A  #0
03320 0756 27 D0       LDQ    GETVL3  SIGN C. VALUE CO. CC ALJ PC
03330 0756 70 061B     NEG    VALUE   SIGN 1. TAKE 2'S COMP, ADJ PT
03340 075B 20 CE       BRA    GETVL3
03350 075D 39         GETVL6 RTS
03360 *
03370 ****
03380 *          END OF SUBROUTINE GETVLU
03390 ****
03400 *
03410 *
03420 *FUNCTION : GETTIM
03430 *INPUTS : BITVLU (CURR BIT VALUE PT'ED TO)
03440 *OUTPUTS : DELT (CURRENT TIME COUNT)
03450 *CALLS : GETBIT, ERRCR
03460 *DESTROYS : A,X,CC REGISTERS
03470 *PURPOSE : THIS ROUTINE, WHEN CALLED AFTER 3 GETVLU
03480 * RETURNS THE VALUE OF THE RUN LENGTH TIME COUNTER I
03490 * VARIABLE DELT.
03500 *
03510 *
03520 075E BD 07ED GETTIM JSR    GETEIT  GET DELIM BIT
03530 0761 B6 061F       LDA A  BITVLU
03540 0764 81 00       CMP A  #0      IS BIT 0?
03550 0766 27 06       BEQ    GETTI1  YES. PROPER DELIM. CONTINUE
03560 0768 CE 09AA     LDX    #SYNCTI  NO. SYNC ERROR. PRINT & ERR O
03570 076B 7E 07F5     JMP    ERROR
03580 076E 7F 031A GETTI1 CLR    DELT   INSURE TIME CNT INITIALLY CLR
03590 0771 86 07       LDA A  #7      INIT CNT FOR 7 BIT TIME
03600 0773 B7 061C     STA A  CNT
03610 0776 BD 07ED GETTI2 JSR    GETBIT  GET DELTA T BIT
03620 0779 B6 061F     LDA A  BITVLU
03630 077C BA 061A     ORA A  DELT   SET LSB BIT ACCORD TO T BIT
03640 077F B7 061A     STA A  DELT
03650 0782 78 061A     ASL    DELT   SHIFT DELT FOR NEXT BIT
03660 0785 7A 061C     DEC    CNT   DONE WITH TIME YET?
03670 0788 26 EC       BNE    GETTI2  NO. KEEP FETCHING TIME BITS
03680 078A B6 061A     LDA A  DELT   YES. GET TIME VALUE
03690 078D 26 06       BNE    GETTI3  IS TIME CNT 0 ?
03700 078F CE 0A3F     LDX    #TINERR YES. ERROR. PRNT INR MSG
03710 0792 7E 07F5     JMP    ERROR  AND ERROR OFF
03720 0795 39         GETTI3 RTS
03730 *
03740 ****
03750 *          END GETTIM SUBROUTINE
03760 ****
03770 *
03780 *
03790 *FUNCTION : OUTDA
03800 *INPUTS : DATA VALUE TO BE D/A'ED IN ACC A
03810 *OUTPUTS : VALUE IN ACC A VIA D/A CH 0
03820 *CALLS : NOTHING

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03830      *DESTROYS : B,CC
03840      *PURPOSE : THIS ROUTINE OUTPUTS THE DATA IN
03850      * ACCUMULATOR A TO THE D/A FOR CONVERSION AND
03860      * TRANSMISSION. THE 8 BIT VALUE IS SHIFTED
03870      * TO FORM THE 12 BIT OPERAND REQUIRED BY THE
03880      * D/A (ST6800).
03890      *
03900      *
03910 0796 36    OUTDA PSH A      SAVE VLU IN A ACC
03920 0797 B7 062C STA A DABUF+1  PUT VALUE IN D/A OUT BUFF
03930 079A 7F 062B CLR DABUF      CLR MSB BYTE OF D/A OUT BUFF
03940 079D 78 062C ASL DABUF+1  LEFT SHIFT 2 BYTE D/A WORD
03950 07A0 79 062B ROL DABUF      FOR CONVERSION OF 8 BIT
03960 07A3 78 062C ASL DABUF+1  VALUE TO 12 BIT
03970 07A6 79 062B ROL DABUF
03980 07A9 78 062C ASL DABUF+1
03990 07AC 79 062B ROL DABUF
04000 07AF 78 062C ASL DABUF+1
04010 07B2 79 062B ROL DABUF
04020 07B5 FE 062B LDX DABUF      NOW LD INDEX WITH 12 BIT VLU
04030 07B8 FF E500 STX DACZRO    AND OUTPUT TO D/A
04040 07BB 32    PUL A          RETRIEV ENTY A ACC VLUE
04050 07BC 39    RTS
04060      *
04070      ****
04080      *
04090      *
04100      *
04110      *
04120      *FUNCTION : GETBIT
04130      *INPUTS : BUFPTR,BITPTR POINTERS
04140      *OUTPUTS : BITVLU, DONIST FLAGS
04150      *CALLS :NOTHING
04160      *DESTROYS :A,CC
04170      *PURPOSE : THIS ROUTINE CHECKS THE STATE OF THE BIT
04180      * POINTED TO BY THE BUFPTR,BITPTR PAIR AND SETS THE
04190      * THE BITVLU FLAG ACCORDINGLY. THIS ROUTINE CHECK
04200      * FOR END OF MEMORY AND FORCES EXIT OF RECONSTRUCTIO
04210      * ROUTINES IF EOF IS DETECTED.
04220      *
04230      *
04240 07BD 7F 061F GETBIT CLR    BITVLU      INSURE BITVLU STARTS RESET
04250 07C0 FE 0611 LDX    BUFPTR     GET MEM WORD POINTER
04260 07C3 F6 0613 LDA B  BITPTR     GET BIT PTR IN THAT MEM W
04270 07C6 E5 00  BIT B  0,X        CHECK IF BIT SLT
04280 07C8 27 03  BEQ    GETBIT1    NO IT WAS RESET. INSUT BITVLU
04290 07CA 7C 061F INC    BITVLU     YES. BIT WAS SET. SUT BITVLU
04300 07CD 74 0613 GETBIT1 LSR    BITPTR     NOW UPDATE POINTERS
04310 07D0 24 22  BCC    GETRTS
04320 07D2 76 0613 ROR    BITPTR
04330 07D5 08    INX
04340 07D6 FF 0611 STX    BUFPTR
04350 07D9 F6 0611 LDA B  BUFPTR     CHECK FOR END OF MEMORY
04360 07DC C1 00  CMP B  #$80

```

```

04370 07DE 26 14      BNE    GETRTS  NO. CONTINUE
04380 07E0 0F          SEI     NOW TURN OFF INTR AND RET
04390 07E1 CE 4000      LDX    #$4000
04400 07E4 FF E500      STX    DACZRO
04410 07E7 FE 061D      LDX    VECSAV
04420 07EA FF FFF8      STX    IRQVEC
04430 07ED 7F 062A      CLR    DONTST YFS. CLR DONTST FLAG AND RTI
04440 07F0 BE 062D      LDS    STKSAV RETRIEVE INTR STACK PTR
04450 07F3 3B          RTI
04460 07F4 39          GETRTS RTS
04470 *
04480 ****
04490 *           END OF SUBROUTINE GETBIT
04500 ****
04510 *
04520 *
04530 *FUNCTION : ERROR
04540 *INPUTS : MESSAGE STRING POINTED TO BY INDEX RFG
04550 *OUTPUTS : ERROR MESSAGE TO CONSOLE DEVICE
04560 *CALLS : OUTNCR,KEYBD0,DISPLA
04570 *DESTROYS : A,B,X,CC
04580 *PURPOSE : THIS ROUTINE PRINTS ERROR MESSAGES UPON
04590 * ERROR DETECTION BY THE DECOMPRESSION ROUTINES.
04600 *
04610 *
04620 07F5 BD CA8F ERROR JSR    OUTNCR
04630 07F8 CE 09DF      LDX    #ANYKEY GET ANYKEY MESSAGE
04640 07FB BD CA8F      JSR    OUTNCR PRINT "TO CONTINUE PRESS ANY"
04650 07FE BD CA2C      JSR    KEYBD0
04660 0801 0F          SEI    NOW CLR INTR CONDITION
04670 0802 CE 4000      LDX    #$4000 AND RETURN TO DISPLAY
04680 0805 FF E500      STX    DACZRO
04690 0808 FE 061D      LDX    VECSAV RETRIEV IRQ VEC ADDR
04700 080B FF FFF8      STX    IRQVEC
04710 080E 7E 0100      JMP    DISPLA RETURN TO "DISPLAY"
04720 *
04730 ****
04740 *           END ERROR HANDLING ROUTINE
04750 ****
04760 *
04770 *           MESSAGE STRINGS
04780 *
04790 0811 1A0D GMSG   FDB    $1A0D,$0A07
04800 0815 44 FCC    /DECOMPRESSION OVERLAY/
04810 082A 0D0A FDB    $0D0A,$0D0A
04820 082E 43 FCC    /CURRENT CHANNEL FLAG IS /
04830 0846 0001 CHNASC RMB    1
04840 0847 0D0A FDB    $0D0A,$0D0A
04850 084B 44 FCC    /DO YOU WISH TO EXECUTE THIS MODULE (
04860 0877 04 FCB    4
04870 0878 0D0A COAGIN FDB    $0D0A,$0A07
04880 087C 44 FCC    /DECOMPRESSION AND DISPLAY COMPLETE/
04890 089E 0D0A FDB    $0D0A,$0D0A
04900 08A2 45 FCC    /ENTER CONTINUATION COMMAND/

```

04910 08LC 0D0A	FDB	\$0D0A
04920 08BE 20	FCC	/ X=DISPLAY CHANNEL X ON D-A CHANNE
04930 08C7 0D0A	FDB	\$0D0A
04940 08E7 20	FCC	/ Y=DISPLAY CHANNEL Y OF D-A CHANNE
04950 090E 0D0A	FDB	\$0D0A
04960 0910 20	FCC	/ Z=DISPLAY CHANNEL Z ON D-A CHANNE
04970 0937 0D0A	FDB	\$0D0A,\$0D0A
04980 093B 41	FCC	/ANY OTHER KEY RETURNS CONTROL TO "DI
04990 0966 0D0A	FDB	\$0D0A,\$0D0A
05000 096A 45	FCC	/ENTER COMMAND NOW=/
05010 097C 04	FCB	4
05020 097D 1A07	SYNCMS	FDB \$1A07,\$0D0A
05030 0981 53	FCC	/SYNC ERROR DETECTED BY TOLAN-A DECOD
05040 09A7 0D0A	FDB	\$0D0A
05050 09A9 04	FCB	4
05060 09AA 1A07	SYNCTI	FDB \$1A07,\$0D0A
05070 09AE 53	FCC	/SYNC ERROR DETECTED BY /
05080 09C5 54	FCC	/TOLAN-A DURING TIME FETCH/
05090 09DC 04	FCB	4
05100 09DF 0D0A	ANYKEY	FDB \$0D0A
05110 09E1 54	FCC	/TO CONTINUE PRESS ANY KEY/
05120 09FA 04	FCB	4
05130 09FB 1A07	OVFLMS	FDB \$1A07,\$0D0A
05140 09FF 41	FCC	/ACCEL VALUE OVERFLOW DETECTED BY TOL
05150 0A2F 0D0A	FDB	\$0D0A
05160 0A31 20	FCC	/ IN CHANNEL /
05170 0A3D 0001	OVFLCH	RMB 1
05180 0A3E 04	FCB	4
05190 0A3F 1A07	TIMERR	FDB \$1A07,\$0D0A
05200 0A43 54	FCC	/TIME COUNT ERROR DETECTED BY TOLAN-A
05210 0A6F 0D0A	FDB	\$0D0A
05220 0A71 54	FCC	/TIME CNT (DELT)=0/
05230 0A82 04	FCB	4
05240	*	*****
05250	*	*****
05260	*	*****
05270	*	END OF DECPRS OVERLAY ROUTINES
05280	*	*****
05290	*	*****
05300	*	*****
05310		END

## ENTROPY

```
0100 REM ****
0110 REM *
0120 REM *      EKG ENTROPY CALCULATION PROGRAM
0130 REM *
0140 REM ****
0150 REM *
0160 REM *      THIS PROGRAM READS THE FREQUENCY
0170 REM * OF OCCURENCE DATA IN MEMORY FROM AN
0180 REM * EKG DATA COLLECTION AND CALCULATES
0190 REM * THE ENTROPY OF THE X,Y,AND Z DATA
0200 REM * SOURCES.  THE CALCULATED ENTROPY IS
0210 REM * STORED BACK IN THE MEMORY FILE HEADER
0220 REM * IN ASCII.  THIS HEADER CAN THEN BE
0230 REM * INSERTED INTO THE DISK FILE USING
0240 REM * MINIDOS.
0250 REM *
0260 REM ****
0270 REM
0280 DIM D(255),N$(8),T(255),M9(6),T$(3)
0290 STRING= 8
0291 DIGITS= 4
0292 T$(1)="X"
0293 T$(2)="Y"
0294 T$(3)="Z"
0295 LINE= 80
0296 L7=6
0297 E1=13507
0313 REM
0325 REM ****
0326 REM * NOW GET FILENAME FROM MEMORY & PRINT IT
0327 REM ****
0328 REM
0330 FOR I=1 TO 8
0340 J=13313+I
0350 N$(I)=CHR$(PEEK(J))
0360 A$=A$+N$(I)
0370 NEXT I
0372 P$=CHR$(12)
0373 PRINT P$
0374 PRINT
0375 PRINT
0376 PRINT
0377 PRINT
0378 PRINT
0380 PRINT
0385 PRINT "EKG ENTROPY CALCULATION"
0386 PRINT
0387 PRINT
0390 PRINT "FILENAME.....";A$
0400 PRINT
0401 PRINT
0402 GOSUB 1696
0403 GOSUB 1900
0404 IF G$ = "S" THEN GO TO 1120
```

## ENTROPY

```
0410 REM
0411 REM ****
0412 REM * B1=MEM FILE DESTINATION FOR ASCII RESULTS
0413 REM * 13856=$3620 (HEX)=FILE MEM USED BY BASIC. DATA LITPLACTD
0414 REM * 13567=34FF (HEX)=XPDF BUFFER-1
0415 REM * 14079=36FF (HEX)=YPDF BUFFER-1
0416 REM * 14591=38FF (HEX)=ZPDF BUFFER-1
0417 REM * 13500=34BC (HEX)=BASSAV. MEM BUF FOR $3620 DATA SAVE
0419 REM ****
0420 REM
0421 REM
0424 REM
0425 REM ****
0426 REM * NOW GET VLU IN BASSAV & STR TO $3620 (HEX)
0427 REM ****
0428 REM
0429 B=13856
0430 C1=1
0432 N=PEEK(13500)
0433 POKE( B,N)
0435 FOR I=1 TO 255
0436 T(I)=0
0437 NEXT I
0438 T9=0
0439 S2=0
0440 S=0
0441 REM
0442 REM ****
0443 REM * NOW BEGIN LOOP TO CALC X,Y,Z & TOTAL ENTROPY
0444 REM ****
0445 REM
0450 IF C1 <> 1 THEN GO TO 480
0460 K=13567
0480 IF C1 <> 2 THEN GO TO 510
0490 K=14079
0510 IF C1 <> 3 THEN GO TO 540
0520 K=14591
0540 IF C1>= 4 THEN GO TO 800
0542 REM
0543 REM ****
0544 REM * GET 2 BYTE PDF DATA & MERGE INTO FLTING POINT NUMBER
0545 REM ****
0546 REM
0550 FOR I=1 TO 509 STEP 2
0560 M=(I+1)/2
0570 D(M)=256*PEEK(I+K)+PEEK(I+K+1)
0572 T(M)=T(M)+D(M)
0590 S=D(M)+S
0592 S2=D(M)+S2
0600 NEXT I
0605 D9=256*PEEK(K+511)+PEEK(K+512)
0606 S=S+D9
0607 T9=T9+D9
0608 S2=S2+D9
```

## ENTROPY

```
0610 E=0
0612 REM ****
0613 REM * NOW CALCULATE PROBABILITY OF OCCURENCE OF SPECIFIC
0614 REM * VLU'S FROM NUM OF OCCURENCES. THEN CALC ENTROPY SUM.
0615 REM ****
0616 REM ****
0617 REM
0620 FOR I=1 TO 255
0630 D(I)=D(I)/S
0640 IF D(I)=0 THEN GO TO 660
0650 E=E+(D(I)*(-(LOG(D(I))/.693147)))
0660 NEXT I
0663 D9=D9/S
0664 IF D9=0 THEN GO TO 666
0665 E=E+(D9*(-(LOG(D9)/.693147)))
0666 IF E=0 THEN GOTO 672
0667 C2=8/E
0668 GOTO 683
0672 C2=1
0673 GOTO 680
0674 REM ****
0675 REM * PRINT RESULTS OF CALC TO TERMINAL (PRINTER).
0676 REM * ****
0677 REM ****
0678 REM
0680 PRINT
0681 PRINT "MAXIMUM COMPRESSION FOR CHANNEL ";T$(C1); " IS INFINITE"
0682 GOTO 690
0683 PRINT
0684 PRINT "MAX COMPRESSION RATIO FOR CHANNEL ";T$(C1); " IS ";C2;" ; 1"
0690 PRINT
0700 PRINT "EKG LEAD ";T$(C1); " ENTROPY = ";E;" BITS."
0710 PRINT
0712 REM ****
0713 REM * NOW CONVERT CALCULATED ENTROPY TO ASCII & STORE BACK
0714 REM * IN MEMORY BUFFER
0715 REM *
0716 REM ****
0717 REM
0720 LET E$=STR$(E)
0740 GOSUB 1750
0785 C1=C1+1
0790 GOTO 440
0800 REM ****
0801 REM * WITH X,Y,Z ENTROPY CALCULATED, NOW CALCULATE TOTAL,
0802 REM * COMBINED ENTRCPY BY ADDING X,Y,Z B'N COUNTS & DIVIDING
0803 REM * BY TOTAL SAMPLE COUNT
0804 REM *
0805 REM ****
0806 REM
0807 E5=0
0810 FOR I=1 TO 255
0820 T(I)=T(I)/S2
0822 IF T(I)=0 THEN GO TO 840
```

## ENTROPY

```
0830 E5=E5+(T(I)*(-(LOG(T(I))/.693147)))
0840 NEXT I
0850 T9=C3/S2
0852 IF T9=0 THEN GO TO 870
0860 E5=E5+(T9*(-(LOG(T9)/.693147)))
0870 C3=8/E5
0872 REM ****
0873 REM ****
0874 REM * PRINT COMBINED ENTROPY TO TERM & STR RESULT IN MEM BUFFER
0875 REM ****
0876 REM
0880 PRINT
0890 PRINT "MAX COMPRESSION RATIO FOR 3 LEAD EKG SYSTEM ";C3;" : 1"
0910 PRINT
0920 PRINT "3 LEAD EKG SOURCE ENTROPY = ";E5;" BITS."
0930 PRINT
0940 LET E$=STR$(E5)
0950 GOSUB 1750
1000 LET E$=STR$(C3)
1005 GOSUB 1750
1010 REM ****
1020 REM ****
1030 REM * NOW GET OTHER STATISTICAL VARIABLES &
1040 REM * CALCULATE CHANNEL MAX,MINS,COMPRESSION RATIO
1050 REM * OBTAINED, AND COMPRESSION TIME EFFICIENCY
1060 REM ****
1070 REM
1080 REM ****
1082 REM * START WITH ACHIEVED COMPRESSION RATIO. THEN
1090 REM * PRINT THE TIME COMPRESSION EFFICIENCY. THIS IS
1091 REM * THE PERCENTAGE OF THE TIME AVAILABLE THAT WAS USED
1092 REM * TO COLLECT,COMPRESS, & CALCULATE STAT VARIABLES.
1100 REM ****
1110 REM
1120 DIGITS= 4
1130 L7=6
1140 M4=PEEK(13483)+PEEK(13482)*256+PEEK(13481)*256*256
1150 D4=PEEK(13487)+PEEK(13486)*256+PEEK(13485)*256*256
1160 D4=PEEK(13484)*256*256*256+D4
1170 C4=D4/M4
1180 LET E$=STR$(C4)
1190 GOSUB 1750
1200 PRINT
1201 PRINT "COMPRESSION RATIO ACHIEVED = ";C4;" : 1"
1210 IF G$="S" THEN GOTO 1260
1211 REM
1212 REM ****
1213 REM * CALCULATE OVERALL COMPRESSION EFFICIENCY
1214 REM ****
1215 REM
1216 DIGITS= 1
1217 C6=((C4-1)/(C3-1))*100
1218 PRINT
1219 PRINT "ACHIEVED COMPRESSION EFFICIENCY = ";C6;" %"
```

## ENTROPY

```
1220 LET E$=STR$(C6)
1221 L7=4
1222 GOSUB 1750
1229 REM
1230 REM ****
1240 REM * NOW CALCULATE COMPRESSION TIME EFFECIENCY
1250 REM ****
1252 REM
1260 S4=PEEK(13461)+PEEK(13460)*256
1262 L4=PEEK(13462)
1264 L5=PEEK(13459)+PEEK(13458)*256+PEEK(13457)*256*256
1266 L5=PEEK(13456)*256*256*256+L5
1270 T4=(1-(L5/(L4*S4)))*100
1272 DIGITS= 1
1274 L7=4
1276 LET E$=STR$(T4)
1278 GOSUB 1750
1280 PRINT
1282 PRINT "COLLECTION TIME EFFICIENCY = ";T4;" %"
1284 PRINT
1290 REM
1291 REM ****
1292 REM * CALCULATE COLLECTION DURATION
1293 REM ****
1294 REM
1295 T6=S4/S8
1296 PRINT
1297 PRINT "COLLECTION DURATION = ";T6;" SECONDS"
1298 PRINT "AT A SAMPLE RATE OF ";S8;" HZ"
1299 LET E$=STR$(T6)
1300 GOSUB 1750
1301 REM
1302 REM ****
1310 REM * RETRIEVE DATA MAX & MINS & CALC VOLTS
1320 REM ****
1330 REM
1339 DIGITS= 5
1340 FOR I=6 TO 1 STEP -1
1350 M9(I)=PEEK(13460+I*3)
1360 IF M9(I) > 127 THEN GOTO 1390
1370 M9(I)=M9(I)*.0390625
1380 GOTO 1400
1390 M9(I)=-(10-M9(I)*.0390625)
1400 NEXT I
1401 E1=13408
1402 FOR I4=5 TO 1 STEP -2
1403 LET E$=STR$(M9(I4))
1404 L7=7
1405 GOSUB 1750
1406 I3=I4+1
1407 LET E$=STR$(M9(I3))
1408 GOSUB 1750
1409 NEXT I4
1410 REM
```

## ENTROPY

```
1420 REM ****
1430 REM * NOW PRINT MAX & MINS TO TERM (PRINTER)
1440 REM ****
1450 REM
1460 PRINT
1470 FOR I=1 TO 6 STEP 2
1480 PRINT
1490 PRINT "CHANNEL ";T$(I+1)/2;" MINIMUM = ";M9(7-I);" VOLTS."
1500 PRINT
1510 PRINT "CHANNEL ";T$(I+1)/2;" MAXIMUM = ";M9(6-I);" VOLTS."
1520 PRINT
1530 NEXT I
1540 REM
1550 REM ****
1560 REM * FINISH WITH A PRINT OF NUM OF BITS USED TO
1570 REM * STORE X,Y,Z IN MEM WITH THIS COMPRESSION TYPE
1580 REM ****
1590 REM
1600 X3=PEEK(13490)+PEEK(13489)*256+PEEK(13488)*256*256
1610 Y3=PEEK(13493)+PEEK(13492)*256+PEEK(13491)*256*256
1620 Z3=PEEK(13496)+PEEK(13495)*256+PEEK(13494)*256*256
1630 T3=PEEK(13499)+PEEK(13498)*256+PEEK(13497)*256*256
1632 DIGITS= 0
1640 PRINT
1650 PRINT "NUMBER OF BITS USED TO STORE CHANNEL X =";X3
1652 PRINT
1660 PRINT "NUMBER OF BITS USED TO STORE CHANNEL Y =";Y3
1662 PRINT
1670 PRINT "NUMBER OF BITS USED TO STORE CHANNEL Z =";Z3
1672 PRINT
1680 PRINT "NUMBER OF BITS USED TO STORE TIME =";T3
1682 PRINT
1683 PRINT
1690 GOTO 2000
1691 REM
1692 REM ****
1693 REM * SUB TO PRINT OUT COMPRESSION TYPE
1694 REM ****
1695 REM
1696 H=PEEK(13312)*256+PEEK(13313)
1697 IF H=20035 THEN HS(1)="NOT COMP"
1698 IF H=20035 THEN HS(2)="RESSED "
1699 IF H=21569 THEN HS(1)="TOLAN A "
1700 IF H=21569 THEN HS(2)=" "
1701 IF H=21570 THEN HS(1)="TOLAN B "
1702 IF H=21570 THEN HS(2)=" "
1703 IF H=17487 THEN HS(1)="DOWER "
1704 IF H=17487 THEN HS(2)=" "
1705 IF H=21584 THEN HS(1)="TURNING "
1706 IF H=21584 THEN HS(2)="POINT "
1707 IF H=18766 THEN HS(1)="2ND ORD "
1708 IF H=18766 THEN HS(2)="INTERPOL"
1709 PRINT "COMPRESSION USED.....";HS(1);HS(2)
1710 PRINT
```

ENTROPY

```
1711 RETURN
1720 REM
1730 : ****
1732 REM * SUBROUTINE TO WRITE ASCII DATA IN ES TO MEM
1733 REM * LOCATION POINTED TO BY E1. WILL WRITE 8 CHAR
1734 REM * STRINGS TO MEMORY.
1739 REM ****
1740 REM
1750 FOR I=1 TO L7
1760 F$=MID$(E$,I,L7)
1770 F=ASC(F$)
1780 POKE( E1,F)
1781 E1=E1+1
1790 NEXT I
1800 POKE( E1,4)
1801 E1=E1+1
1810 RETURN
1820 REM
1830 REM ****
1840 REM * SUBROUTINE TO PROMPT COMMAND FOR SHORT RUN OR LONG
1850 REM * RUN. SHORT DOES NOT CALCULATE ENTROPY OR COMPRESSION
1860 REM * RATIOS.
1870 REM ****
1880 REM
1900 PRINT
1910 PRINT "ENTER S FOR SHORT RUN (NO ENTROPY CALCULATED)"
1920 PRINT "ENTER L FOR LONG RUN (WITH ENTROPY CALCULATED)"
1930 INPUT G$
1932 PRINT "ENTER COLLECTION SAMPLE RATE (IE. 500)"
1934 INPUT O$
1935 S8=VAL(O$)
1940 RETURN
1950 REM
1960 REM ****
1970 REM * END OF ENTROPY CALCULATION PROGRAM
1980 REM ****
1990 REM
2000 PRINT
2010 PRINT
2020 PRINT "ENTROPY CALCULATION & STATISTICS PRINTOUT COMPLETE"
2030 END
```

#### Appendix D

This appendix contains a listing of the data compressed by the TOLAN-A8 module for the thesis experiment.

EKG SAMPLE COLLECTION STATISTICS : PAGE 1  
FILENAME . . . . . TAI359PA  
SUBJECT . . . . . PARRELL  
SAMPLING RATE . . . . . 500 HZ  
DATE OF COLLECTION . . . . . 23 OCT 80  
TIME OF COLLECTION . . . . . 1359  
COMPRESSION USED . . . . . TOLAN-A  
CHANNEL X ENTROPY . . . . . 2.4871 BITS  
CHANNEL Y ENTROPY . . . . . 2.4751 BITS  
CHANNEL Z ENTROPY . . . . . 2.4217 BITS  
TOTAL SOURCE ENTROPY . . . . . 2.4796 BITS  
PRESS RETURN FOR PAGE 2 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS : PAGE 2  
APPROX MAX COMPRESSION  
RATIO POSSIBLE . . . . . 3.2263 : 1  
COMPRESSION RATIO  
ACHIEVED . . . . . 2.2595 : 1  
ACHIEVED COMPRESSION  
EFFICIENCY . . . . . 56.5 % OF MAXIMUM  
COMPRESSION TIME  
EFFICIENCY OBTAINED . . . . . 77.8 % SMP INTERVAL  
COLLECTION DURATION . . . . . 26.2 SECONDS  
CHANNEL X MAXIMUM . . . . . 2.50000 VOLTS  
CHANNEL X MINIMUM . . . . . -0.1171 VOLTS  
CHANNEL Y MAXIMUM . . . . . 2.50000 VOLTS  
CHANNEL Y MINIMUM . . . . . -0.1171 VOLTS  
CHANNEL Z MAXIMUM . . . . . 2.50000 VOLTS  
CHANNEL Z MINIMUM . . . . . -0.1562 VOLTS  
COMMENTS . . . . . X,Y,Z IN COMMON LEAD 1  
PRESS RETURN FOR PAGE 3 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS: PAGE 3  
NUMBER OF SAMPLES TAKEN (SAMPNO)= 3336 (HEX)  
MAXIMUM LOOP COUNT PER INTERRUPT (LPCAL)= 27 (HEX)  
TOTAL WAITING LOOP COUNTS DURING COLLECTION (LOOPCT)= 0001B9DB (HEX)  
TIME EFFICIENCY = (1-(LOOPCT\*(SAMPNO\*LPCAL)))\*100  
CHANNEL MAXIMUMS AND MINIMUMS  
XMAX= 40 (HEX) AT SAMPLE NUMBER 0DA7 (HEX)  
XMIN= FD (HEX) AT SAMPLE NUMBER 0171 (HEX)  
YMAX= 40 (HEX) AT SAMPLE NUMBER 0DA7 (HEX)  
YMIN= FD (HEX) AT SAMPLE NUMBER 0171 (HEX)  
ZMAX= 40 (HEX) AT SAMPLE NUMBER 0DA7 (HEX)  
ZMIN= FC (HEX) AT SAMPLE NUMBER 0171 (HEX)  
COMPRESSION STATISTICS:  
NUMBER OF MEMORY BITS AVAILABLE = 021FFC (HEX)  
NUMBER OF BITS AVAILABLE TO VAR LEN CODER = 00040000 (HEX)  
TOTAL NUMBER OF DATA BITS STORED = 0004CD10 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL X = 006652 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Y = 0066A8 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Z = 006160 (HEX)  
NUMBER OF BITS USED TO STORE TIME OR OTHER PARAMETER = 010000 (HEX)  
COMPRESSION RATIO = TOTAL DATA BITS STORED PER MM BITS AVAILABLE

## FILE TA1359PA

CHANNEL X AMPLITUDE DISTRIBUTION  
BYTE VALUE : NUMBER OF OCCURRENCES

80.....	0000 (HEX)
81.....	0000 (HEX)
82.....	0000 (HEX)
83.....	0000 (HEX)
84.....	0000 (HEX)
85.....	0000 (HEX)
86.....	0000 (HEX)
87.....	0000 (HEX)
88.....	0000 (HEX)
89.....	0000 (HEX)
8A.....	0000 (HEX)
8B.....	0000 (HEX)
8C.....	0000 (HEX)
8D.....	0000 (HEX)
8E.....	0000 (HEX)
8F.....	0000 (HEX)
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F0	.....	0000	(HEX)
F1	.....	0000	(HEX)
F2	.....	0001	(HEX)
F3	.....	0001	(HEX)
F4	.....	0003	(HEX)
F5	.....	000C	(HEX)
F6	.....	0012	(HEX)
F7	.....	0016	(HEX)
F8	.....	000F	(HEX)
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FA	.....	000E	(HEX)
FB	.....	000E	(HEX)
FC	.....	001B	(HEX)
FD	.....	004D	(HEX)
FE	.....	0260	(HEX)
FF	.....	0961	(HEX)
00	.....	06A2	(HEX)
01	.....	09A5	(HEX)
02	.....	024A	(HEX)
03	.....	005F	(HEX)
04	.....	0027	(HEX)
05	.....	0011	(HEX)
06	.....	0003	(HEX)
07	.....	0008	(HEX)
08	.....	0008	(HEX)
09	.....	0004	(HEX)
0A	.....	0001	(HEX)
0B	.....	0001	(HEX)
0C	.....	0006	(HEX)
0D	.....	0003	(HEX)
0E	.....	0003	(HEX)
0F	.....	000A	(HEX)
10	.....	0009	(HEX)
11	.....	0000	(HEX)
12	.....	0003	(HEX)
13	.....	0002	(HEX)
14	.....	0000	(HEX)
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18	.....	0000	(HEX)
19	.....	0000	(HEX)
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1B	.....	0000	(HEX)
1C	.....	0000	(HEX)
1D	.....	0000	(HEX)
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59.....	0000	(HEX)
5A.....	0000	(HEX)
5B.....	0000	(HEX)
5C.....	0000	(HEX)
5D.....	0000	(HEX)
5E.....	0000	(HEX)
5F.....	0000	(HEX)
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61.....	0000	(HEX)
62.....	0000	(HEX)
63.....	0000	(HEX)
64.....	0000	(HEX)
65.....	0000	(HEX)
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67.....	0000	(HEX)
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69.....	0000	(HEX)
6A.....	0000	(HEX)
6B.....	0000	(HEX)
6C.....	0000	(HEX)
6D.....	0000	(HEX)
6E.....	0000	(HEX)
6F.....	0000	(HEX)
70.....	0000	(HEX)
71.....	0000	(HEX)
72.....	0000	(HEX)
73.....	0000	(HEX)
74.....	0000	(HEX)
75.....	0000	(HEX)
76.....	0000	(HEX)
77.....	0000	(HEX)
78.....	0000	(HEX)
79.....	0000	(HEX)
7A.....	0000	(HEX)
7B.....	0000	(HEX)
7C.....	0000	(HEX)
7D.....	0000	(HEX)
7E.....	0000	(HEX)
7F.....	0000	(HEX)

EKG SAMPLE COLLECTION STATISTICS : PAGE 1  
FILENAME . . . . . TA1413LU  
SUBJECT . . . . . LUTZ  
SAMPLING RATE . . . . . 500 HZ  
DATE OF COLLECTION . . . . . 23 OCT 80  
TIME OF COLLECTION . . . . . 1413  
COMPRESSION USED . . . . . TOLAN-A  
CHANNEL X ENTROPY . . . . . 3.7826 BITS  
CHANNEL Y ENTROPY . . . . . 3.7950 BITS  
CHANNEL Z ENTROPY . . . . . 3.8149 BITS  
TOTAL SOURCE ENTROPY . . . . . 3.8028 BITS  
PRESS RETURN FOR PAGE 2 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS : PAGE 2  
APPROX MAX COMPRESSION  
RATIO POSSIBLE . . . . . 2.1036 : 1  
COMPRESSION RATIO  
ACHIEVED . . . . . 1.2519 : 1  
ACHIEVED COMPRESSION  
EFFICIENCY . . . . . 22.8 % OF MAXIMUM  
COMPRESSION TIME  
EFFICIENCY OBTAINED . . . . . 95.4 % SMP INTERVAL  
COLLECTION DURATION . . . . . 14.5 SECONDS  
CHANNEL X MAXIMUM . . . . . 1.79687 VOLTS  
CHANNEL X MINIMUM . . . . . -1.0156 VOLTS  
CHANNEL Y MAXIMUM . . . . . 1.75781 VOLTS  
CHANNEL Y MINIMUM . . . . . -1.0156 VOLTS  
CHANNEL Z MAXIMUM . . . . . 1.75781 VOLTS  
CHANNEL Z MINIMUM . . . . . -1.0156 VOLTS  
COMMENTS . . . . . X,Y,Z IN COMMON. LEAD 1  
PRESS RETURN FOR PAGE 3 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS: PAGE 3  
NUMBER OF SAMPLES TAKEN (SAMPNO) = 1C60 (HEX)  
MAXIMUM LOOP COUNT PER INTERRUPT (LPCAL) = 27 (HEX)  
TOTAL WAITING LOOP COUNTS DURING COLLECTION (LOOPCT) = 000031E3 (HEX)  
TIME EFFICIENCY = (1-(LOOPCT\*(SAMPNO\*LPCAL)))\*100  
CHANNEL MAXIMUMS AND MINIMUMS  
XMAX= 2E (HEX) AT SAMPLE NUMBER 0E7F (HEX)  
XMIN= E6 (HEX) AT SAMPLE NUMBER 11C4 (HEX)  
YMAX= 2D (HEX) AT SAMPLE NUMBER 0E7F (HEX)  
YMIN= E6 (HEX) AT SAMPLE NUMBER 11C4 (HEX)  
ZMAX= 2D (HEX) AT SAMPLE NUMBER 0E7F (HEX)  
ZMIN= E6 (HEX) AT SAMPLE NUMBER 11C4 (HEX)  
COMPRESSION STATISTICS:  
NUMBER OF MEMORY BITS AVAILABLE = 021FF0 (HEX)  
NUMBER OF BITS AVAILABLE TO VAR LEN CODER = 00032040 (HEX)  
TOTAL NUMBER OF DATA BITS STORED = 0002A900 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL X = 007786 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Y = 0077F8 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Z = 00764C (HEX)  
NUMBER OF BITS USED TO STORE TIME OR OTHER PARAMETERS = 00CA50 (HEX)  
COMPRESSION RATIO = TOTAL DATA BITS STORED PER MEM. BITS AVAILABLE

FILE TA1413LU  
CHANNEL X AMPLITUDE DISTRIBUTION  
DATA VALUE                            NUMBER OF OCCURENCES

80.....	0000 (HEX)
81.....	0000 (HEX)
82.....	0000 (HEX)
83.....	0000 (HEX)
84.....	0000 (HEX)
85.....	0000 (HEX)
86.....	0000 (HEX)
87.....	0000 (HEX)
88.....	0000 (HEX)
89.....	0000 (HEX)
8A.....	0000 (HEX)
8B.....	0000 (HEX)
8C.....	0000 (HEX)
8D.....	0000 (HEX)
8E.....	0000 (HEX)
8F.....	0000 (HEX)
90.....	0000 (HEX)
91.....	0000 (HEX)
92.....	0000 (HEX)
93.....	0000 (HEX)
94.....	0000 (HEX)
95.....	0000 (HEX)
96.....	0000 (HEX)
97.....	0000 (HEX)
98.....	0000 (HEX)
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9A.....	0000 (HEX)
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9C.....	0000 (HEX)
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A0.....	0000 (HEX)
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AB.....	0000 (HEX)
AC.....	0000 (HEX)
AD.....	0000 (HEX)
AE.....	0000 (HEX)
AF.....	0000 (HEX)
B0.....	0000 (HEX)
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BB.....0000 (HEX)  
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BD.....0000 (HEX)  
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EC.....	0000	(HEX)
ED.....	0000	(HEX)
EE.....	0000	(HEX)
EF.....	0001	(HEX)
F0.....	0000	(HEX)
F1.....	0000	(HEX)
F2.....	0001	(HEX)
F3.....	0006	(HEX)
F4.....	0007	(HEX)
F5.....	0012	(HEX)
F6.....	0012	(HEX)
F7.....	0022	(HEX)
F8.....	0035	(HEX)
F9.....	0050	(HEX)
FA.....	0082	(HEX)
FB.....	00CC	(HEX)
FC.....	0136	(HEX)
FD.....	01D9	(HEX)
FE.....	02EF	(HEX)
FF.....	03D0	(HEX)
00.....	0164	(HEX)
01.....	040A	(HEX)
02.....	02BB	(HEX)
03.....	01E4	(HEX)
04.....	0132	(HEX)
05.....	00AD	(HEX)
06.....	006B	(HEX)
07.....	0058	(HEX)
08.....	0041	(HEX)
09.....	001F	(HEX)
0A.....	0015	(HEX)
0B.....	0013	(HEX)
0C.....	000C	(HEX)
0D.....	000A	(HEX)
0E.....	0002	(HEX)
0F.....	0002	(HEX)
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17.....	0001	(HEX)
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19.....	0000	(HEX)
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1C.....	0000	(HEX)
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49	.....	0000	(HEX)
4A	.....	0000	(HEX)
4B	.....	0000	(HEX)
4C	.....	0000	(HEX)
4D	.....	0000	(HEX)
4E	.....	0000	(HEX)
4F	.....	0000	(HEX)
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76.....	0000	(HEX)
77.....	0000	(HEX)
78.....	0000	(HEX)
79.....	0000	(HEX)
7A.....	0000	(HEX)
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7F.....	0000	(HEX)

EKG SAMPLE COLLECTION STATISTICS : PAGE 1  
FILENAME. . . . . TA1448L  
SAMPLING . . . . . LOWEST IN  
SAMPLING RATE . . . . . 500 HZ  
DATE OF COLLECTION. . . . . 23 OCT 80  
TIME OF COLLECTION. . . . . 1449  
COMPRESSION USED. . . . . TOLAN-A  
CHANNEL X ENTROPY . . . . . 2.9299 BITS  
CHANNEL Y ENTROPY . . . . . 2.9243 BITS  
CHANNEL Z ENTROPY . . . . . 2.9564 BITS  
TOTAL SOURCE ENTROPY. . . . . 2.9436 BITS  
PRESS RETURN FOR PAGE 2 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS : PAGE 2  
APPROX MAX COMPRESSION  
RATIO POSSIBLE. . . . . 2.7176 : 1  
COMPRESSION RATIO  
ACHIEVED. . . . . 1.6008 : 1  
ACHIEVED COMPRESSION  
EFFICIENCY. . . . . 34.9 % OF MAXIMUM  
COMPRESSION TIME  
EFFICIENCY OBTAINED . . . . 90.6 % SMP INTERVAL  
COLLECTION DURATION . . . . 18.5 SECONDS  
CHANNEL X MAXIMUM . . . . . 1.79687 VOLTS  
CHANNEL X MINIMUM . . . . . -0.7421 VOLTS  
CHANNEL Y MAXIMUM . . . . . 1.79687 VOLTS  
CHANNEL Y MINIMUM . . . . . -0.7421 VOLTS  
CHANNEL Z MAXIMUM . . . . . 1.79687 VOLTS  
CHANNEL Z MINIMUM . . . . . -0.7812 VOLTS  
COMMENTS. . . . . X,Y,Z IN COMMON. LEAD 1  
PRESS RETURN FOR PAGE 3 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS: PAGE 3  
NUMBER OF SAMPLES TAKEN (SAMPNO) = 2448 (HEX)  
MAXIMUM LOOP COUNT PER INTERRUPT (LPCAL) = 27 (HEX)  
TOTAL WAITING LOOP COUNTS DURING COLLECTION (LOOPCT) = 0000845C (HEX)  
TIME EFFICIENCY =  $(1 - (LOOPCT * (SAMPNO * LPCAL))) * 100$   
CHANNEL MAXIMUMS AND MINIMUMS  
XMAX= 2E (HEX) AT SAMPLE NUMBER 0D0A (HEX)  
XMIN= ED (HEX) AT SAMPLE NUMBER 1D0A (HEX)  
YMAX= 2E (HEX) AT SAMPLE NUMBER 0D0A (HEX)  
YMIN= ED (HEX) AT SAMPLE NUMBER 1D0A (HEX)  
ZMAX= 2E (HEX) AT SAMPLE NUMBER 0D0A (HEX)  
ZMIN= EC (HEX) AT SAMPLE NUMBER 1D0A (HEX)  
COMPRESSION STATISTICS:  
NUMBER OF MEMORY BITS AVAILABLE = 021FF0 (HEX)  
NUMBER OF BITS AVAILABLE TO VAR LEN CODER = 0003A700 (HEX)  
TOTAL NUMBER OF DATA BITS STORED = 000366C0 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL X = 0069AA (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Y = 00693E (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Z = 006739 (HEX)  
NUMBER OF BITS USED TO STORE TIME OR OTHER PARAMETER = 00EB98 (HEX)  
COMPRESSION RATIO = TOTAL DATA BITS STORED PER MEM BITS AVAILABLE

FILE TA1448LT  
CHANNEL X AMPLITUDE DISTRIBUTION  
DATA VALUE                    NUMBER OF OCCURRENCES

00	.....0000 (HEX)
81	.....0000 (HEX)
82	.....0000 (HEX)
83	.....0000 (HEX)
84	.....0000 (HEX)
85	.....0000 (HEX)
86	.....0000 (HEX)
87	.....0000 (HEX)
88	.....0000 (HEX)
89	.....0000 (HEX)
8A	.....0000 (HEX)
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8C	.....0000 (HEX)
8D	.....0000 (HEX)
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8F	.....0000 (HEX)
90	.....0000 (HEX)
91	.....0000 (HEX)
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95	.....0000 (HEX)
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97	.....0000 (HEX)
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99	.....0000 (HEX)
9A	.....0000 (HEX)
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9D	.....0000 (HEX)
9E	.....0000 (HEX)
9F	.....0000 (HEX)
A0	.....0000 (HEX)
A1	.....0000 (HEX)
A2	.....0000 (HEX)
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A5	.....0000 (HEX)
A6	.....0000 (HEX)
A7	.....0000 (HEX)
A8	.....0000 (HEX)
A9	.....0000 (HEX)
AA	.....0000 (HEX)
AB	.....0000 (HEX)
AC	.....0000 (HEX)
AD	.....0000 (HEX)
AE	.....0000 (HEX)
AF	.....0000 (HEX)
B0	.....0000 (HEX)
B1	.....0000 (HEX)
B2	.....0000 (HEX)

B3.....0000 (HEX)  
B4.....0000 (HEX)  
B5.....0000 (HEX)  
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B7.....0000 (HEX)  
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C0.....0000 (HEX)  
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C8.....0000 (HEX)  
C9.....0000 (HEX)  
CA.....0000 (HEX)  
CB.....0000 (HEX)  
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CD.....0000 (HEX)  
CE.....0000 (HEX)  
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D0.....0000 (HEX)  
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D3.....0000 (HEX)  
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D7.....0000 (HEX)  
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DA.....0000 (HEX)  
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E9.....	0000 (HEX)
EA.....	0000 (HEX)
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F0.....	0000 (HEX)
F1.....	0000 (HEX)
F2.....	0000 (HEX)
F3.....	0000 (HEX)
F4.....	0000 (HEX)
F5.....	0002 (HEX)
F6.....	0001 (HEX)
F7.....	0003 (HEX)
F8.....	0007 (HEX)
F9.....	001C (HEX)
FA.....	0034 (HEX)
FB.....	008F (HEX)
FC.....	0180 (HEX)
FD.....	03F7 (HEX)
FE.....	06A2 (HEX)
FF.....	030D (HEX)
00.....	0714 (HEX)
01.....	03F0 (HEX)
02.....	018A (HEX)
03.....	008A (HEX)
04.....	002A (HEX)
05.....	0013 (HEX)
06.....	0006 (HEX)
07.....	0002 (HEX)
08.....	0002 (HEX)
09.....	0002 (HEX)
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0B.....	0000 (HEX)
0C.....	0000 (HEX)
0D.....	0000 (HEX)
0E.....	0000 (HEX)
0F.....	0000 (HEX)
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13.....	0000 (HEX)
14.....	0000 (HEX)
15.....	0000 (HEX)
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18.....	0000 (HEX)
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1A.....	0000 (HEX)
1B.....	0000 (HEX)
1C.....	0000 (HEX)
1D.....	0000 (HEX)
1E.....	0000 (HEX)

1F	.....	0000	(HEX)
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21	.....	0000	(HEX)
22	.....	0000	(HEX)
23	.....	0000	(HEX)
24	.....	0000	(HEX)
25	.....	0000	(HEX)
26	.....	0000	(HEX)
27	.....	0000	(HEX)
28	.....	0000	(HEX)
29	.....	0000	(HEX)
2A	.....	0000	(HEX)
2B	.....	0000	(HEX)
2C	.....	0000	(HEX)
2D	.....	0000	(HEX)
2E	.....	0000	(HEX)
2F	.....	0000	(HEX)
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31	.....	0000	(HEX)
32	.....	0000	(HEX)
33	.....	0000	(HEX)
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35	.....	0000	(HEX)
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37	.....	0000	(HEX)
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39	.....	0000	(HEX)
3A	.....	0000	(HEX)
3B	.....	0000	(HEX)
3C	.....	0000	(HEX)
3D	.....	0000	(HEX)
3E	.....	0000	(HEX)
3F	.....	0000	(HEX)
40	.....	0000	(HEX)
41	.....	0000	(HEX)
42	.....	0000	(HEX)
43	.....	0000	(HEX)
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45	.....	0000	(HEX)
46	.....	0000	(HEX)
47	.....	0000	(HEX)
48	.....	0000	(HEX)
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4A	.....	0000	(HEX)
4B	.....	0000	(HEX)
4C	.....	0000	(HEX)
4D	.....	0000	(HEX)
4E	.....	0000	(HEX)
4F	.....	0000	(HEX)
50	.....	0000	(HEX)
51	.....	0000	(HEX)
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53	.....	0000	(HEX)
54	.....	0000	(HEX)

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57.....0000 (HEX)  
58.....000C (HEX)  
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5A.....0000 (HEX)  
5B.....0000 (HEX)  
5C.....0000 (HEX)  
5D.....0000 (HEX)  
5E.....0000 (HEX)  
5F.....0000 (HEX)  
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6A.....0000 (HEX)  
6B.....0000 (HEX)  
6C.....0000 (HEX)  
6D.....0000 (HEX)  
6E.....0000 (HEX)  
6F.....0000 (HEX)  
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7A.....0000 (HEX)  
7B.....0000 (HEX)  
7C.....0000 (HEX)  
7D.....0000 (HEX)  
7E.....0000 (HEX)  
7F.....0000 (HEX)

EKG SAMPLE COLLECTION STATISTICS : PAGE 1  
FILENAME . . . . . TAI545T  
COLLECTED . . . . . 1980.10.23  
SAMPLING RATE . . . . . 500 Hz  
DATE OF COLLECTION . . . . . 23 OCT 80  
TIME OF COLLECTION . . . . . 1546  
COMPRESSION USED . . . . . TOLAN-A  
CHANNEL X ENTROPY . . . . . 3.2801 BITS  
CHANNEL Y ENTROPY . . . . . 3.2716 BITS  
CHANNEL Z ENTROPY . . . . . 3.3188 BITS  
TOTAL SOURCE ENTROPY . . . . . 3.3026 BITS  
PRESS RETURN FOR PAGE 2 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS : PAGE 2  
APPROX MAX COMPRESSION  
RATIO POSSIBLE . . . . . 2.4223 : 1  
COMPRESSION RATIO  
ACHIEV'D . . . . . 1.3617 : 1  
ACHIEVED COMPRESSION  
EFFICIENCY . . . . . 25.4 % OF MAXIMUM  
COMPRESSION TIME  
EFFICIENCY OBTAINED . . . . . 93.9 % SMP INTERVAL  
COLLECTION DURATION . . . . . 15.8 SECONDS  
CHANNEL X MAXIMUM . . . . . 3.59375 VOLTS  
CHANNEL X MINIMUM . . . . . -1.4453 VOLTS  
CHANNEL Y MAXIMUM . . . . . 3.59375 VOLTS  
CHANNEL Y MINIMUM . . . . . -1.4453 VOLTS  
CHANNEL Z MAXIMUM . . . . . 3.55468 VOLTS  
CHANNEL Z MINIMUM . . . . . -1.4453 VOLTS  
COMMENTS . . . . . X,Y,Z, IN COMMON, LEAD 1  
PRESS RETURN FOR PAGE 3 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS: PAGE 3  
NUMBER OF SAMPLES TAKEN (SAMPN)= 1EDD (HEX)  
MAXIMUM LOOP COUNT PER INTERRUPT (LPCAL)= 26 (HEX)  
TOTAL WAITING LOOP COUNTS DURING COLLECTION (LOOPCT)= 0000470E (HEX)  
TIME EFFICIENCY =  $(1 - (LOOPCT \cdot (SAMPN \cdot LPCAL))) \cdot 100$   
CHANNEL MAXIMUMS AND MINIMUMS  
XMAX= 5C (HEX) AT SAMPLE NUMBER 03DD (HEX)  
XMIN= DB (HEX) AT SAMPLE NUMBER 145D (HEX)  
YMAX= 5C (HEX) AT SAMPLE NUMBER 03DD (HEX)  
YMIN= DB (HEX) AT SAMPLE NUMBER 145D (HEX)  
ZMAX= 5B (HEX) AT SAMPLE NUMBER 03DD (HEX)  
ZMIN= DB (HEX) AT SAMPLE NUMBER 145D (HEX)  
COMPRESSION STATISTICS:  
NUMBER OF MEMORY BITS AVAILABLE = 021FF0 (HEX)  
NUMBER OF BITS AVAILABLE TO VAR LEN CODER = 00034970 (HEX)  
TOTAL NUMBER OF DATA BITS STORED = 0002E4B8 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL X = 0071DF (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Y = 007199 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Z = 007013 (HEX)  
NUMBER OF BITS USED TO STORE TIME OR OTHER PARAMETER = 00D220 (HEX)  
COMPRESSION RATIO = TOTAL DATA BITS STORED PER MEM BITS AVAILABLE

EKG SAMPLE COLLECTION STATISTICS : PAGE 1

FILENAME. . . . . TA1548T  
POLLING. . . . . 1000000  
SAMPLING RATE. . . . . 300.000  
DATE OF COLLECTION. . . . . 23 OCT 80  
TIME OF COLLECTION. . . . . 1548  
COMPRESSION USED. . . . . TOLAN-A  
CHANNEL X ENTROPY. . . . . 3.0171 BITS  
CHANNEL Y ENTROPY. . . . . 3.0114 BITS  
CHANNEL Z ENTROPY. . . . . 3.0373 BITS  
TOTAL SOURCE ENTROPY. . . . . 3.0355 BITS  
PRESS RETURN FOR PAGE 2 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS : PAGE 2

APPROX MAX COMPRESSION

RATIO POSSIBLE. . . . . 2.6354 : 1

COMPRESSION RATIO

ACHIEVED. . . . . 1.4998 : 1

ACHIEVED COMPRESSION

EFFICIENCY. . . . . 30.5 % OF MAXIMUM

COMPRESSION TIME

EFFICIENCY OBTAINED. . . . . 91.4 % SMP INTERVAL

COLLECTION DURATION. . . . . 17.4 SECONDS

CHANNEL X MAXIMUM. . . . . 2.77343 VOLTS

CHANNEL X MINIMUM. . . . . -1.3671 VOLTS

CHANNEL Y MAXIMUM. . . . . 2.73437 VOLTS

CHANNEL Y MINIMUM. . . . . -1.3671 VOLTS

CHANNEL Z MAXIMUM. . . . . 2.73437 VOLTS

CHANNEL Z MINIMUM. . . . . -1.3671 VOLTS

COMMENTS. . . . . X,Y,Z IN COMMON, LEAD AVL

PRESS RETURN FOR PAGE 3 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS: PAGE 3

NUMBER OF SAMPLES TAKEN (SAMPHO)= 21FE (HEX)

MAXIMUM LOOP COUNT PER INTERRUPT (LPCAL)= 27 (HEX)

TOTAL WAITING LOOP COUNTS DURING COLLECTION (LOOPCT)= 00007134 (HEX)

TIME EFFICIENCY = (1-(LOOPCTS(SAMPHO\*LPCAL)))\*100

CHANNEL MAXIMUMS AND MINIMUMS

XMAX= 47 (HEX) AT SAMPLE NUMBER 0A9E (HEX)

XMIN= DD (HEX) AT SAMPLE NUMBER 1447 (HEX)

YMAX= 46 (HEX) AT SAMPLE NUMBER 095A (HEX)

YMIN= DD (HEX) AT SAMPLE NUMBER 1447 (HEX)

ZMAX= 46 (HEX) AT SAMPLE NUMBER 095A (HEX)

ZMIN= DD (HEX) AT SAMPLE NUMBER 1447 (HEX)

COMPRESSION STATISTICS:

NUMBER OF MEMORY BITS AVAILABLE = 021FF0 (HEX)

NUMBER OF BITS AVAILABLE TO VAR LEN CODER = 00036F C (HEX)

TOTAL NUMBER OF DATA BITS STORED = 00032F00 (HEX)

NUMBER OF BITS USED TO STORE CHANNEL X = 006EE8 (HEX)

NUMBER OF BITS USED TO STORE CHANNEL Y = 006E90 (HEX)

NUMBER OF BITS USED TO STORE CHANNEL Z = 006E9E (HEX)

NUMBER OF BITS USED TO STORE TIME OR OTHER PARAMETER = 000B48 (HEX)

COMPRESSION RATIO = TOTAL DATA BITS STORED PER MEM BITS AVAILABLE

EKG SAMPLE COLLECTION STATISTICS : PAGE 1

FILE NAME . . . . . TA1559B  
SAMPLING RATE . . . . . 500 Hz  
DATE OF COLLECTION . . . . . 23 OCT 80  
TIME OF COLLECTION . . . . . 1559  
COMPRESSION USED . . . . . TOLAN-A  
CHANNEL X ENTROPY . . . . . 3.0165 BITS  
CHANNEL Y ENTROPY . . . . . 3.0038 BITS  
CHANNEL Z ENTROPY . . . . . 3.0261 BITS  
TOTAL SOURCE ENTROPY . . . . . 3.0233 BITS  
PRESS RETURN FOR PAGE 2 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS : PAGE 2

APPROX MAX COMPRESSION  
RATIO POSSIBLE . . . . . 2.6460 : 1  
COMPRESSION RATIO  
ACHIEVED . . . . . 1.5280 : 1  
ACHIEVED COMPRESSION  
EFFICIENCY . . . . . 32.0 % OF MAXIMUM  
COMPRESSION TIME  
EFFICIENCY OBTAINED . . . . . 93.0 % SMP INTERVAL  
COLLECTION DURATION . . . . . 17.7 SECONDS  
CHANNEL X MAXIMUM . . . . . 2.92968 VOLTS  
CHANNEL X MINIMUM . . . . . 0.19531 VOLTS  
CHANNEL Y MAXIMUM . . . . . 2.89062 VOLTS  
CHANNEL Y MINIMUM . . . . . 0.19531 VOLTS  
CHANNEL Z MAXIMUM . . . . . 2.89062 VOLTS  
CHANNEL Z MINIMUM . . . . . 0.15625 VOLTS  
COMMENTS . . . . . X,Y,Z IN COMMON, LEAD 1  
PRESS RETURN FOR PAGE 3 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS: PAGE 3

NUMBER OF SAMPLES TAKEN (SAMPN)= 22A2 (HEX)  
MAXIMUM LOOP COUNT PER INTERRUPT (LPCAL)= 26 (HEX)  
TOTAL WAITING LOOP COUNTS DURING COLLECTION (LOOPCT)= 00005B1E (HEX)  
TIME EFFICIENCY =  $(1 - (LOOPCT * (SAMPN * LPCAL))) * 100$   
CHANNEL MAXIMUMS AND MINIMUMS  
XMAX= 4B (HEX) AT SAMPLE NUMBER 1FA2 (HEX)  
XMIN= 05 (HEX) AT SAMPLE NUMBER 00BD (HEX)  
YMAX= 4A (HEX) AT SAMPLE NUMBER 1FA2 (HEX)  
YMIN= 05 (HEX) AT SAMPLE NUMBER 005C (HEX)  
ZMAX= 4A (HEX) AT SAMPLE NUMBER 1FA2 (HEX)  
ZMIN= 04 (HEX) AT SAMPLE NUMBER 01D5 (HEX)  
COMPRESSION STATISTICS:  
NUMBER OF MEMORY BITS AVAILABLE = 021FF0 (HEX)  
NUMBER OF BITS AVAILABLE TO VAR LEN CODE R = 0003A780 (HEX)  
TOTAL NUMBER OF DATA BITS STORED = 00333F30 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL X = 006BE5 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Y = 006AC9 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Z = 006857 (HEX)  
NUMBER OF BITS USED TO STORE TIME OR OTHER PARAMETER = 00E9E0 (HEX)  
COMPRESSION RATIO = TOTAL DATA BITS STORED PER MEM BITS AVAILABLE

:EKG SAMPLE COLLECTION STATISTICS : PAGE 1  
FILENAME . . . . . TA1511S  
SAMPLING RATE . . . . . 500 Hz  
DATE OF COLLECTION . . . . . 1511  
TIME OF COLLECTION . . . . . 1511  
COMPRESSION USED . . . . . TOLAN-A  
CHANNEL X ENTROPY . . . . . 3.3255 BITS  
CHANNEL Y ENTROPY . . . . . 3.3341 BITS  
CHANNEL Z ENTROPY . . . . . 3.3560 BITS  
TOTAL SOURCE ENTROPY . . . . . 3.3457 BITS  
PRESS RETURN FOR PAGE 2 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS : PAGE 2  
APPROX MAX COMPRESSION  
RATIO POSSIBLE . . . . . 2.3911 : 1  
COMPRESSION RATIO  
ACHIEVED . . . . . 1.4319 : 1  
ACHIEVED COMPRESSION  
EFFICIENCY . . . . . 31.0 % OF MAXIMUM  
COMPRESSION TIME  
EFFICIENCY OBTAINED . . . . . 94.0 % SMP INTERVAL  
COLLECTION DURATION . . . . . 16.6 SECONDS  
CHANNEL X MAXIMUM . . . . . 1.32812 VOLTS  
CHANNEL X MINIMUM . . . . . -2.0703 VOLTS  
CHANNEL Y MAXIMUM . . . . . 1.32812 VOLTS  
CHANNEL Y MINIMUM . . . . . -2.0703 VOLTS  
CHANNEL Z MAXIMUM . . . . . 1.32812 VOLTS  
CHANNEL Z MINIMUM . . . . . -2.0703 VOLTS  
COMMENTS . . . . . X,Y,Z IN COMMON. LEAD AVL  
PRESS RETURN FOR PAGE 3 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS: PAGE 3  
NUMBER OF SAMPLES TAKEN (SAMPN0) = 2074 (HEX)  
MAXIMUM LOOP COUNT PER INTERRUPT (LPCAL) = 27 (HEX)  
TOTAL WAITING LOOP COUNTS DURING COLLECTION (LOOPCT) = 00004ADC (HEX)  
TIME EFFICIENCY =  $(1 - (LOOPCT * (SAMPN0 * LPCAL))) * 100$   
CHANNEL MAXIMUMS AND MINIMUMS  
XMAX= 22 (HEX) AT SAMPLE NUMBER 1171 (HEX)  
XMIN= CB (HEX) AT SAMPLE NUMBER 021D (HEX)  
YMAX= 22 (HEX) AT SAMPLE NUMBER 1423 (HEX)  
YMIN= CB (HEX) AT SAMPLE NUMBER 021D (HEX)  
ZMAX= 22 (HEX) AT SAMPLE NUMBER 1423 (HEX)  
ZMIN= CR (HEX) AT SAMPLE NUMBER 00E8 (HEX)  
COMPRESSION STATISTICS:  
NUMBER OF MEMORY BITS AVAILABLE = 021FF0 (HEX)  
NUMBER OF BITS AVAILABLE TO VAR LEN CODER = 000390C (HEX)  
TOTAL NUMBER OF DATA BITS STORED = 00030AE0 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL X = 00717A (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Y = 00713D (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Z = 006127 (HEX)  
NUMBER OF BITS USED TO STORE TIME OR OTHER PARAMETERS = 001000 (HEX)  
COMPRESSION RATIO = TOTAL DATA BITS STORED PER MEM BITS AVAILABLE

EKG SAMPLE COLLECTION STATISTICS : PAGE 1  
FILENAME . . . . . TA1520B  
SUBJCT . . . . . BALSALO  
SAMPLING RATE . . . . . 500 Hz  
DATE OF COLLECTION . . . . . 23 OCT 80  
TIME OF COLLECCTION . . . . . 1520  
COMPRESSION USED . . . . . TOLAN-A  
CHANNEL X ENTROPY . . . . . 2.6012 BITS  
CHANNEL Y ENTROPY . . . . . 2.5808 BITS  
CHANNEL Z ENTROPY . . . . . 2.6229 BITS  
TOTAL SOURCE ENTROPY . . . . . 2.6104 BITS  
PRESS RETURN FOR PAGE 2 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS : PAGE 2  
APPROX MAX COMPRESSION  
RATIO POSSIBLE . . . . . 3.0645 : 1  
COMPRESSION RATIO  
ACHIEVED . . . . . 1.7247 : 1  
ACHIEVED COMPRESSION  
EFFICIENCY . . . . . 35.1 % OF MAXIMUM  
COMPRESSION TIME  
EFFICIENCY OBTAINED . . . . . 90.9 % SMP INTERVAL  
COLLECTION DURATION . . . . . 20.0 SECONDS  
CHANNEL X MAXIMUM . . . . . 2.07031 VOLTS  
CHANNEL X MINIMUM . . . . . -0.5468 VOLTS  
CHANNEL Y MAXIMUM . . . . . 2.07031 VOLTS  
CHANNEL Y MINIMUM . . . . . -0.5468 VOLTS  
CHANNEL Z MAXIMUM . . . . . 2.07031 VOLTS  
CHANNEL Z MINIMUM . . . . . -0.5468 VOLTS  
COMMENTS . . . . . X,Y,Z IN COMMON. LFA D 2  
PRESS RETURN FOR PAGE 3 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS: PAGE 3  
NUMBER OF SAMPLES TAKEN (SAMENO) = 2717 (HEX)  
MAXIMUM LOOP COUNT PER INTERRUPT (LPCAL) = 27 (HEX)  
TOTAL WAITING LOOP COUNTS DURING COLLECTION (LOOPCT) = 00008A6D (HEX)  
TIME EFFICIENCY =  $(1 - (LOOPCT * (SAMENO * LPCAL))) * 100$   
CHANNEL MAXIMUMS AND MINIMUMS  
XMAX= 35 (HEX) AT SAMPLE NUMBER 1F9E (HEX)  
XMIN= F2 (HEX) AT SAMPLE NUMBER 0N63 (HEX)  
YMAX= 35 (HEX) AT SAMPLE NUMBER 1F9E (HEX)  
YMIN= F2 (HEX) AT SAMPLE NUMBER 054P (HEX)  
ZMAX= 35 (HEX) AT SAMPLE NUMBER 1F9E (HEX)  
ZMIN= F2 (HEX) AT SAMPLE NUMBER 054B (HEX)  
COMPRESSION STATISTICS:  
NUMBER OF MEMORY BITS AVAILABLE = 021FF0 (HEX)  
NUMBER OF BITS AVAILABLE TO VAR LEN CODER = 0002F0 (HEX)  
TOTAL NUMBER OF DATA BITS STORED = 0002AA28 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL X = 0066DE (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Y = 006666 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Z = 0063CC (HEX)  
NUMBER OF BITS USED TO STORE TIME OR OTHER INFORMATION = 00F630 (HEX)  
COMPRESSION RATIO = TOTAL DATA BITS STORED DIVIDED BY BITS AVAILABLE

EKG SAMPLE COLLECTION STATISTICS : PAGE 1  
FILENAME. . . . . . . . . . TAI439ST  
SUBJECT . . . . . . . . . . STROUP  
SAMPLING RATE . . . . . . . . 500 HZ  
DATE OF COLLECTION. . . . . . . . 23 OCT 80  
TIME OF COLLECTION. . . . . . . . 1439  
COMPRESSION USED. . . . . . . . TOLAN-A  
CHANNEL X ENTROPY . . . . . . . . 3.2669 BITS  
CHANNEL Y ENTROPY . . . . . . . . 3.2756 BITS  
CHANNEL Z ENTROPY . . . . . . . . 3.3094 BITS  
TOTAL SOURCE ENTROPY. . . . . . . . 3.2912 BITS  
PRESS RETURN FOR PAGE 2 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS : PAGE 2  
APPROX MAX COMPRESSION  
RATIO POSSIBLE. . . . . . . . . . 2.4307 : 1  
COMPRESSION RATIO  
ACHIEVED. . . . . . . . . . 1.4314 : 1  
ACHIEVED COMPRESSION  
EFFICIENCY. . . . . . . . . . 30.1 % OF MAXIMUM  
COMPRESSION TIME  
EFFICIENCY OBTAINED . . . . . 93.4 % SMP INTERVAL  
COLLECTION DURATION . . . . . 16.6 SECONDS  
CHANNEL X MAXIMUM . . . . . . . . 3.35937 VOLTS  
CHANNEL X MINIMUM . . . . . . . . -0.5468 VOLTS  
CHANNEL Y MAXIMUM . . . . . . . . 3.35937 VOLTS  
CHANNEL Y MINIMUM . . . . . . . . -0.5468 VOLTS  
CHANNEL Z MAXIMUM . . . . . . . . 3.35937 VOLTS  
CHANNEL Z MINIMUM . . . . . . . . -0.5859 VOLTS  
COMMENTS. . . . . . . . . . X,Y,Z IN COMMON. LEAD 1  
PRESS RETURN FOR PAGE 3 OF STATISTICS=

EKG SAMPLE COLLECTION STATISTICS: PAGE 3  
NUMBER OF SAMPLES TAKEN (SAMENO)= 2071 (HEX)  
MAXIMUM LOOP COUNT PER INTERRUPT (LPCAL)= 27 (HEX)  
TOTAL WAITING LOOP COUNTS DURING COLLECTION (LOOPCT)= 0000526C (HEX)  
TIME EFFICIENCY = (1-(LOOPCTS(SAMENO\*LPCAL)))\*100  
CHANNEL MAXIMUMS AND MINIMUMS  
XMAX= 56 (HEX) AT SAMPLE NUMBER 1472 (HEX)  
XMIN= F2 (HEX) AT SAMPLE NUMBER 1E5A (HEX)  
YMAX= 56 (HEX) AT SAMPLE NUMBER 1472 (HEX)  
YMIN= F2 (HEX) AT SAMPLE NUMBER 0FD8 (HEX)  
ZMAX= 56 (HEX) AT SAMPLE NUMBER 1472 (HEX)  
ZMIN= F1 (HEX) AT SAMPLE NUMBER 1E5A (HEX)  
COMPRESSION STATISTICS:  
NUMBER OF MEMORY BITS AVAILABLE = 021FF0 (HEX)  
NUMBER OF BITS AVAILABLE TO VAR LFM CODER = 00037700 (HEX)  
TOTAL NUMBER OF DATA BITS STORED = 00030A98 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL X = 006E82 (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Y = 006EDC (HEX)  
NUMBER OF BITS USED TO STORE CHANNEL Z = 006CF0 (HEX)  
NUMBER OF BITS USED TO STORE TIME OR OTHER PARAMETERS = 00DPC0 (HEX)  
COMPRESSION RATIO = TOTAL DATA BITS STORED PER UNIT TIME AVAILABLE

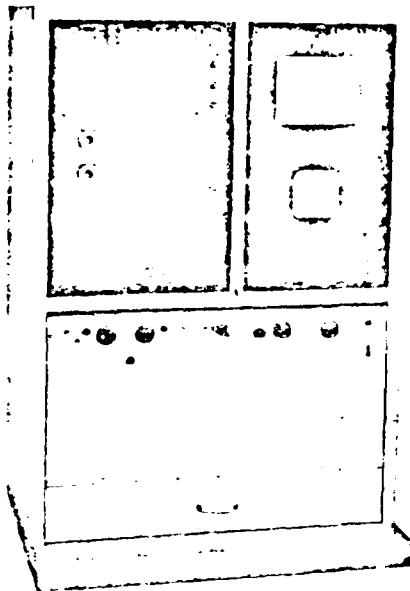
Appendix E

This appendix contains a photocopy of the specifications for the pertinent equipment used in this thesis.

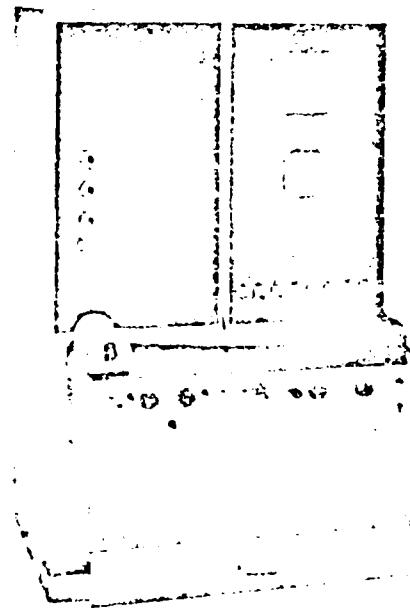
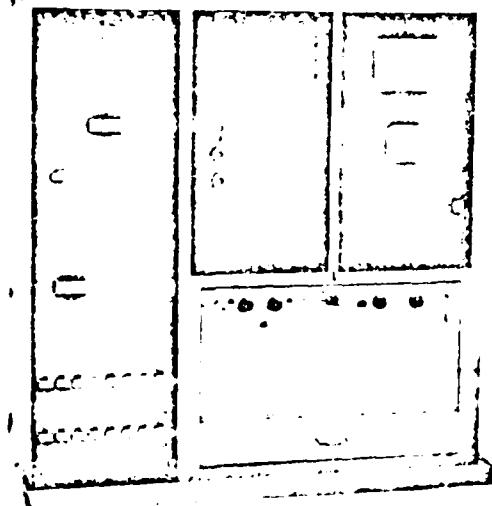
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# RESEARCH RECORDER



## OPERATING INSTRUCTIONS



ELECTRONICS FOR MEDICINE INC

30 VIRGINIA ROAD, WHITE PLAINS, N.Y. 10603 • TEL. 944-4187

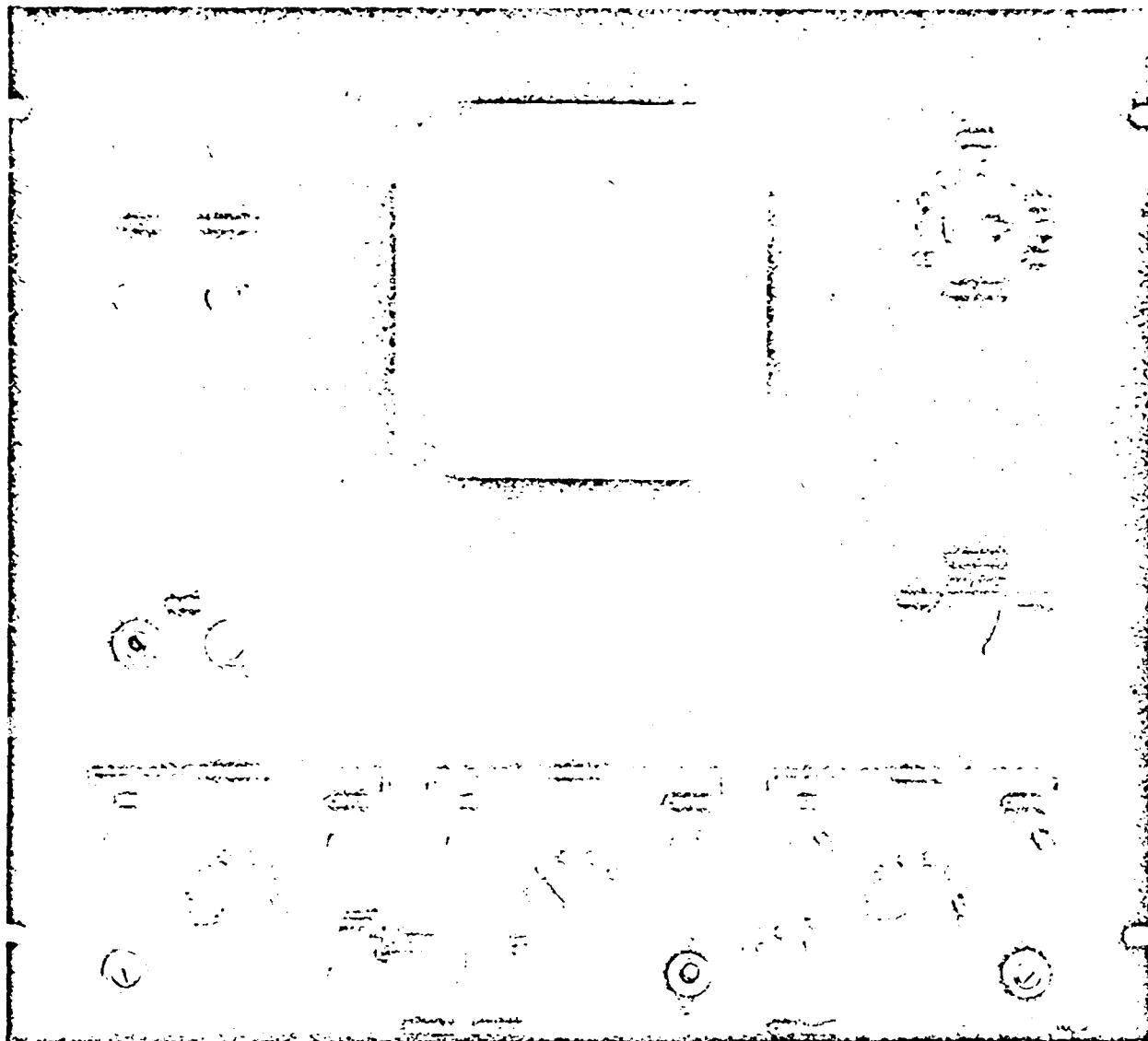
## **DR-8 RESEARCH RECORDER MONITOR**

**The Model DR-8 Research Recorder is designed to operate from a 100-130 volt, 60 Hz line. Power consumption is about 700 watts. The circuit is protected by a 10 ampere circuit breaker which is incorporated in the power ON-OFF switch on the camera panel.**

**It is essential for safety and to minimize 60 Hz interference in the tracings that the Research Recorder be connected to a good "ground". A third wire is brought out at the power plug for this purpose. A good "ground" can usually be obtained at a cold water pipe.**

**The instrument combines as many as eight amplifiers with a cathode ray tube recording camera and two cathode ray tube monitors. The multitrace monitor displays any number of traces up to eight. The single monitor displays any one of the traces enlarged or reduced in size. It is also used for the presentation of loops, such as the vectoreardiogram or the lung compliance and lung resistance loops, and for balancing strain gauge amplifiers. All phenomena can be monitored continuously. Recording is done at will, merely by turning a switch. The display on the cathode ray tube in the camera can be identical to that on either the multitrace tube or the single trace tube. It is therefore possible to photograph either scalar tracings or loops. In actual operation, size adjustments and preliminary measurements are made by observing the monitor screens without recording. Thus, the camera runs for only short intervals.**

**The multitrace tube has only one beam, but because the amplifiers are connected to it through an electronic switch operating at a rate of 45,000 Hz/sec., it is possible to display as many as eight traces simultaneously. The beam switches rapidly from one trace to the next, with blanking of the cathode ray tube taking place during the switching period. Thus, each beam is on screen only 1/8th of the time, although giving the impression of being continuous. A switch is provided to make possible more frequent sampling for high speed phenomena, by reducing the number of channels displayed.**



**POWER** switch, located on the right side of camera panel, turns the Research Recorder on. A minute or two is required before the traces on the monitor screens are stabilized. Adjustments may now be made.

**INTENSITY** and **FOCUS** control adjust the cathode ray tube beams. The upper two controls are for the multitrace tube, the lower two for the single trace tube. The intensity is set according to the brightness of the room in which the instrument is being used. The traces will be sharper at lower intensities. Focus is adjusted for the sharpest spot and should be done when the beams are approximately 1/3rd of the distance from either side of the tube to allow best overall focusing.

**FUNCTION** selector switch determines the presentation on the single trace tube; the multitrace tube will display scalar traces continuously. The single trace tube will display the trace from any amplifier when the FUNCTION selector switch is in the SWEEP position. In the BALANCE position, a pattern is displayed on the single trace screen which permits rapid balancing of the pressure transducers. In LOOPS, any channel can be presented on the horizontal axis of the single trace tube, and at the same time any other channel can be presented on the vertical axis. This position is used for the VCG - ops (X-Y plot).

SWEET switch determines the speed at which the traces move across the screen. The proper setting depends on the repetition rate of the phenomena displayed. Nine monitoring speeds are available: 2.5, 5, 10, 25, 50, 75, 100, 150 and 200 mm/sec.

The MULTITRACE MONITOR permits the display of up to eight channels at one time. The positions of the individual traces on the multitrace tube are controlled by the CENTER knob on the individual channels. Each of the traces can be positioned anywhere on the screen, superimposed on other traces, or can be moved off screen in either direction if it is not being used. The amplitude of the traces on the multitrace screen will depend upon the setting of the AMPLITUDE controls on the individual amplifiers.

The CHANNELS switch determines the number of amplifiers that will be displayed. The multitrace presentation is the result of an 8 channel electron switch, whose basic frequency is 48 K Hz. In the 1-8 position, each amplifier will be sampled equally 6,000 times second, or every 166 microseconds. In the 1-5 position, channels 1, 2 and 3 are sampled at 12,000 times second or every 83 microseconds. Channels 4 and 5 are sampled at the normal 6,000 times and channels 6, 7 and 8 are not displayed at all. Placing the multitrace switch at 1-6 will sample channels 1 and 2 at the 12,000 rate and 3, 4, 5 and 6 at the normal 6,000 rate. Channels 7 and 8 will not appear on the multitrace screen. Amplifiers providing fast rise time phenomena, such as heart sounds, action and nerve potentials, should be in positions 1, 2 and 3 to take advantage of the faster sampling rate and to minimize the dotting effect.

When the SUBTRACT SELECTOR switch is in the OFF position, the built-in baseline, gradient and subtract functions are inoperative. Turning switch to BASELINE will allow built in baseline to be positioned to any point on screen with POSITION control. An event MARK button and jack (for marking from a remote position) are provided. To display either baseline or gradient 1-5, 1-6, 1-8 switch must be in 1-8, since baseline or gradient trace replace channel 8. When BASELINE or GRADIENT position is selected channel 8 amplifier will not be seen.

A PRESSURE GRADIENT can be displayed on the multitrace screen together with the two pressures. The adjustment procedure is as follows:

- 1) Attach pressure transducers to pressure amplifiers; allow them to warm up, balance the amplifiers. Position the traces to baseline from which measurements are to be made; equalize the amplitudes. (See SGA-SGM instructions)
- 2) Turn the SUBTRACT SELECTOR switch to "GRADIENT". The gradient trace will now be substituted for Channel 8 and its position controlled by the baseline POSITION control.
- 3) Turn HORIZONTAL SIZE and SUBTRACT SIZE to zero (counterclockwise)

- 4) Select the channel to be subtracted from with the Horizontal switch; the channel to be subtracted with the Subtract switch.
- 5) Position the gradient trace to the baseline position of the two pressure channels. There are now 3 or 4 lines superimposed i.e., 2 pressure amplifiers, the difference trace (gradient) and optionally a baseline, which can be a baseline, marker trace, or any unused pressure amplifier.
- 6) Turn the +/- switch to +.
- 7) Turn HORIZONTAL SIZE slowly clockwise; note in which direction trace moves. If it moves upwards, turn HORIZONTAL CENTER counter-clockwise to return trace to its original position. Advance HORIZONTAL SIZE fully clockwise and readjust HORIZONTAL CENTER, if necessary. When the center is optimally adjusted, turning HORIZONTAL SIZE should not affect the position of trace. Return HORIZONTAL SIZE to zero. Repeat procedure with SUBTRACT SIZE and CENTER controls, returning SIZE to zero after obtaining balance.
- 8) Turn CALIBRATE control on pressure amplifier which was selected by HORIZONTAL SWITCH so that its trace moves approximately full screen; adjust HORIZONTAL SIZE until gradient trace coincides with the pressure trace. Remove the CALIBRATE signal from the pressure amplifier and both lines will return to their original positions.
- 9) Turn the CALIBRATE on that amplifier selected by SUBTRACT switch until the pressure trace moves approximately full screen. Adjust SUBTRACT SIZE until lines are superimposed. Return CALIBRATE switch to zero and both lines will return to their original positions.
- 10) The "Gradient" trace will now represent the sum of the two pressures. To observe the gradient, turn the - - switch to -. Readjust gradient POSITION, if necessary.

For the H and V sections of SUBTRACT SELECTOR switch, see loops.

SINGLE TRACE MONITOR presentation is determined by the FUNCTION switch. Any individual amplifier can be displayed when the FUNCTION switch is in the SWEEP position. Turn the VERTICAL selector switch to the channel number representing the desired amplifier. Make certain this amplifier is visible on multitrace screen. The position of the selected trace can be controlled by the CENTER control on the channel, or by the VERTICAL CENTER control. The amplitude on the screen is adjusted by the VERTICAL SIZE control and can be 2 1/2 times the amplitude of the corresponding trace on the multitrace tube, or it

can be reduced as low as zero amplitude. When the SIZE control is at zero (completely counterclockwise), the CENTER control will have no effect. If trace is more than 1" from mid-screen, (SIZE control fully counterclockwise) turn screwdriver adjustment below VERTICAL CENTER to return baseline to mid-screen. Turning the SIZE control clockwise may move trace off screen. Note in which direction it moves (up or down) and return trace to screen with VERTICAL CENTER control. If trace moves upward, turn CENTER counterclockwise; if trace moves downward, turn CENTER clockwise.

The BALANCE position of the FUNCTION switch will display a lissajous figure, which is used as an aid in balancing a transducer in a pressure amplifier.

Turn VERTICAL selector switch to pressure amplifier (with transducer connected) to be balanced. The horizontal deflection represents the transducer excitation voltage, the vertical represents the output voltage from the transducer through a pressure amplifier.

The vertical SIZE and CENTER controls do not have any effect in the BALANCE position. When the transducer is brought into balance, there will not be any output voltage, and the lissajous presentation will be a straight horizontal line.

In LOOPS position, the VERTICAL selector switch selects the channel to be presented on the vertical axis, and the vertical SIZE and CENTER controls are used to adjust vertical amplitude and position.

The HORIZONTAL switch selects the channel to be presented on the horizontal axis, and the horizontal SIZE and CENTER controls are used to adjust horizontal amplitude and position. When the HORIZONTAL (or VERTICAL) SIZE control is turned to minimum counterclockwise, the CENTER controls will not have any effect. A dot should appear at mid-screen. If it is more than 1" from mid-screen, adjust screwdriver controls below VERTICAL CENTER or HORIZONTAL SIZE until dot is at mid-screen.

Advance HORIZONTAL SIZE clockwise and adjust HORIZONTAL CENTER to return dot to its original position near mid-screen. Proper adjustment of center is achieved when it is possible to turn horizontal SIZE from minimum counterclockwise to maximum clockwise without moving dot when there isn't any applied signal. Do the same with the VERTICAL and SUBTRACT switches: return the subtract SIZE to minimum counterclockwise.

The H or V positions of SUBTRACT SELECTOR make it possible to add or subtract the signals from two channels on either the horizontal or vertical axis when measuring loops. This is especially useful when measuring Lung Compliance or Lung resistance. For example: Let us suppose it is desired to subtract from the horizontal axis. The HORIZONTAL switch selects the amplifier to be subtracted from. Then, the SUBTRACT selector switch is used to select the channel to be subtracted. + - switch is turned to -. The SUBTRACT SELECTOR is turned from off to horizontal. The SUBTRACT, SIZE and CENTER controls will now also control the horizontal amplitude and position but in a direction opposite to that of the VERTICAL controls. See instructions for Lung Resistance and Lung Compliance loops.

Timing and direction marking can be introduced to the loop patterns by turning the camera selector to LOOPS and choosing the appropriate TIMING rate. The camera ON-OFF switch need not be ON to display timing on loop pattern, even at when photographing loop. Camera timing should be in .004 or .02 to properly display LC/LR loops; or .004 for VCG.

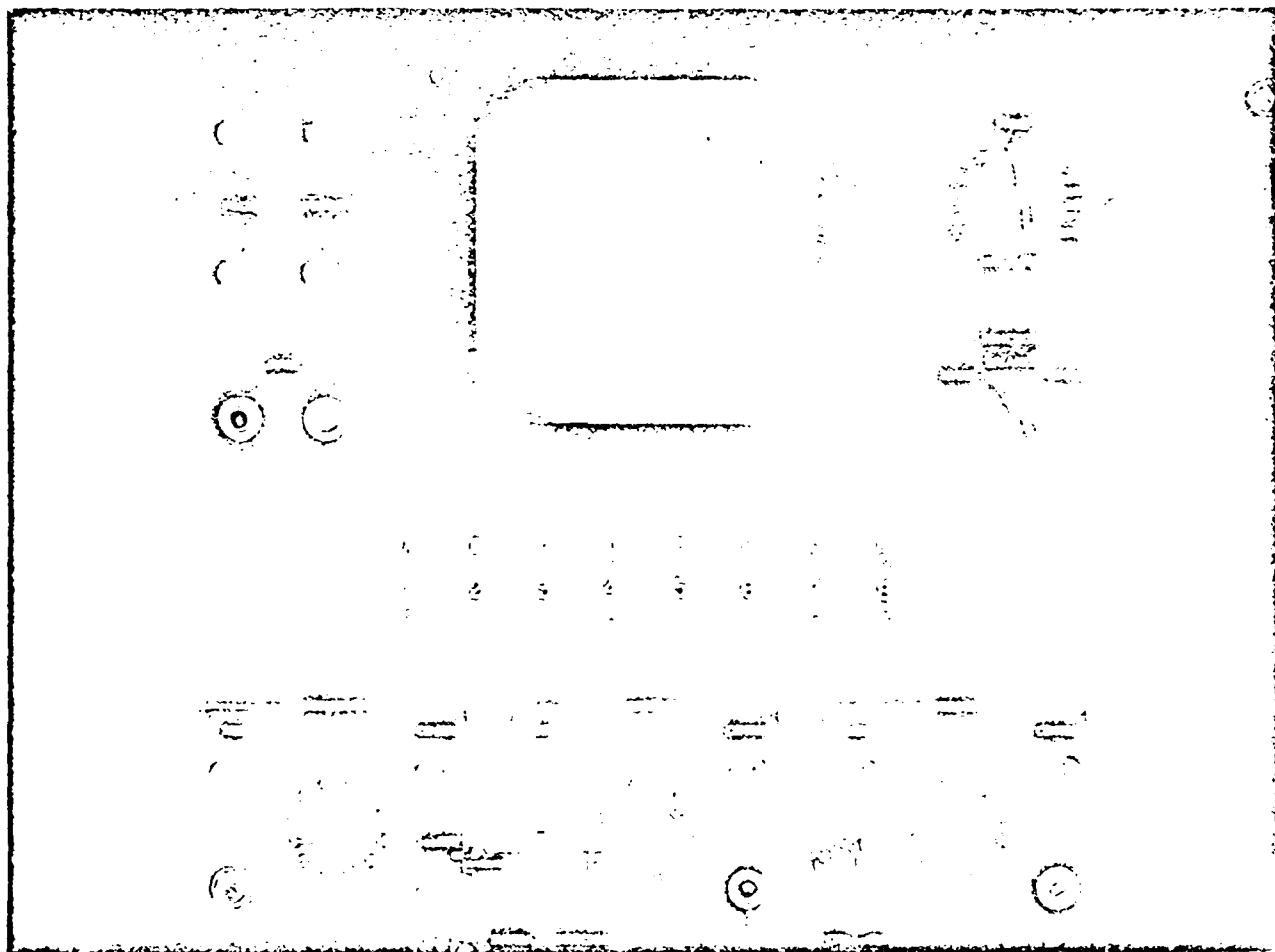
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MONITOR  
MODEL DR-12

The DR-12 recorder is a modified version of the DR-8. The DR-8 instruction plus the following information covers the operation of the DR-12.

The amplifiers in the DR-12 are numbered from 1 through 12. Ten of these, 1-10, starting from the top, are located on the left side of the recorder; numbers 11 and 12 are located beneath the monitor panel. The eight displayable channels of the recorder are lettered from A through H with a dial type switch supplied for each channel which permits any of the 12 amplifiers to be selected on any of the 8 channels. Since it is possible to select the same amplifier on more than one channel, the user should be certain that each of the dials is set to a different number or the same trace will appear on the screen more than once.



The A-E, A-F, A-H switch on the monitor panel can be used to sample two or three channels more frequently than the others. The electronic switching rate is 48,000 Hz/second; each of the 8 channels is sampled 6,000 times per second when the switch is in the A-H position. In the A-E position, channels A, B, and C are sampled 12,000 times/second. The channels D and E are sampled 6,000 times/second. Only

5 channels are operative in this position. In the A-F position, channels A and B are sampled 12,000 times/second while channels C, D, E and F are sampled 6,000 times/second. Operation is limited to 6 channels in this position. A and B can only be sampled at these higher rates by selecting them on channels A or B.

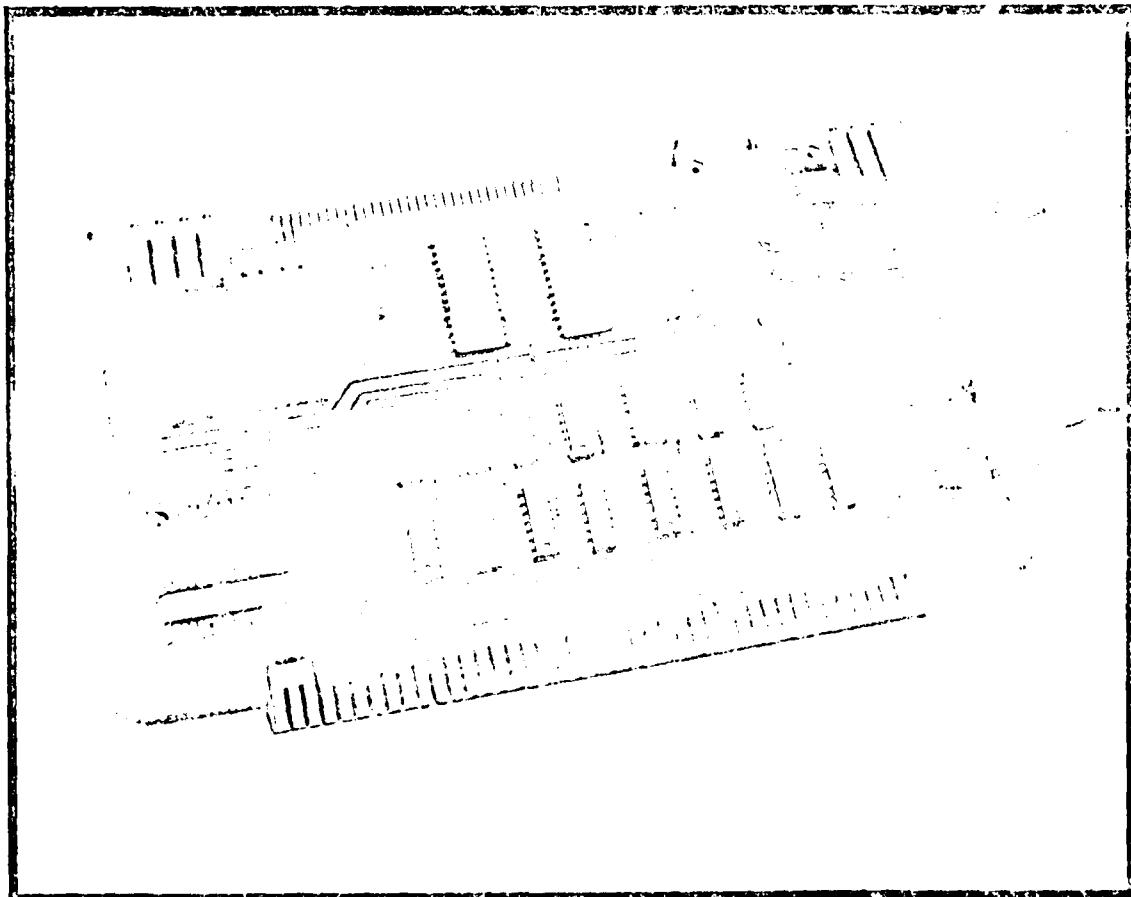
SUBTRACT SELECTOR switch marked BASELINE, GRADIENT, OFF, H and V, when positioned on BASELINE or GRADIENT, will remove channel H (A-E, A-F, A-H switch must be on A-H) from display on multitrace screen and replace it with a baseline or gradient trace. With switch in BASELINE, the POSITION control, below the switch, can be used to set the baseline to any desired point. A MARK push button and remote MARK jack can be used as an event marker. A hand/foot switch can be patched into MARK jack, to mark events from a remote point. Placing switch on GRADIENT position will permit the display of gradients using built in gradient amplifier (POSITION control is the same as for baseline). The subtraction is done horizontally. See section on "subtraction" in DR-8 instructions. The H and V position of the SUBTRACT SELECTOR switch, permit subtraction on vector screen, either vertically or horizontally in the OFF position, subtract or GRADIENTS and BASELINE are inoperative. See DR-8 instructions.

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SINETRAC SERIES  
MODEL ST-6800  
A/D-D/A PERIPHERAL SYSTEMS  
INSTRUCTION MANUAL  
Part No. 58-12140-25

JANUARY 1979



SPECIFICATIONS

DATA ACQUISITION SECTION

Specifications

Typical @ +25°C, dynamic conditions, unless otherwise specified.

Analog Inputs

Number of Channels	32 Single-ended or 16 differential
Channel Expansion	Up to 128 single ended or 112 differential using ADC-Slave Expander Cards (Model ST-6800 ADX)
Full Scale Input Ranges	0 to +5 Volts 0 to +10 Volts -5 to +5 Volts -10 to +10 Volts
Current Input Channel Range (8 Channels)	4-20 mA type
Common Mode Range	+10 Volts
Input Overvoltage	+35 Volts Max. continuous
Input Impedance	100 Megohms differential or to ground
Input Bias Current	3nA typ., 10nA max.
Input Capacitance	5pF, OFF CHANNEL, 100 pF ON CHANNEL to ground

Performance

Accuracy @ +25C	Within $\pm 0.025\%$ of input range
Resolution	12 Binary bits (1 part in 4096)
Nonlinearity	$\pm 1/2$ LSB maximum
Differential non-linearity	$\pm 1/2$ LSB maximum
Gain Error	Adjustable to zero
Offset or Zero Error	Adjustable to zero
Gain Temperature Drift (Bipolar)	Within $\pm 10$ ppm of FSR/°C
Zero Temperature Drift (Unipolar)	Within $\pm 5$ ppm of FSR/°C max.
Common Mode Rejection	70 dB min, DC to 1 kHz with $1k\Omega$ unbalance
Power Supply Rejection	100 dB to +5V bus

Dynamic Characteristics

Typical Data Transfer

I/O Period (Total)	36 microseconds
Throughput Period	20 microseconds
Acquisition Time	8 microseconds
A/D Conversion Time	12 microseconds
Aperture Time	100 nanoseconds
Sample/Hold Switch	.01% Max.
Feedthrough	
MUX Crosstalk from OFF Channels	.007% @ 1 kHz, $R_s = 1\text{K}$

Digital Outputs

Output Coding

Straight Binary (Unipolar) } Jumper  
Offset Binary (Bipolar) } Selected  
2's Complement (Bipolar) }

Output Format

2-Byte group electrically compatible to Motorola's EXORciser bidirectional bus. Sign extension is jumper selected on bits 12 thru 15 for 2's complement units. Bits 12 thru 15 are logic zero for all other units.

Channel Addressing

Random channel addressing may be started by external interrupt input for event operation or by internal program control.

Base Address

Prewired by PC Board jumpers for one of 128 base addresses.

Data Distribution Section (D/A Analog Outputs)

Number of Channels

2 (Expandable only by stand-alone ST-6800DA Boards.)

Resolution

12 Bits

Full Scale Output

0 to +5 Volts

Voltage Ranges

0 to +10 Volts  
-5 to +5 Volts  
-10 to +10 Volts

Input Coding

Straight Binary (Unipolar)  
Offset Binary (Bipolar)  
2's Complement (Bipolar)

Output Impedance

.05 ohm

Output Current

$\pm 5 \text{ mA min}$

ST-6800  
M68-Dhi705

Performance

Non linearity	$\pm 1/2$ LSB, maximum
Differential Nonlinearity	$\pm 1/2$ LSB, maximum
Gain Error	Adjustable to zero using pot. for each channel
Offset or Zero Error	Adjustable to zero using pot for each channel
Gain Temperature Drift	$\pm 20$ ppm of output/ $^{\circ}$ C
Zero Temperature Drift	(Unipolar output) $\pm 5$ ppm of FSR/ $^{\circ}$ C
Offset Temperature Drift	(Bipolar output) $\pm 10$ ppm of FSR/ $^{\circ}$ C
Settling Time (20V change)	4 microseconds to $\pm 1/2$ LSB
Slew Rate	20V/usec
Power Supply Rejection	$\pm 0.02\%$ of FSR per 1% variation

Power Consumption

1.2 ampstypical @ +5 vdc supplied from MPU bus connector. On-board DC to DC Convert. supplies  $\pm 15$  vdc to linear circuits.

Physical

Operating Temperature Range	0 $^{\circ}$ to +70 $^{\circ}$ C
Storage Temperature Range	-25 $^{\circ}$ C to +85 $^{\circ}$ C
Card Size	9.75"W x 5.75"H x .062"D

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M. Allen

AVT/55318

# IBRUSH MARK II

## OPERATING INSTRUCTIONS

Brush Recorder Mark II  
Model RD 2522 20

## General Information

### Specifications

Sensitivity .....	10 millivolts per chart line (mm). Full scale deflection from chart center : 200 millivolts.
Sensitivity steps .....	.01, .02, .05, .1, .2, .5, 1, 2, 5 and 10 volts per chart line (mm). Maximum attenuator error 1% with balanced input.
Measurement range	
Single-ended input (5 megohm) .....	.010 volt to 400 volts
Balanced input (10 megohm) .....	.010 volt to 400 volts, side to side, allowable voltage off-ground at any attenuator step is 1000 x volts per chart line switch setting up to 500 volts off-ground.
Common mode rejection .....	Better than 1000 to 1, attenuator set in .01 volt per chart line position.
Zero line stability .....	Less than 1.4 chart line (mm) per hour. Total drift over eight hour period not more than 1.2 chart line (mm).
Noise .....	Not noticeable on chart with shorted input.
Frequency response .....	The recorded peak-to-peak amplitude of a constant voltage sine wave will be within 1 chart line (mm) of a nominal 10 lines from D.C. to 100 cps.
Maximum amplitude .....	40 lines (mm) peak-to-peak, D.C. to 40 cps. 20 lines (mm) peak-to-peak, D.C. to 70 cps. 10 lines (mm) peak-to-peak, D.C. to 100 cps.

## General Information

### Specifications

Pen Bias . . . . .	Permits positioning of pen on chart, $\pm 20$ chart lines (mm). Effective for either single-ended or balanced input.
Trace linearity . . . . .	D. C. within 2% full chart width. A. C. within 3% full chart width, any frequency within limits of maximum amplitude for electric writing.
Trace width . . . . .	0.006 with Model RA 2E22-31 pen.
Writing method . . . . .	Electric stylus.
Number of recording channels . . . . .	2
Number of event channels . . . . .	1 actuated by external switch. 1 actuated by panel switch.
Channel width . . . . .	40 mm (40 divisions).
Chart Supply . . . . .	150 feet
Chart speeds . . . . .	1, 5, 25, and 125 mm per second
Chart speed regulation . . . . .	Synchronous motor, direct drive.

## General Information

### Specifications

Operating Temperature range, ambient . . . . . 0°C to 55°C.

Power requirements . . . . . 105-125 volts, 60 cps, 135 watts at 115 volts.

Transistors . . . . . 1-5651 and 2-5687

Tubes . . . . . 1-6BW4, 2-12B4A, 2-12AT7, 2-12AX7

#### Input terminals

Front (signal) . . . . . Binding posts.

Rear (event Marker) . . . . . Binding posts.

### Supplies

DESCRIPTION	BRUSH PART NO.
Chart Paper, 2 channel	RA 2922-22
Electric Styli (4)	RA 2822-31
Gram Gage Assembly	227416-910
Pen Mounting Tool	126717

# CHAPTER 1 GENERAL DESCRIPTION

## 1-1 INTRODUCTION

The M6800 EXORciser (Mo8SDT) is a system development tool used in the design and development of M6800 Microcomputer Systems. The EXORciser Debug and the user's system are built around the M6800 Microcomputer Family of Parts. The M6800 Microcomputer Family of Parts are discussed in the following documents.

- M6800 Microprocessor Programming Manual
- M6800 Microprocessor Applications Manual
- M6800 Microprocessor Family of Parts Data Sheets

The EXORciser may be configured in a variety of applications and with various EXORciser options. This manual, rather than discussing every possible configuration and option, discusses only the basic EXORciser with each option except the Wirewrap and Extender Modules discussed in a supplement to this manual. The basic EXORciser is discussed in Paragraph 1-5 and the options are identified in Paragraph 1-6.

This manual provides general information, installation instructions, applications information, operat-

ing procedures, and theory of operation, for Motorola's M6800 EXORciser. The M6800 EXORciser is illustrated in Figure 1-1.

## 1-2 EXORciser FEATURES

The features of the basic EXORciser include:

- Flexible, adaptable, and expandable design development tool
- Easy to use
- Provides the Microprocessing Unit capability for both the EXORciser and the user's system
- Saves system design and development time
- Decreases system design and development costs
- Evaluates and debugs the user's program
- Evaluates and debugs the user's system hardware

## 1-3 EXORciser SPECIFICATIONS

Table 1-1 identifies the basic EXORciser specifications.

TABLE 1-1. Basic EXORciser Specifications

CHARACTERISTICS	SPECIFICATIONS
Power Requirements	95-135 205-250 VAC 47-420 Hz, 250W
Word Size	
Data	8 bits
Address	16 bits
Instructions	8, 16, or 24 bits
Memory Size	64k bytes
Instruction Set	72 variable length instructions
Clock cycle time	Selectable, 1μs or an external clock between 1 and 10 μs
Interrupt	Maskable real mode interrupt
Physical Characteristics	
Table top	
Length	19.25 in
Depth	17.50 in
Height	2.00 in
Rack Mountable	
Length	19.00 in
Depth	17.00 in
Height	2.00 in
Baud Rates (Switch Selectable)	110, 150, 300, 600, 1200, 2400, 4800, and 9600

## 1.4 FUNCTIONAL DESCRIPTION

The M6800 EXORciser, illustrated in Figure 1-2, is a development system built off the shelf system. It is produced by MC6800 Microcomputer Family of Parts. The EXORciser may be easily tailored to meet the user's needs in the design and development of his system. Its modular design reduces the time required to develop a system and at the same time, provides great flexibility in configuring an emulation (functional representation) of the user's system. The EXORciser's EXbug Firmware, through its debug and program control features, minimizes the time required to develop user's systems. The EXbug firmware provides the EXORciser with the capability to:

- Display the contents of the MPU registers
- Step through the program
- Trace through user's programs to locate problem areas
- Stop the program on a selected program step
- Provide an oscilloscope trigger signal on a selected program step
- Abort from the user's program and return to the EXbug control program on command
- Reinitialize the EXORciser on command
- Change the contents of memory

The user communicates with the EXORciser in one of two ways:

- Through a RS 232C or TTY terminal
- Through the EXORciser front panel controls and indicators

The terminal device permits the user to communicate directly with the EXORciser's EXbug Firmware. The EXORciser's front panel permits the user to apply power to the EXORciser, to abort (exit) the EXORciser from a routine, and to initialize and restart the EXORciser. The EXORciser's unique front panel was designed to incorporate future EXORciser options such as data keys and displays.

## 1-5 BASIC EXORciser SYSTEM (FIGURES 1-2 AND 1-3)

The basic EXORciser (M68SDT) consists of the MPU Module, the Debug Module, the Baud Rate Module, the Power Supply, and the chassis. These modules are built aboard the M6800 Microcomputer Family of Parts (MC6800 Microprocessing Unit, MC6820 Peripheral Interface Adapter, MC6850 Asynchronous Communications Interface Adapter, MCM6810 Random Access Memory, and MCM6830 Read Only Memory devices).

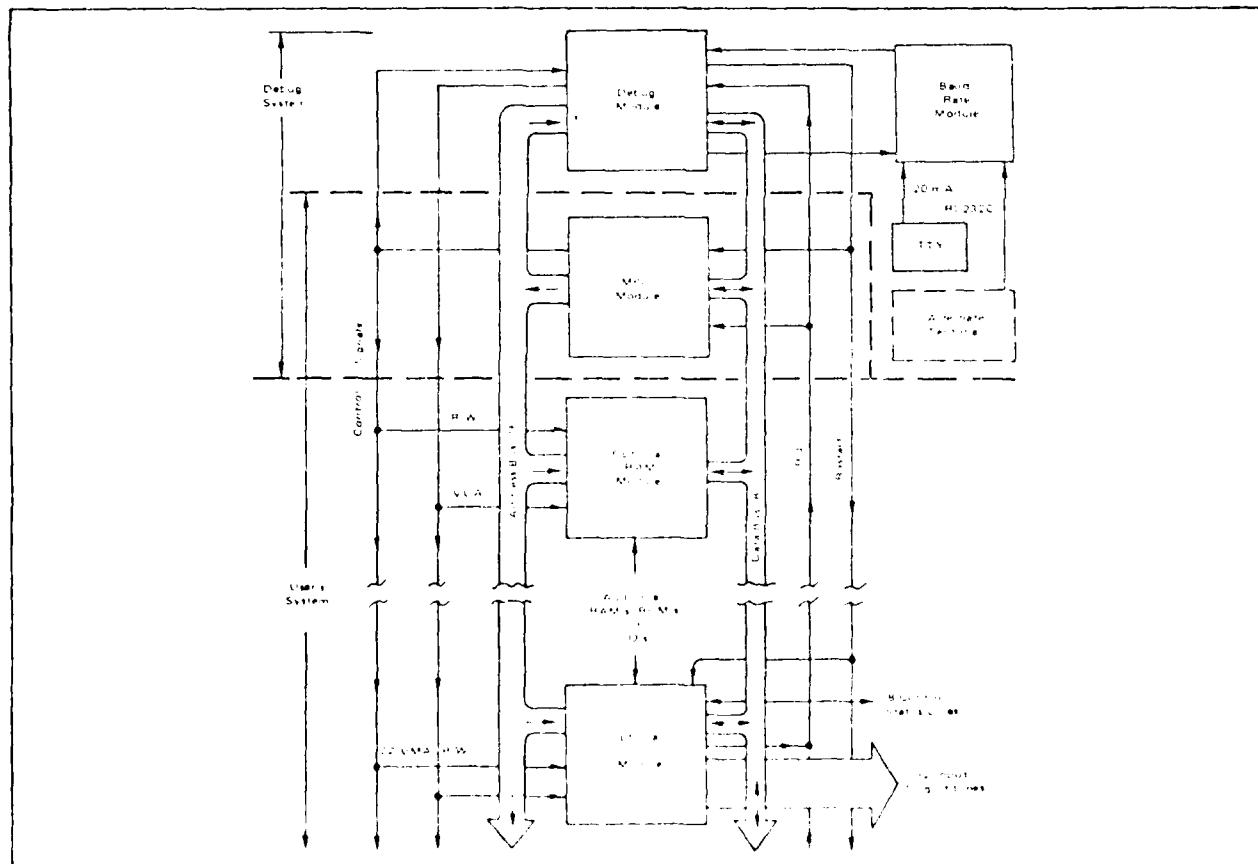


FIGURE 1-2 EXORciser Simplified Block Diagram

The MPU Module (Figure 1-4) incorporates the MC6800 Microcomputer Unit (MPU) and the system clock. This module, as illustrated in Figure 1-2, serves a dual function in that the module provides the MPU and clock for the EXORciser, the EXORciser ROM, and the memory system. The MPU Module also initiates an EXORciser restart operation and initializes the EXORciser. The MC6800 Microprocessing Unit is an 8-bit parallel device capable of addressing 64k bytes of memory. In addition,

the MPU addresses its input and output devices as memory. The MPU also provides the EXORciser with 72 variable length instructions and the capability of sequential line to read the EXORciser ROM. The EXORciser ROM is a 128 word by 8 bit ROM which contains the software required to run the EXORciser. The EXORciser can be supplied with either a 1 MHz or 2 MHz system (delayed); the MPU Module appears exactly like a MC6800 Microprocessing Unit with unlimited TTL bus drive capability.

FIGURE 1-3. Basic EXORciser System

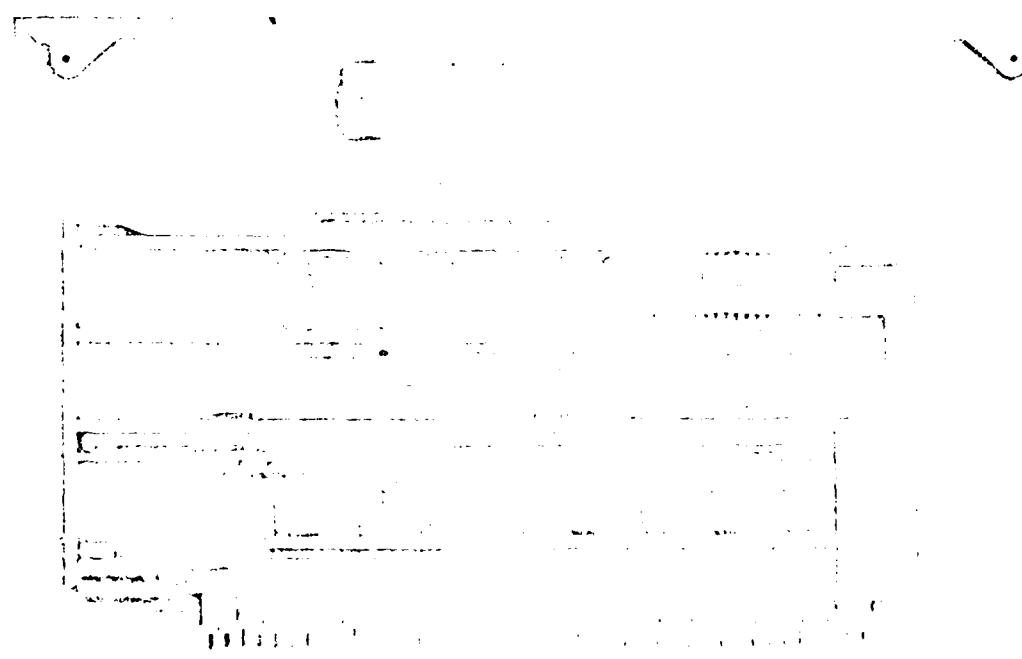


FIGURE 1-4. MPU Module

Vita

Melvin D. Townsend was born 5 April 1952 in Colorado Springs, Colorado. Upon graduation from Gen. William Mitchell High School in 1970, he enrolled in undergraduate engineering study at Colorado State University. In December, 1974, he was awarded the BSEE degree and commissioned as a second lieutenant, USAF, via AFROTC. Entering active duty immediately, he was assigned to the Foreign Technology Division (AFSC), WPAFB, Ohio. After three and one half years working technical intelligence engineering, he was transferred PCA to the AF Avionics Laboratory (AFSC) where he performed R&D engineering on electro-optical target acquisition systems. In June, 1979, he entered the Air Force Institute of Technology. Captain Townsend is a member of Eta Kappa Nu and IEEE.

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20. ABSTRACT (Continue on reverse side if necessary and identify by block number)  EKG data compression techniques were investigated for potential real time implementation on an 8 bit Motorola 6800 microprocessor. Research indicated entropy reduction transform technique such as the Fast Fourier Transform and the discrete Karhunen-Loeve Transform were not feasible for implementation on the 6800. Two redundancy reduction (RR) techniques (TOLAM and DQRR) utilizing 2nd order difference		

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operations in conjunction with variable length encoding were studied in detail. One such technique (GZIP) was found to yield significant results in EKG data. Analysis revealed compression ratios ranging from 1.25:1 to 2.26:1. Investigation of the poor performance of the compression algorithm showed significant degradation of the 2nd order difference "decorrelator" due to a noisy collection environment. It was concluded that real time EKG data compression is feasible on the 6800 but that time compression techniques which store a zero value sequence counter versus the value of zero are not efficient in a high noise environment.

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